## CS 598MEB

## Computational Cancer Genomics

## Lecture 2

## Mohammed El-Kebir

January 20, 2022


## Course Information

Course website:

- https://www.el-kebir.net/teaching/CS598MEB/Spring 2022/CS598MEB.htm|

Piazza: (please sign up)

- https://piazza.com/illinois/spring2022/cs598meb


## Description:

- This course focuses on recent algorithmic methods in cancer genomics, including somatic variant calling, phylogeny inference and identification of driver mutations. Students will study the underlying principles of these methods and the application of these methods to cancer genomics data.


## Lecture Outline

- Recap
- Maximum Parsimony
- Two-state Perfect Phylogeny
- Two-state Perfect Phylogeny Mixtures


## Reading

- Lecture notes


## Hallmarks of Cancer

## Inter-tumor heterogeneity: Every tumor is different!



Figure 1. Acquired Capabilities of Cancer
We suggest that most if not all cancers have acquired the same set of functional capabilities during their development, albeit through various mechanistic strategies.

## Cancer is Caused by Somatic Mutations



## Inter-tumor heterogeneity: Every tumor is different!



## Tumorigenesis: Cell Mutation, Division \& Migration

Clonal Evolution Theory of Cancer [Nowell, 1976]



Intra-Tumor Heterogeneity


Phylogenetic Tree $\boldsymbol{T}$

Intra-tumor heterogeneity: Every tumor cell is different

## Phylogenies are Key to Understanding Cancer



These downstream analyses critically rely on accurate tumor phylogeny inference

## Key challenge in phylogenetics:

Accurate phylogeny inference from data at present time

## Lecture Outline

- Recap
- Maximum Parsimony
- Two-state Perfect Phylogeny
- Two-state Perfect Phylogeny Mixtures


## Reading

- Lecture notes


## Character-Based Tree Reconstruction

- Characters may be morphological features
- Shape of beak \{generalist, insect catching, ...\}
- Number of legs $\{2,3,4, .$.
- Hibernation \{yes, no\}
- Character may be nucleotides/amino acids
- $\{\mathrm{A}, \mathrm{T}, \mathrm{C}, \mathrm{G}\}$
- 20 amino acids
- Values of a character are called states
- We assume discrete states



## Character-Based Phylogeny Reconstruction

## Input characters

## Output optimal tree

Question: What is optimal?

Want: Optimization criterion


## Character-Based Phylogeny Reconstruction

## Input characters

## Output optimal tree

Question: What is optimal?

Want: Optimization criterion

Question: How to optimize this criterion?

Want: Algorithm


## Character-Based Phylogeny Reconstruction: Input

| Characters / states | State 1 | State 2 |
| :--- | :--- | :--- |
| Mouth | Smile | Frown |
| Eyebrows | Normal | Pointed |



Character-Based Phylogeny Reconstruction: Criterion


Question: Which tree is better?

## Character-Based Phylogeny Reconstruction: Criterion


(a) Parsimony Score $=3$

(b) Parsimony Score $=2$

Parsimony: minimize number of changes on edges of tree

## Why Parsimony?

- Ockham's razor: "simplest" explanation for data
- Assumes that observed character differences resulted from the fewest possible mutations
- Seeks tree with the lowest parsimony score, i.e. the sum of all (costs of) mutations in the tree.



## Binary Characters

Characters

| $\begin{aligned} & \mathrm{A} \\ & \stackrel{\sim}{\mathscr{O}} \mathrm{~B} \\ & \stackrel{\mathrm{D}}{\sim} \mathrm{n} \end{aligned}$ | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 1 | 1 | 0 | 0 |
|  | 0 | 0 | 1 | 1 | 0 |
|  | 1 | 1 | 1 | 1 | 0 |
| D | 1 | 1 | 0 | 1 | 1 |

## Characters only have two possible states

Possible Encoding:<br>0 : not-mutated<br>1 : mutated

## A Small and a Large Problem

## Small Maximum Parsimony Phylogeny Problem:

Given $m \times n$ matrix $A=\left[a_{i, j}\right]$ and tree $T$ with $m$ leaves, find assignment of character states to each internal vertex of $T$ with minimum parsimony score.

Large Maximum Parsimony Phylogeny Problem:<br>Given $m \times n$ matrix $A=\left[a_{i, j}\right]$, find a tree $T$ with $m$ leaves labeled according to $A$ and an assignment of character states to each internal vertex of $T$ with minimum parsimony score.

## A Small and a Large Problem

## Small Maximum Parsimony Phylogeny Problem:

Given $m \times n$ matrix $A=\left[a_{i, j}\right]$ and tree $T$ with $m$ leaves, find assignment of character states to each internal vertex of $T$ with minimum parsimony score.

| Large Maximum Parsimony Phylogeny Problem: |
| :---: |
| Given $m \times n$ matrix $A=\left[a_{i, j}\right]$, find a tree $T$ with $m$ leaves labeled |
| according to $A$ and an assignment of character states to each internal |
| vertex of $T$ with minimum parsimony score. |

Question: Are both problems easy (i.e. in P)?

## Large Maximum Parsimony Phylogeny

- This problem is NP-hard
- Heuristics using local search (tree moves)

1. Start with an arbitrary tree $T$.
2. Check "neighbors" of $T$.
3. Move to a neighbor if it provides the best improvement in parsimony/likelihood score.


## Local Search: Nearest-Neighbor Interchange (NNI)



Rearrange four subtrees
defined by one internal edge

## Phylogenies are Key to Understanding Cancer



These downstream analyses critically rely on accurate tumor phylogeny inference

## Key challenge in phylogenetics:

Accurate phylogeny inference from data at present time

## Additional Challenge in Cancer Phylogenetics



## Additional Challenge in Cancer Phylogenetics



## Additional Challenge in Cancer Phylogenetics


normal

## Additional challenge in cancer phylogenetics: <br> Phylogeny inference from mixed bulk samples at present time

## Tumor Phylogeny Inference

Metastatic Colorectal Cancer (Patient CRC2)
[Kim et al., Clin Cancer Res 21(19), 2015]:

- 5 primary samples (P1-P5)
- 2 metastases (M1-M2)
- 412 single-nucleotide variants (SNVs)

| $n$ mutations |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ |  |  |  |
|  | ${ }^{\text {P1 }}$ (1 | 11 | 1 | 1 |  |  |  |
|  | P2 1 | 11 | 1 | 0 |  |  |  |
|  | P3 1 | 10 | 0 | 1 |  |  |  |
|  | P4 0 | 01 | 1 | 1 |  |  |  |
| $\stackrel{\cong}{E}$ | P5 0 | $0 \quad 1$ | 1 | 0 |  |  |  |
|  | m1 1 | 11 | 1 | 0 |  |  |  |
|  | M2 0 | 0 | 1 | 1 |  |  |  |
| Binary Matrix B |  |  |  |  |  |  |  |

Maximum
Parsimony

## Tumor Phylogeny Inference

## Metastatic Colorectal Cancer (Patient CRC2)

[Kim et al., Clin Cancer Res 21(19), 2015]:

- 5 primary samples (P1-P5)
- 2 metastases (M1-M2)
- 412 single-nucleotide variants (SNVs)
- 41 mutate more than once (homoplasy)


Maximum
Parsimony


## Heuristic for Tumor Phylogeny Inference

## Metastatic Colorectal Cancer (Patient CRC2)

[Kim et al., Clin Cancer Res 21(19), 2015]:

- 5 primary samples (P1-P5)
- 2 metastases (M1-M2)
- 412 single-nucleotide variants (SNVs)
- 41 mutate more than once (homoplasy)


Resulting sample tree is not representative of the
 division/mutation history or the migration history

## Lecture Outline

- Recap
- Maximum Parsimony
- Two-state Perfect Phylogeny
- Two-state Perfect Phylogeny Mixtures


## Reading

- Lecture notes


## Somatic Mutations and Cancer

## Clonal theory of cancer (Nowell, 1976)



[^0]
## Somatic Mutations and Cancer



## Progression of Somatic Mutations


$0=$ normal
$1=$ mutated


Root is the normal, founder cell and leaves are cells in tumor.

## Progression of Somatic Mutations

## Single nucleotide mutation

... CGTAATTAG ...

CGTCATTAG ...
$0=$ normal
$1=$ mutated


Root is the normal, founder cell and leaves are cells in tumor.
Infinite sites assumption: each locus mutates only once.

## Infinite Sites Model



## Two-state Perfect Phylogeny

Matrix $M \in\{0,1\}^{n \times m}$ has $n$ taxa and $m$ characters

- Taxon $f$ has state 1 for character $c$ $\Leftrightarrow f$ possesses character $c$

|  | $c_{1}$ | $c_{2}$ | $c_{3}$ | $c_{4}$ | $c_{5}$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $r_{1}$ | 1 | 1 | 0 | 0 | 0 |
| $r_{2}$ | 0 | 0 | 1 | 0 | 0 |
| $r_{3}$ | 1 | 1 | 0 | 0 | 1 |
| $r_{4}$ | 0 | 0 | 1 | 1 | 0 |
| $r_{5}$ | 0 | 1 | 0 | 0 | 0 |

## Definition

A perfect phylogeny for $M$ is a rooted tree $T$ with $n$ leaves such that:
(1) Each taxon labels only one leaf
(2) Each character labels only one edge
(3) Character possessed by a taxon are on unique path to root


## Two-state Perfect Phylogeny - Alternative Definitions

(1) Each taxon labels exactly one leaf
(2) Each character labels exactly one edge
(3) Character possessed by a taxon are on unique path to root

(1) Each taxon labels exactly one leaf
(2) Each node is labeled by $\{0,1\}^{m}$
(3) Nodes labeled with state $i$ for character $c$ form a connected subtree
(1) Each taxon labels exactly one leaf
(2) $T_{c}(i)$ is smallest subtree connecting all leaves labeled with state $i$ for character $c$
(3) $T_{c}(0)$ and $T_{c}(1)$ are disjoint for all $c$


## Two-state Perfect Phylogeny Problem

## Input:

Matrix $M \in\{0,1\}^{n \times m}$ has $n$ taxa and $m$ characters

- Taxon $f$ has state 1 for character $c$ $\Leftrightarrow f$ possesses character $c$

|  | $c_{1}$ | $c_{2}$ | $c_{3}$ | $c_{4}$ | $c_{5}$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $r_{1}$ | 1 | 1 | 0 | 0 | 0 |
| $r_{2}$ | 0 | 0 | 1 | 0 | 0 |
| $r_{3}$ | 1 | 1 | 0 | 0 | 1 |
| $r_{4}$ | 0 | 0 | 1 | 1 | 0 |
| $r_{5}$ | 0 | 1 | 0 | 0 | 0 |

## Problem

Given $M \in\{0,1\}^{n \times m}$ does $M$ have a perfect phylogeny?

## Try it yourself!

Only one of these matrices can be used to build a perfect phylogeny.
(1) As a group, decide on an approach to try to determine which one is which.
(2) Try out your approach to see if you can construct the tree.
(3) What did you learn from your attempt?

Characters

Characters

## The Perfect Phylogeny Problem - Preliminaries

## Problem

Given $M \in\{0,1\}^{n \times m}$ does $M$ have a perfect phylogeny?

## Definition

$I(c)$ is the set of taxa that possess character $c$; and $\sigma(f)$ is the set of characters possessed by taxon $f$.

|  | $c_{1}$ | $c_{2}$ | $c_{3}$ | $c_{4}$ | $c_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $r_{1}$ | 1 | 1 | 0 | 0 | 0 |
| $r_{2}$ | 0 | 0 | 1 | 0 | 0 |
| $r_{3}$ | 1 | 1 | 0 | 0 | 1 |
| $r_{4}$ | 0 | 0 | 1 | 1 | 0 |
| $r_{5}$ | 0 | 1 | 0 | 0 | 0 |


|  |  | $c_{1}(2)$ | $c_{2}(1)$ | $c_{3}(3)$ | $c_{4}(5)$ | $c_{5}(4)$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $r_{1}$ | 1 | 1 | 0 | 0 | 0 |
|  | $r_{2}$ | 0 | 0 | 1 | 0 | 0 |
| 7 | $r_{3}$ | 1 | 1 | 0 | 1 | 0 |
|  | $r_{4}$ | 0 | 0 | 1 | 0 | 1 |
|  | $r_{5}$ | 1 | 0 | 0 | 0 | 0 |

$$
\begin{aligned}
I\left(c_{1}\right) & =\left\{r_{1}, r_{3}\right\} \\
\sigma\left(r_{1}\right) & =\left\{c_{1}, c_{2}\right\}
\end{aligned}
$$

Sort columns of $M$ s.t. $c<d$ iff $|I(c)| \geq|I(d)|$. Break ties arbitrarily.

- Consider rows of $M$ iteratively
- $T_{i}$ is tree of first $i$ rows of $M$
- $T_{1}$ is a path graph
- Terminal nodes $r$ and 1
- $|\sigma(1)|+1$ edges labeled by $\sigma(1)$

$$
c<d \text { iff }|I(c)| \geq|I(d)|
$$

|  | $c_{1}$ | $c_{2}$ | $c_{3}$ | $c_{4}$ | $c_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $r_{1}$ | 1 | 1 | 0 | 0 | 0 |
| $r_{2}$ | 0 | 0 | 1 | 0 | 0 |
| $r_{3}$ | 1 | 1 | 0 | 1 | 0 |
| $r_{4}$ | 0 | 0 | 1 | 0 | 1 |
| $r_{5}$ | 1 | 0 | 0 | 0 | 0 |

- Consider rows of $M$ iteratively

$$
c<d \text { iff }|I(c)| \geq|I(d)|
$$

- $T_{i}$ is tree of first $i$ rows of $M$
- $T_{1}$ is a path graph
- Terminal nodes $r$ and 1
- $|\sigma(1)|+1$ edges labeled by $\sigma(1)$
- $T_{i+1}$ is a supertree of $T_{i}$
- Let $v$ be last node on walk from $r$ matching characters $\sigma(i+1)$
$\star$ Character $d$ is the last match
* Unmatched characters $\tau(i+1)$

|  | $c_{1}$ | $c_{2}$ | $c_{3}$ | $c_{4}$ | $c_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $r_{1}$ | 1 | 1 | 0 | 0 | 0 |
| $r_{2}$ | 0 | 0 | 1 | 0 | 0 |
| $r_{3}$ | 1 | 1 | 0 | 1 | 0 |
| $r_{4}$ | 0 | 0 | 1 | 0 | 1 |
| $r_{5}$ | 1 | 0 | 0 | 0 | 0 |

- Consider rows of $M$ iteratively

$$
c<d \text { iff }|I(c)| \geq|I(d)|
$$

- $T_{i}$ is tree of first $i$ rows of $M$
- $T_{1}$ is a path graph
- Terminal nodes $r$ and 1
- $|\sigma(1)|+1$ edges labeled by $\sigma(1)$
- $T_{i+1}$ is a supertree of $T_{i}$
- Let $v$ be last node on walk from $r$ matching characters $\sigma(i+1)$
$\star$ Character $d$ is the last match
* Unmatched characters $\tau(i+1)$
- Extend $T_{i}$ with path $\Pi$
$\star \Pi$ has terminals $v$ and $i+1$
$\star \Pi$ has $|\tau(i+1)|+1$ edges labeled by $\tau(i+1)$

|  | $c_{1}$ | $c_{2}$ | $c_{3}$ | $c_{4}$ | $c_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $r_{1}$ | 1 | 1 | 0 | 0 | 0 |
| $r_{2}$ | 0 | 0 | 1 | 0 | 0 |
| $r_{3}$ | 1 | 1 | 0 | 1 | 0 |
| $r_{4}$ | 0 | 0 | 1 | 0 | 1 |
| $r_{5}$ | 1 | 0 | 0 | 0 | 0 |



- Consider rows of $M$ iteratively

$$
c<d \text { iff }|I(c)| \geq|I(d)|
$$

- $T_{i}$ is tree of first $i$ rows of $M$
- $T_{1}$ is a path graph
- Terminal nodes $r$ and 1
- $|\sigma(1)|+1$ edges labeled by $\sigma(1)$
- $T_{i+1}$ is a supertree of $T_{i}$
- Let $v$ be last node on walk from $r$ matching characters $\sigma(i+1)$
$\star$ Character $d$ is the last match
$\star$ Unmatched characters $\tau(i+1)$
- Extend $T_{i}$ with path $\Pi$
$\star \Pi$ has terminals $v$ and $i+1$
$\star \Pi$ has $|\tau(i+1)|+1$ edges labeled by $\tau(i+1)$

|  | $c_{1}$ | $c_{2}$ | $c_{3}$ | $c_{4}$ | $c_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $r_{1}$ | 1 | 1 | 0 | 0 | 0 |
| $r_{2}$ | 0 | 0 | 1 | 0 | 0 |
| $r_{3}$ | 1 | 1 | 0 | 1 | 0 |
| $r_{4}$ | 0 | 0 | 1 | 0 | 1 |
| $r_{5}$ | 1 | 0 | 0 | 0 | 0 |



- Consider rows of $M$ iteratively

$$
c<d \text { iff }|I(c)| \geq|I(d)|
$$

- $T_{i}$ is tree of first $i$ rows of $M$
- $T_{1}$ is a path graph
- Terminal nodes $r$ and 1
- $|\sigma(1)|+1$ edges labeled by $\sigma(1)$
- $T_{i+1}$ is a supertree of $T_{i}$
- Let $v$ be last node on walk from $r$ matching characters $\sigma(i+1)$
* Character $d$ is the last match
* Unmatched characters $\tau(i+1)$
- Extend $T_{i}$ with path $\Pi$
$\star \Pi$ has terminals $v$ and $i+1$
$\star \Pi$ has $|\tau(i+1)|+1$ edges labeled by $\tau(i+1)$

|  | $c_{1}$ | $c_{2}$ | $c_{3}$ | $c_{4}$ | $c_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $r_{1}$ | 1 | 1 | 0 | 0 | 0 |
| $r_{2}$ | 0 | 0 | 1 | 0 | 0 |
| $r_{3}$ | 1 | 1 | 0 | 1 | 0 |
| $r_{4}$ | 0 | 0 | 1 | 0 | 1 |
| $r_{5}$ | 1 | 0 | 0 | 0 | 0 |



## Lemma

Let $M_{i} \in 0,1^{i \times m}$ be a submatrix of $M$. If $M$ is conflict-free then $T_{i}$ is a perfect phylogeny for $M_{i}$.

## Lecture Outline

- Recap
- Maximum Parsimony
- Two-state Perfect Phylogeny
- Two-state Perfect Phylogeny Mixtures


## Reading

- Lecture notes


## Sequencing and Tumor Phylogeny Inference



## Sequencing and Tumor Phylogeny Inference



## Tumor Phylogeny Inference: Given frequencies $\boldsymbol{F}$, find phylogeny $\boldsymbol{T}$ and proportions $\boldsymbol{U}$

## Key Challenge in Computational Biology



Translating a biological problem into a computational biology

## Perfect Phylogeny Mixture

## Assumptions:

- Infinite sites assumption:
a character changes state once
- Error-free data

Restricted PP


1-1 $\downarrow$ Equivalent


Restricted PP Matrix B
Rows of $\boldsymbol{U}$ are proportions:

$$
u_{p j} \geq 0 \text { and } \sum_{j} u_{p j} \leq 1
$$

Perfect Phylogeny Theorem
[Estabrook, 1971]
[Gusfield, 1991]

Perfect Phylogeny Mixture: [El-Kebir*, Oesper* et al., 2015] Given $\boldsymbol{F}$, find $\boldsymbol{U}$ and $\boldsymbol{B}$ such that $\boldsymbol{F}=\boldsymbol{U} \boldsymbol{B}$

## Previous Work

## Variant of PPM:

TrAp [Strino et al., 2013], PhyloSub [Jiao et al., 2014]
CITUP [Malikic et al., 2015], BitPhylogeny [Yuan et al., 2015]
LICHeE [Popic et al., 2015], ...
Restricted PP
Tree $\boldsymbol{T}$


1-1 $\downarrow$ Equivalent


Restricted PP Matrix B
Rows of $\boldsymbol{U}$ are proportions:

$$
u_{p j} \geq 0 \text { and } \sum_{j} u_{p j} \leq 1
$$

Perfect Phylogeny Theorem
[Estabrook, 1971]
[Gusfield, 1991]

Perfect Phylogeny Mixture: [El-Kebir*, Oesper* et al., 2015] Given $\boldsymbol{F}$, find $\boldsymbol{U}$ and $\boldsymbol{B}$ such that $\boldsymbol{F}=\boldsymbol{U} \boldsymbol{B}$

## Combinatorial Characterization

- Frequency $f_{p, i}$ is mass of subtree rooted at node that introduced $i$
- Usage $u_{p, i}$ is mass of node that introduced $i$


| $\bigcirc 0$ | $n$ mutations |  |  |  | $\bigcirc$ | clones |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\bigcirc$ | 0 | 0 | $\bigcirc$ |  |  | $\bigcirc$ | 6 | $\bigcirc$ | 6 |  | $\bigcirc$ |
| $\stackrel{0}{\circ} \mathrm{~s}_{1} 0.8$ | 0.8 | 0.8 | 0.0 | 0.0 | 0.0 | $\stackrel{\text { c }}{ }$ | 0.0 | 0.0 | 0. | 0.0 |  | 0.0 |
| ${ }_{\text {¢ }}^{0} \mathrm{~S}_{2} 0.7$ | 0.6 | 0.0 | 0.6 | 0.0 | 0.0 |  | 0.1 | 0.0 | 0. | 0.6 |  | 0.0 |
| $\stackrel{\sim}{\boldsymbol{E}} \mathrm{S}_{3} 0.8$ | 0.0 | 0.0 | 0.0 | 0.6 | 0.4 |  |  | 0.0 | 0. | 0.0 |  | 0.4 |
| Frequency Matrix F |  |  |  |  |  | Mixture Matrix U |  |  |  |  |  |  |



Restricted PP Matrix B

$$
\begin{aligned}
& \text { Rows of } \boldsymbol{U} \text { are proportions: } \\
& \qquad u_{p j} \geq 0 \text { and } \sum_{j} u_{p j} \leq 1
\end{aligned}
$$

Perfect Phylogeny Mixture: [El-Kebir*, Oesper* et al., 2015] Given $\boldsymbol{F}$, find $\boldsymbol{U}$ and $\boldsymbol{B}$ such that $\boldsymbol{F}=\boldsymbol{U} \boldsymbol{B}$

## Combinatorial Characterization

- Frequency $f_{p, i}$ is mass of subtree rooted at node that introduced $i$
- Usage $u_{p, i}$ is mass of node that introduced $i$



Restricted PP Matrix B

## Theorem 1:

$\boldsymbol{T}$ is a solution to the PPM if and only if $\boldsymbol{T}$ is a spanning tree of $\boldsymbol{G}$ satisfying the sum condition

Perfect Phylogeny Mixture: [El-Kebir*, Oesper* et al., 2015] Given $\boldsymbol{F}$, find $\boldsymbol{U}$ and $\boldsymbol{B}$ such that $\boldsymbol{F}=\boldsymbol{U} \boldsymbol{B}$

## Combinatorial Characterization

- Frequency $f_{p, i}$ is mass of subtree rooted at node that introduced $i$
- Usage $u_{p, i}$ is mass of node that introduced $i$



Restricted PP Matrix B

## Theorem 1:

$\boldsymbol{T}$ is a solution to the PPM if and only if $\boldsymbol{T}$ is a spanning tree of $\boldsymbol{G}$ satisfying the sum condition

Perfect Phylogeny Mixture: [El-Kebir*, Oesper* et al., 2015] Given $\boldsymbol{F}$, find $\boldsymbol{U}$ and $\boldsymbol{B}$ such that $\boldsymbol{F}=\boldsymbol{U} \boldsymbol{B}$

## Non-uniqueness of Solutions to PPM



Question 1: Can we determine the number of solutions?

Question 2: Can sample solutions uniformly at random?

## Summary of Lectures 1 \& 2

- DNA, RNA and proteins are sequences
- Central dogma of molecular biology: DNA -> RNA -> protein
- Problem != algorithm
- Key challenge in computational biology is translating a biological problem into a computational problem
- Cancer is a genetic disease caused by somatic mutations
- Inter-tumor heterogeneity and intra-tumor heterogeneity:
- Not only is every tumor different, but so is every tumor cell...
- Reading:
- "Biology for Computer Scientists" by Lawrence Hunter (http://www.el-kebir.net/teaching/CS466/Hunter_BIO_CS.pdf)


[^0]:    "typical tumor": ~10 driver mutations
    100's - 1000's of passenger mutations

