

CS 598MEB

Computational Cancer Biology

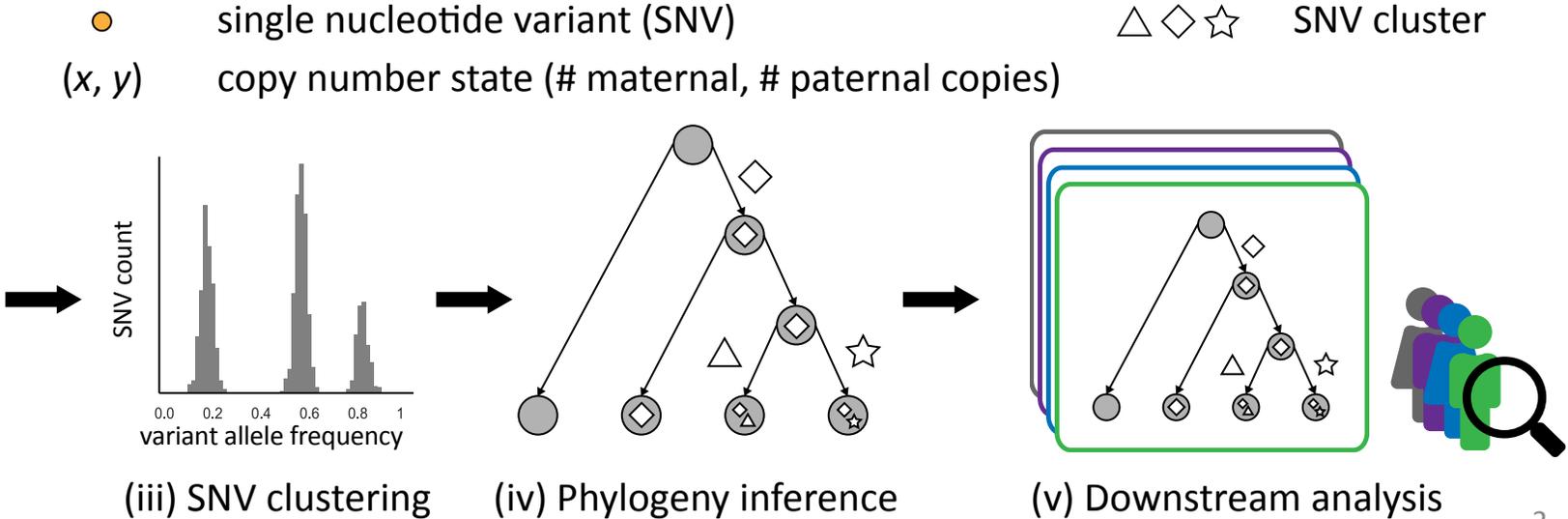
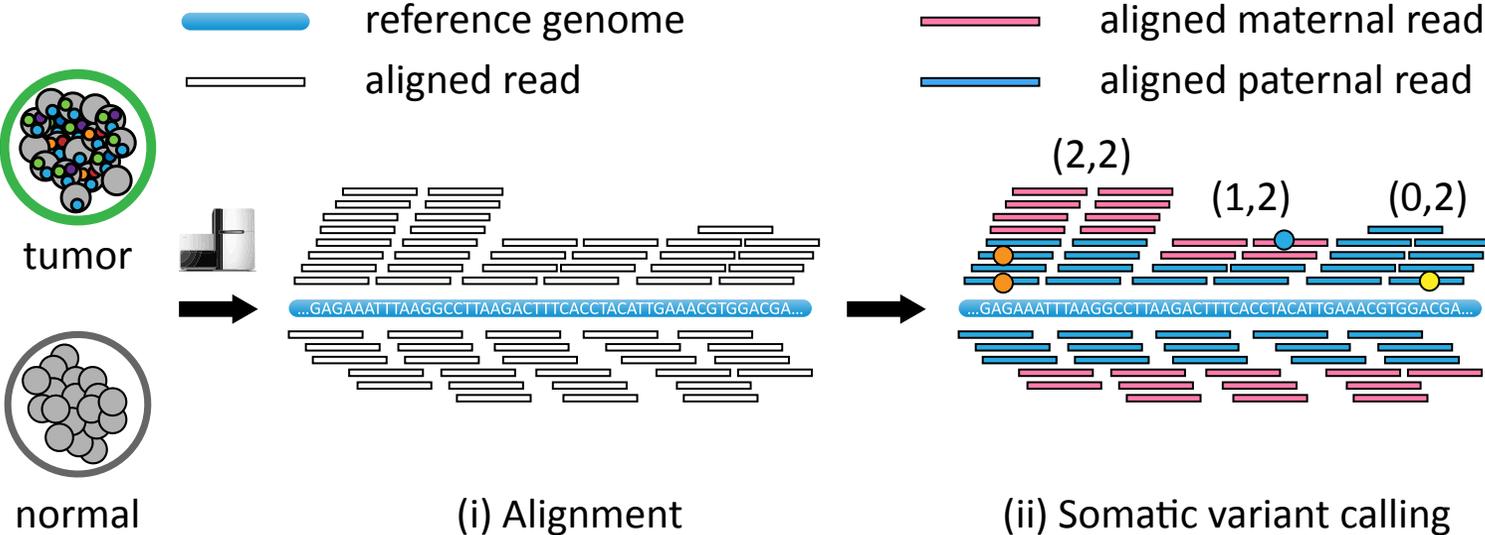
Lecture 15

Mohammed El-Kebir

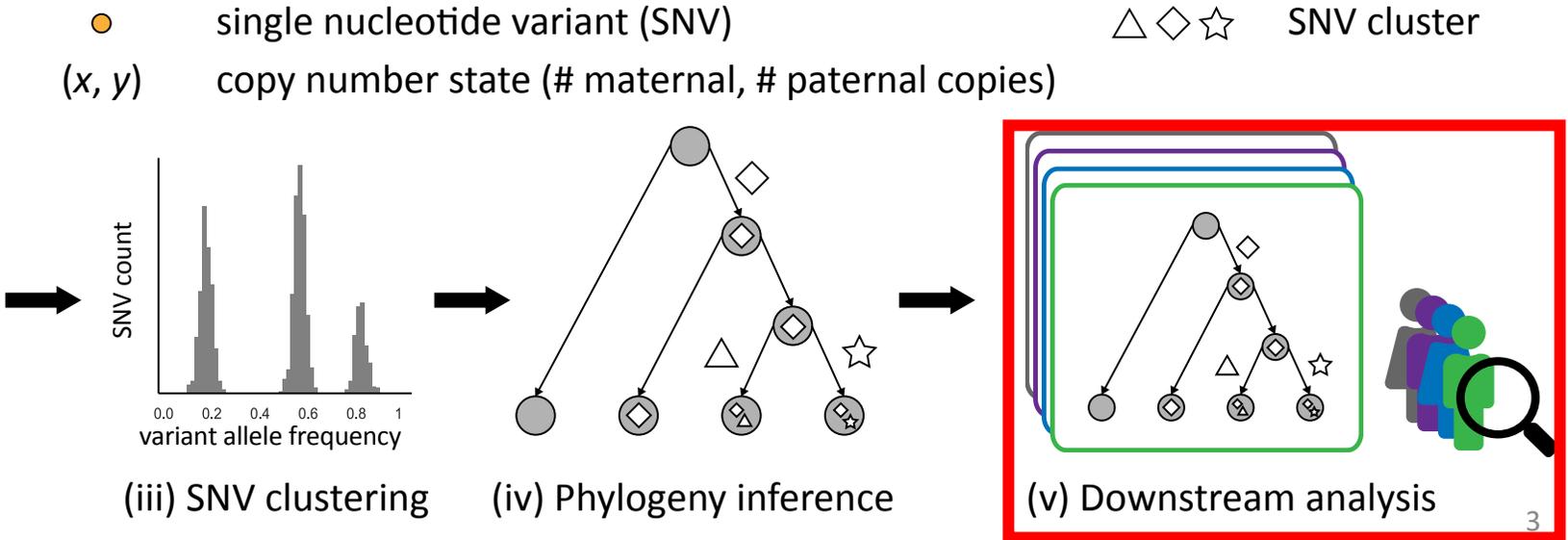
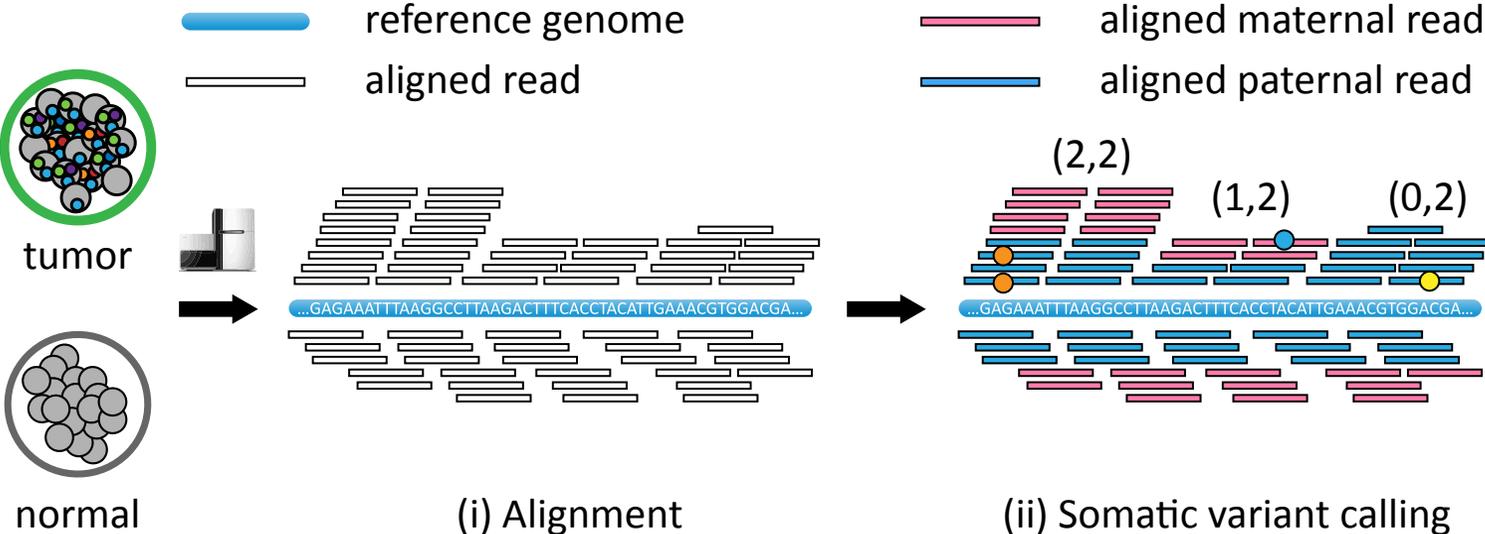
March 18, 2021



Cancer Phylogenetics Pipeline

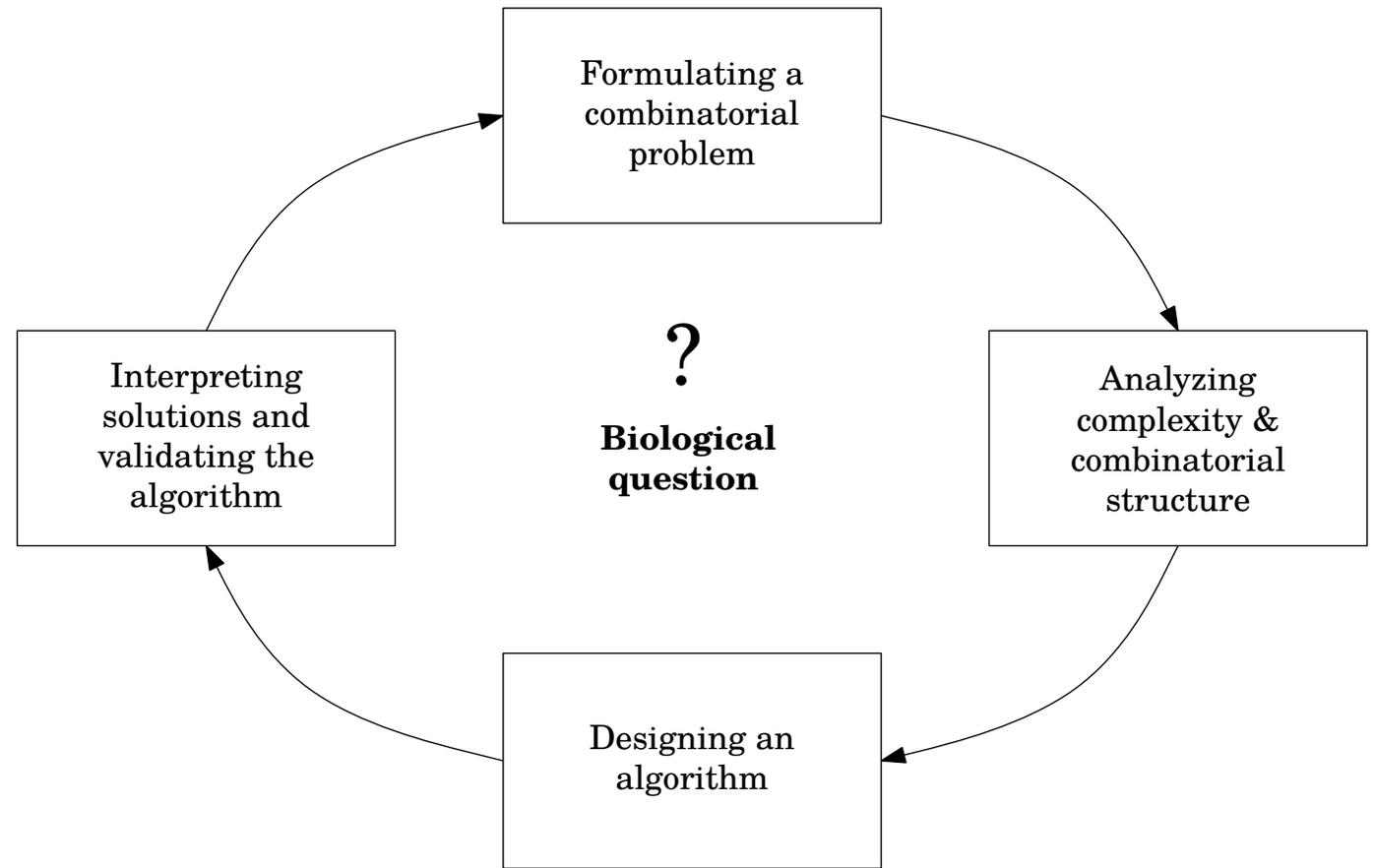


Cancer Phylogenetics Pipeline



Outline

- Metastasis
- Maximum parsimony
- Problem statement
- Complexity
- Algorithm & results
- Problem variants



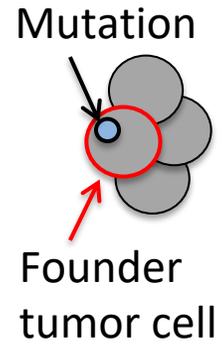
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Tumorigenesis: (i) Cell Mutation

Clonal Theory of Cancer

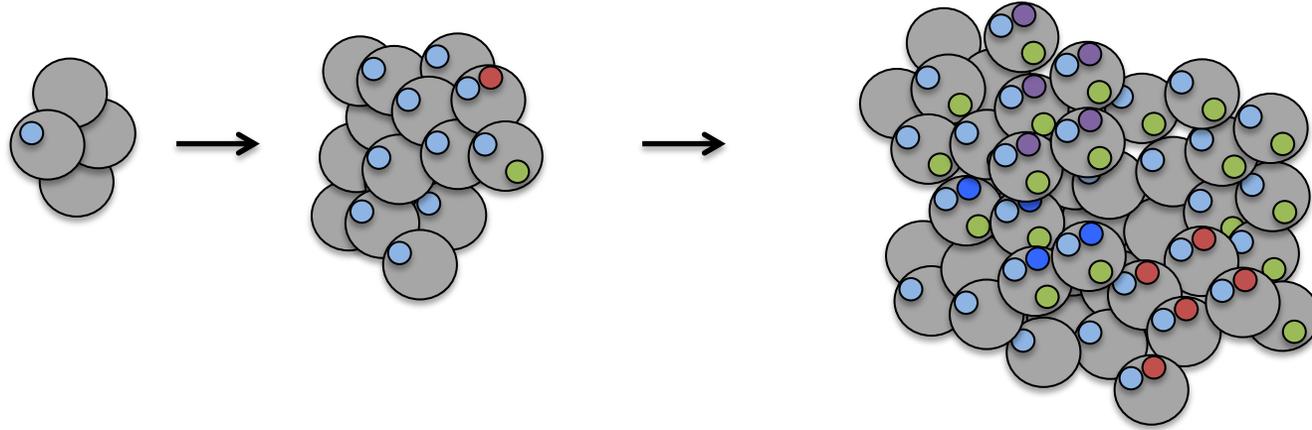
[Nowell, 1976]



Tumorigenesis: (i) Cell Mutation, (ii) Cell Division

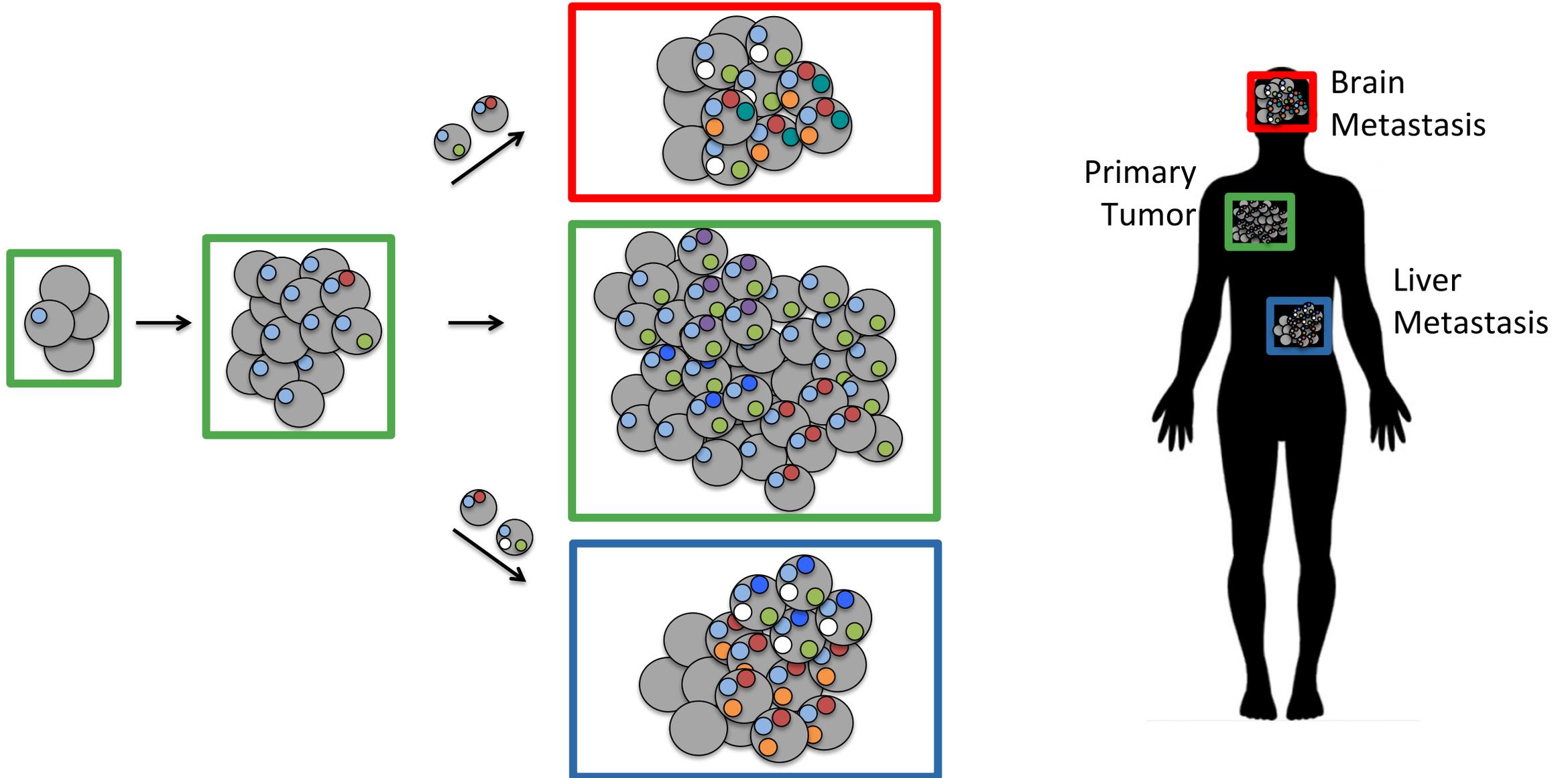
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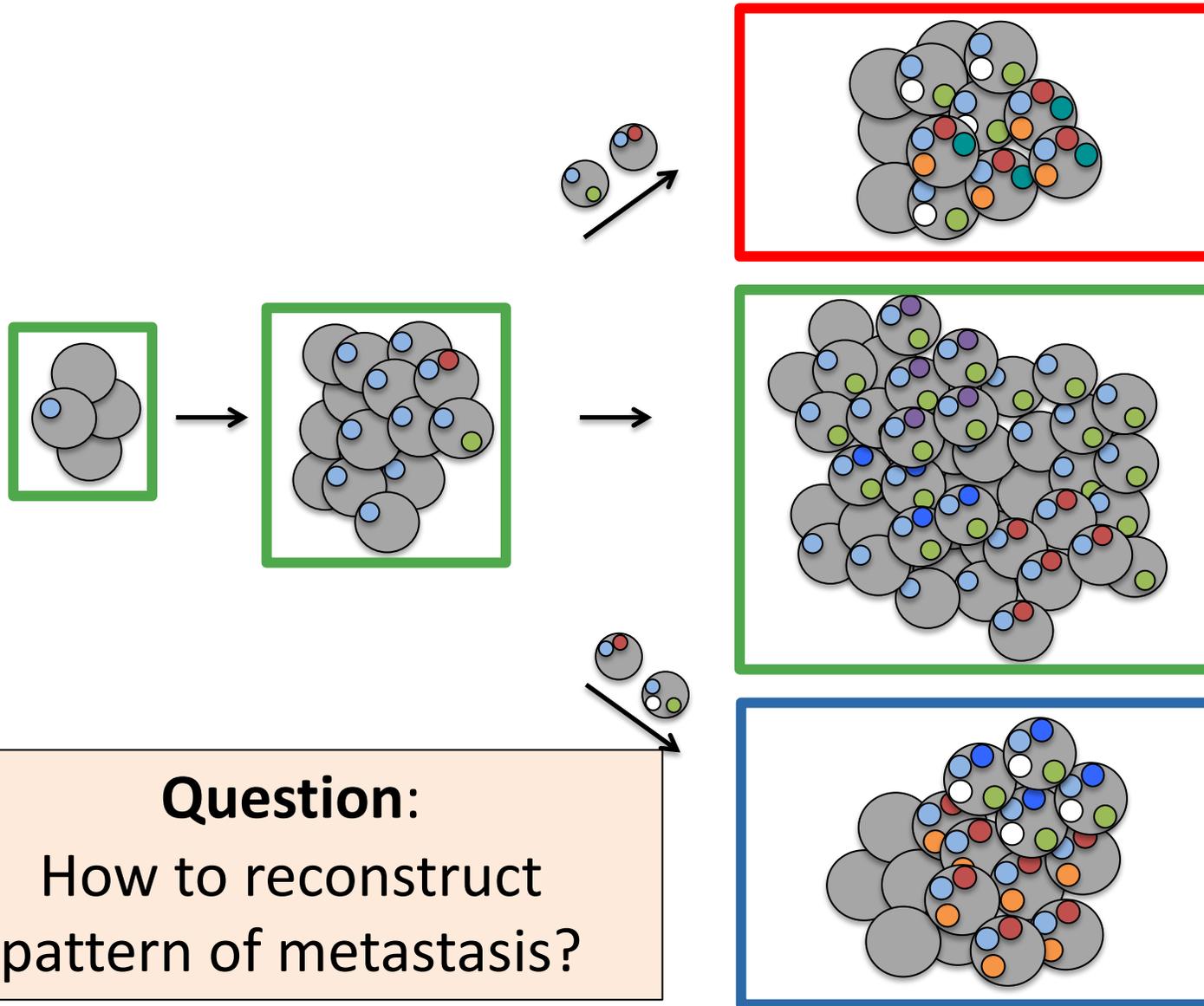


Heterogeneous Tumor

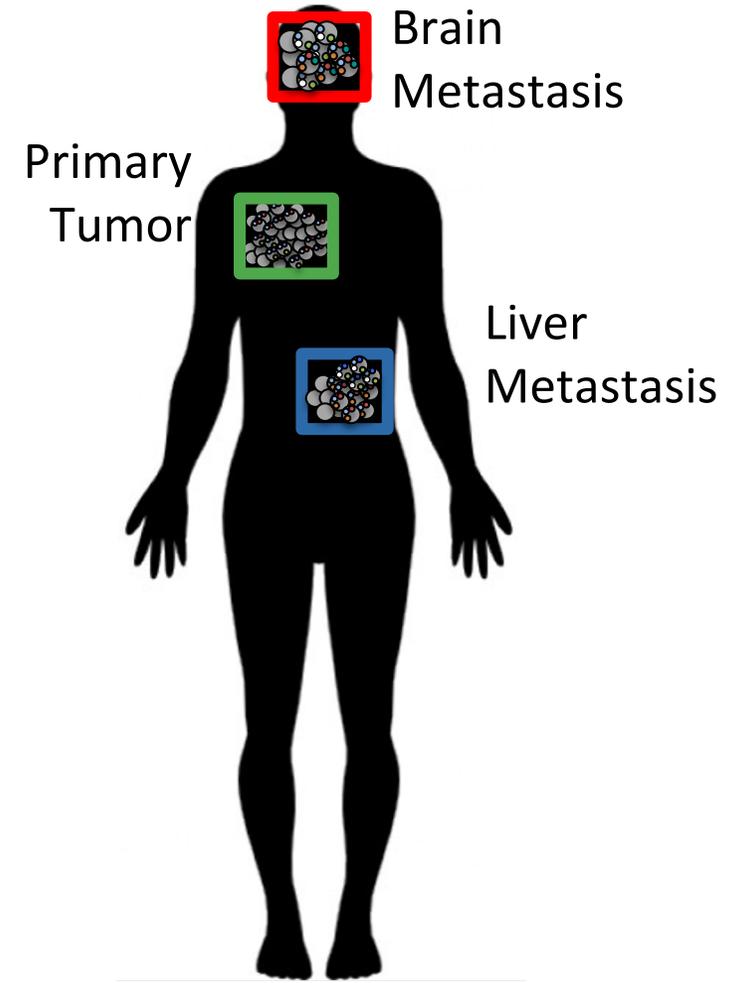
Tumorigenesis: (i) Cell Division, (ii) Mutation & (iii) Migration



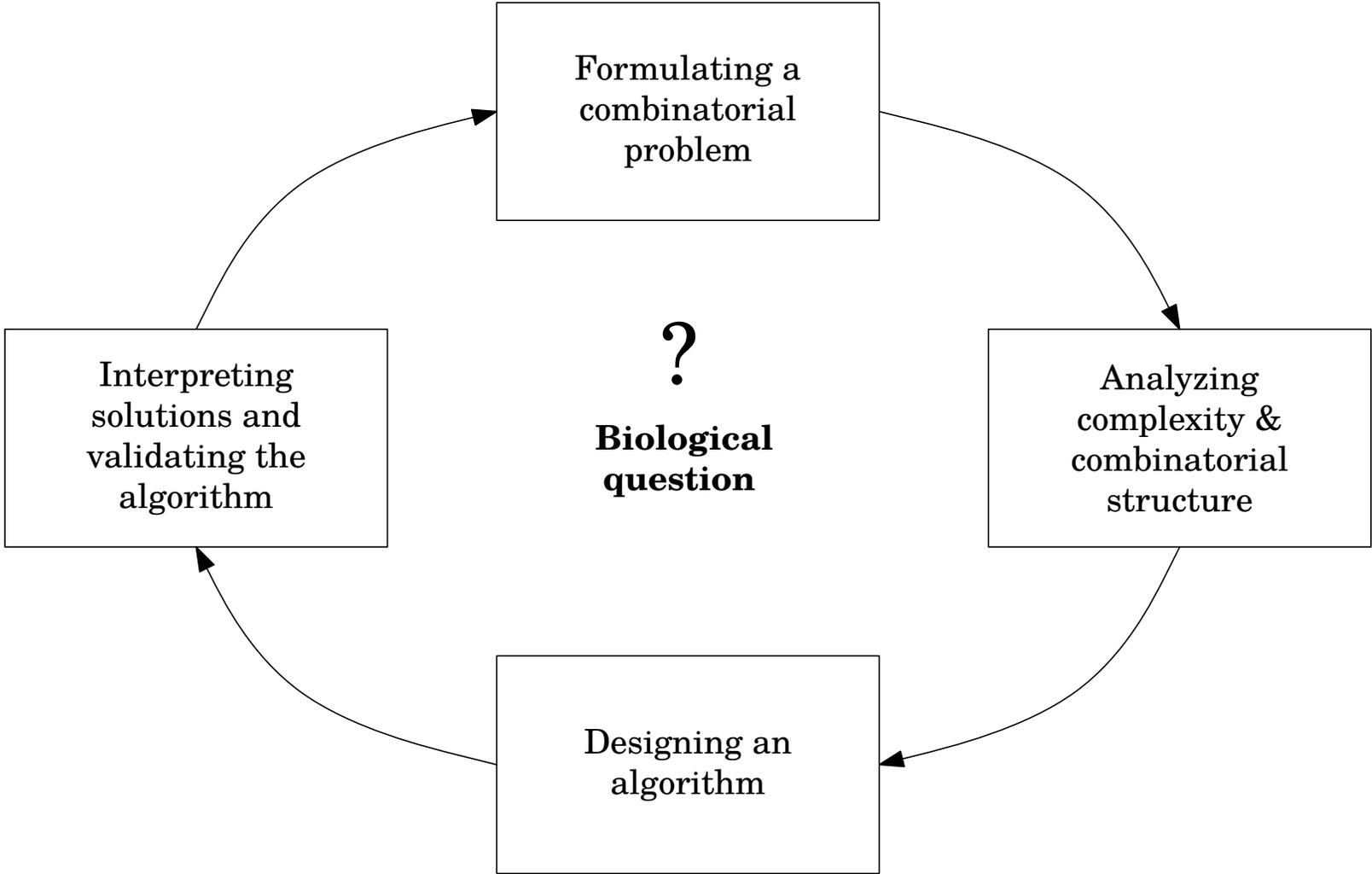
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Question:
How to reconstruct
pattern of metastasis?

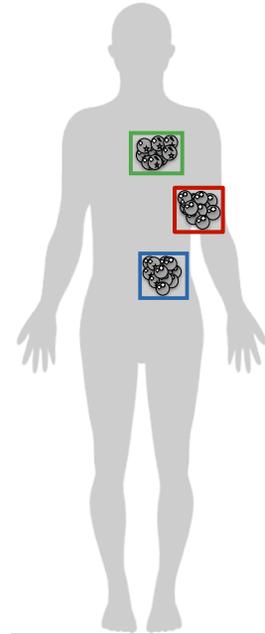
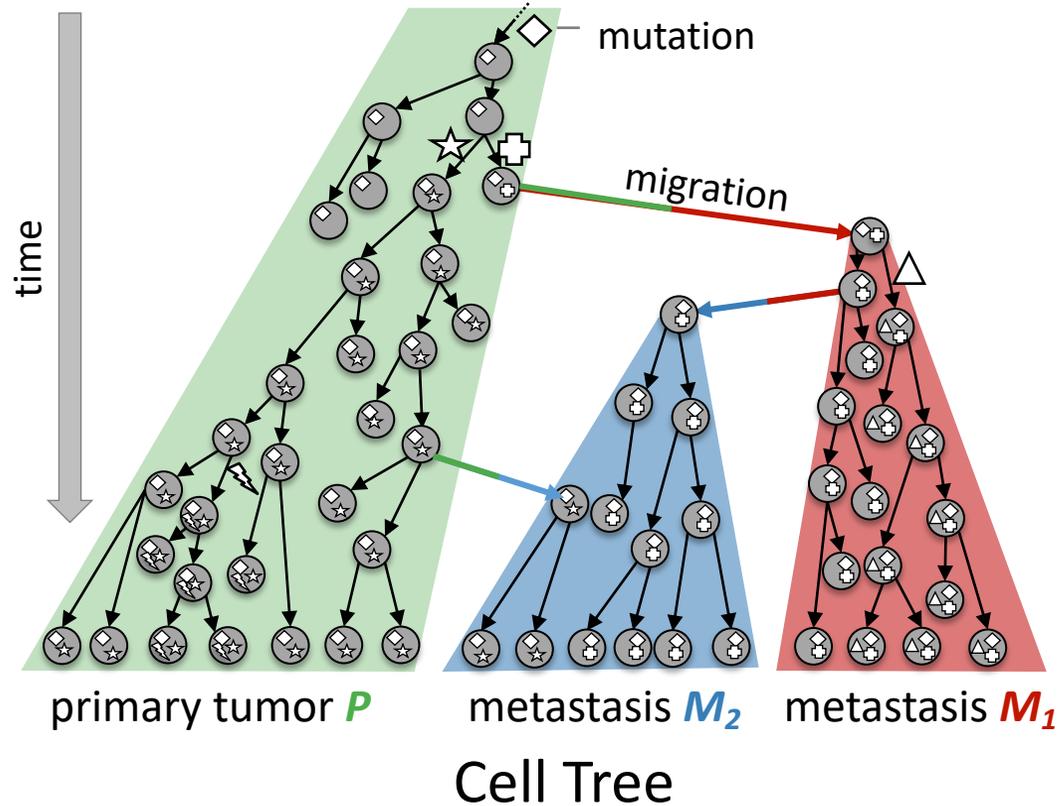


Key Challenge in Computational Biology

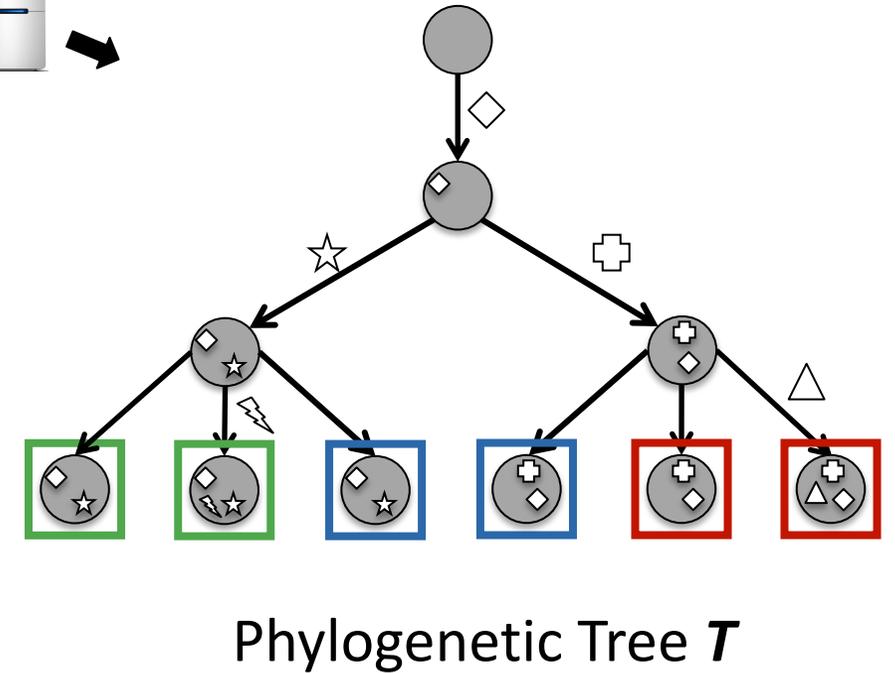
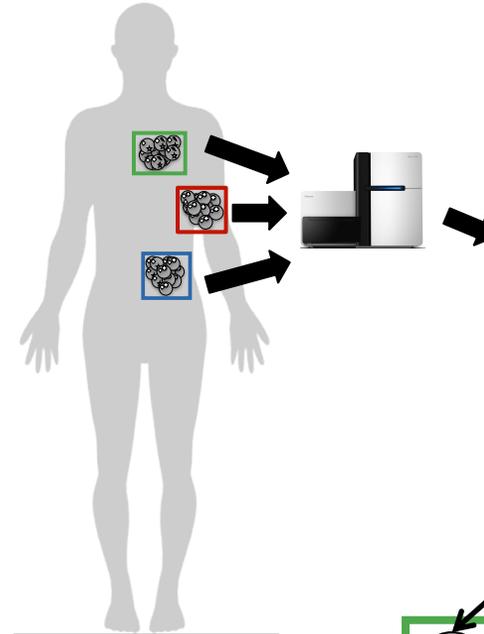
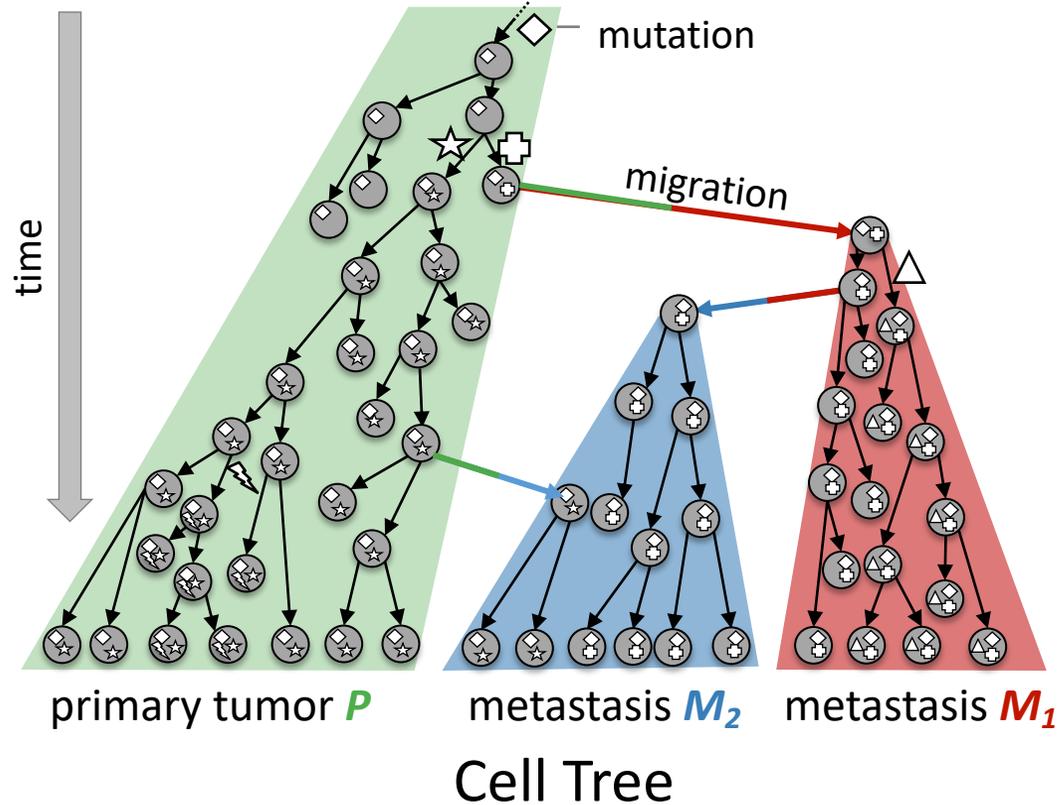


Translating a biological problem into computer science

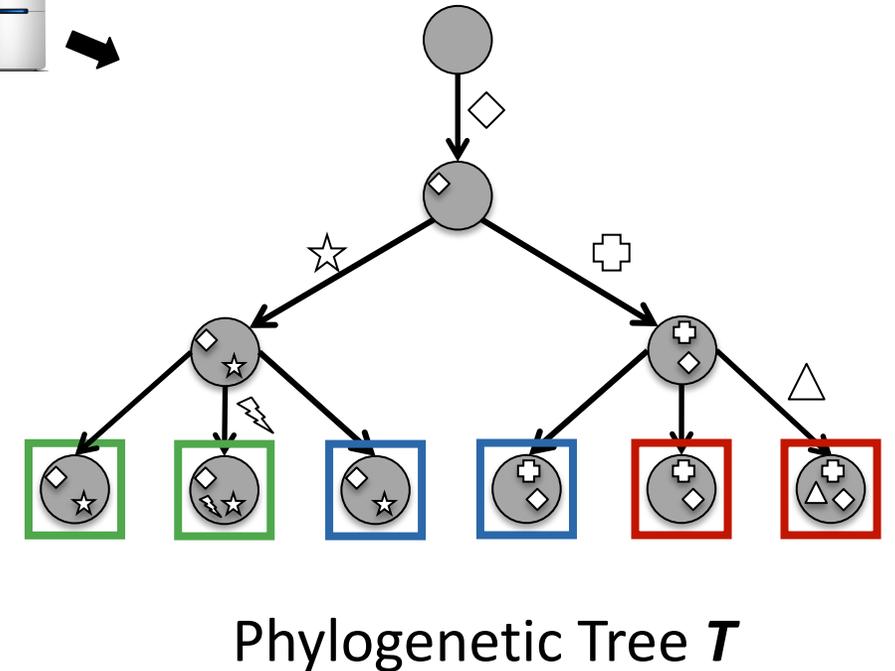
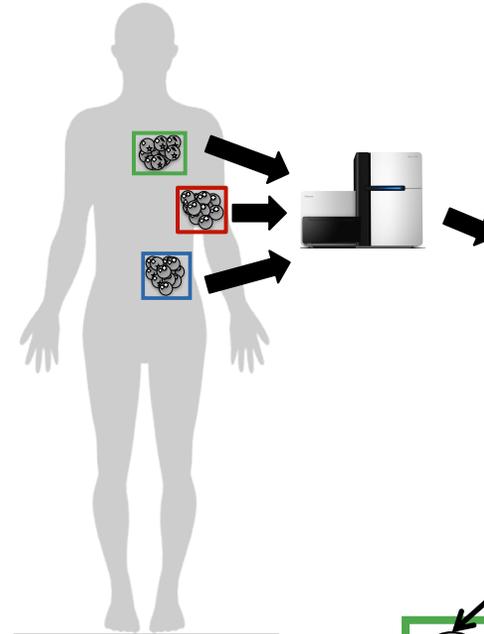
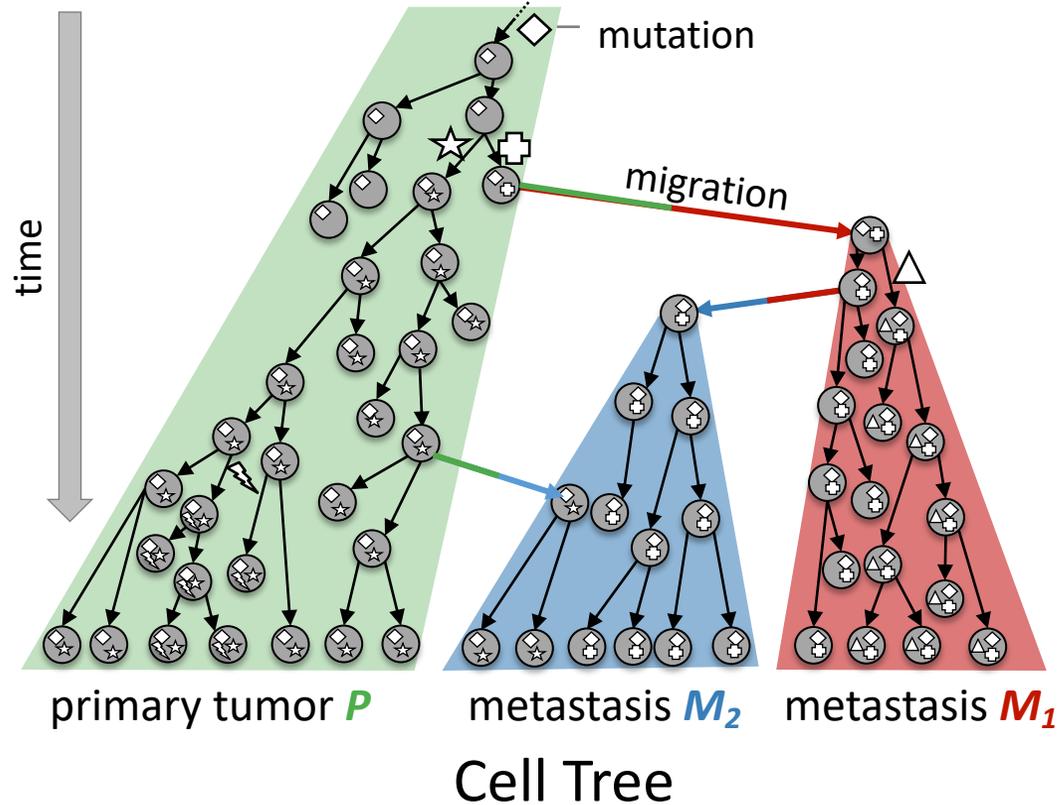
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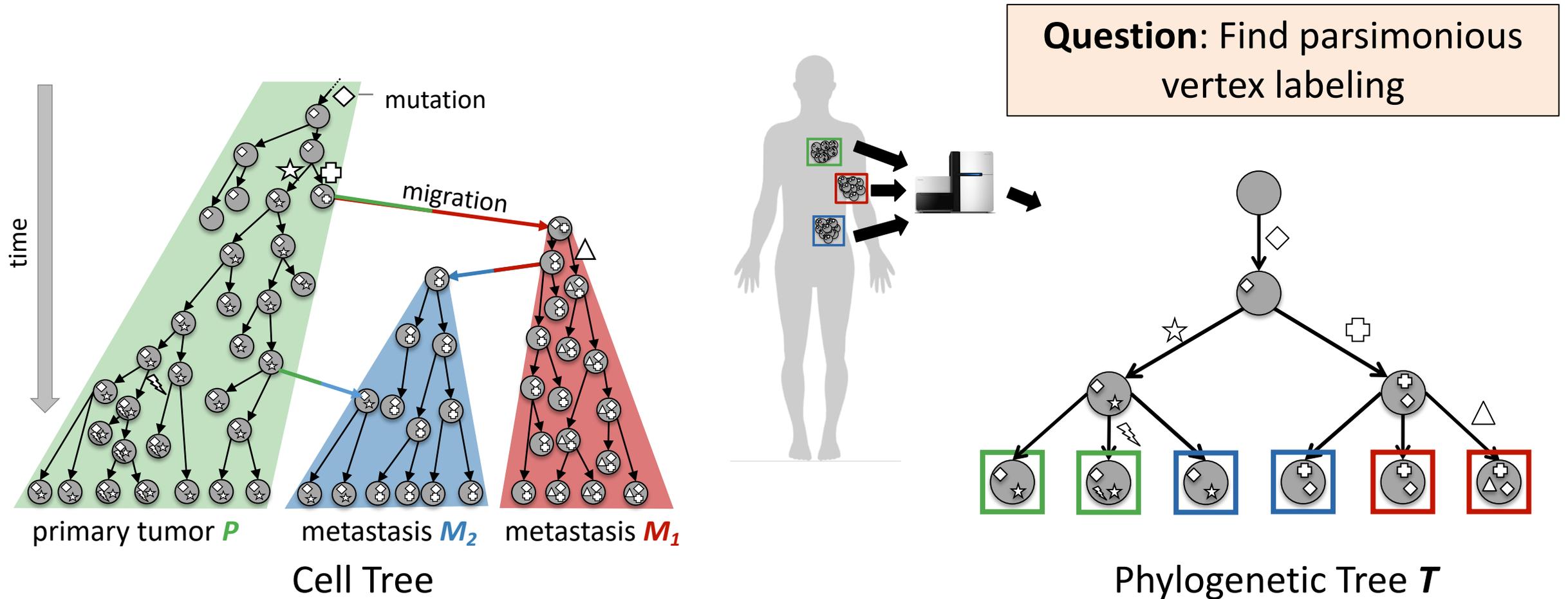


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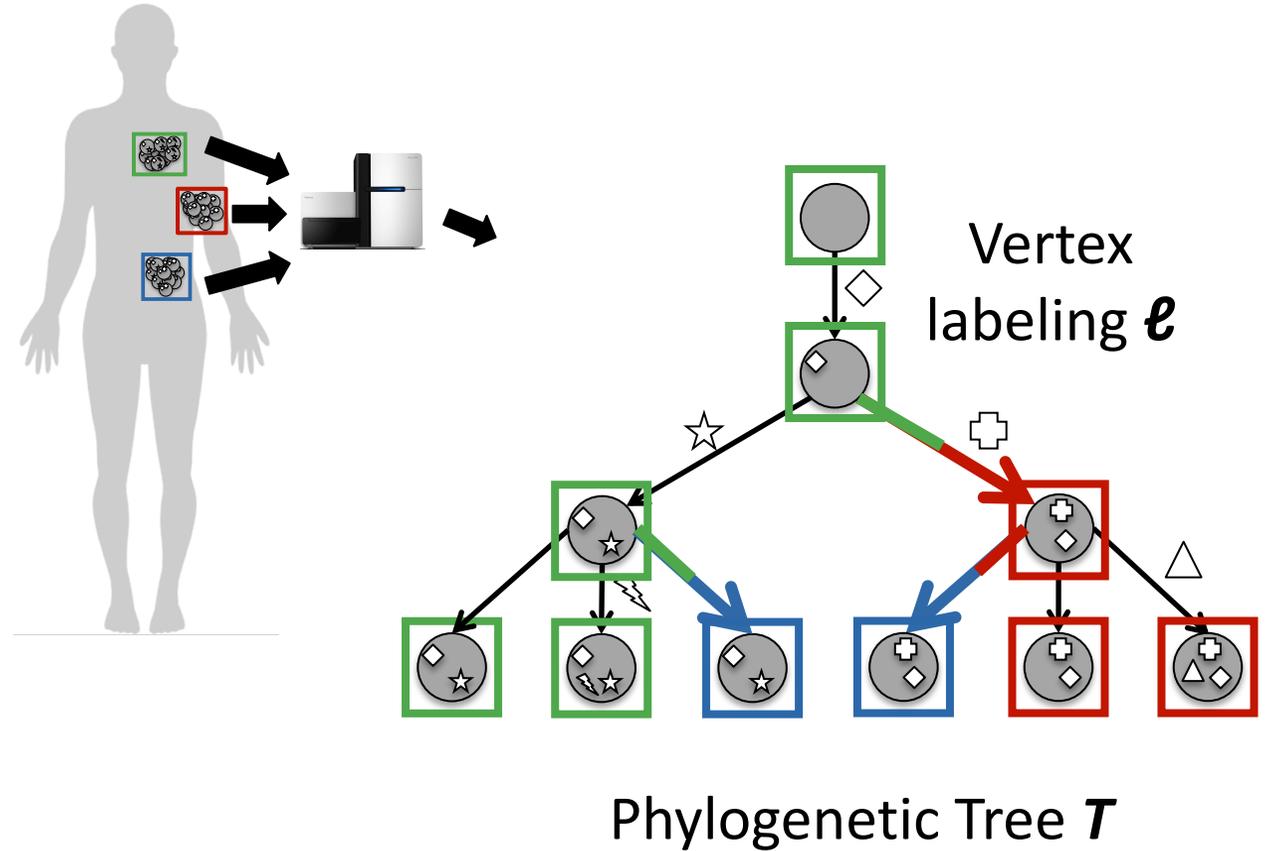
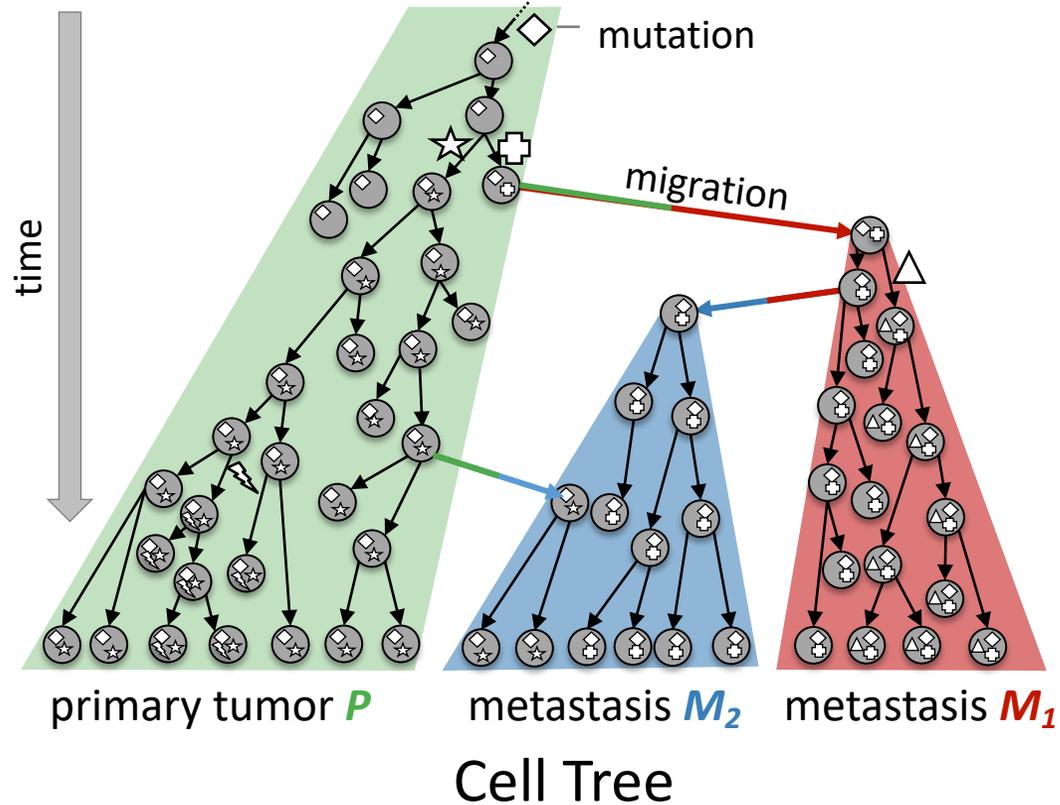
Goal: Given phylogenetic tree T , find *parsimonious* vertex labeling ℓ with fewest migrations

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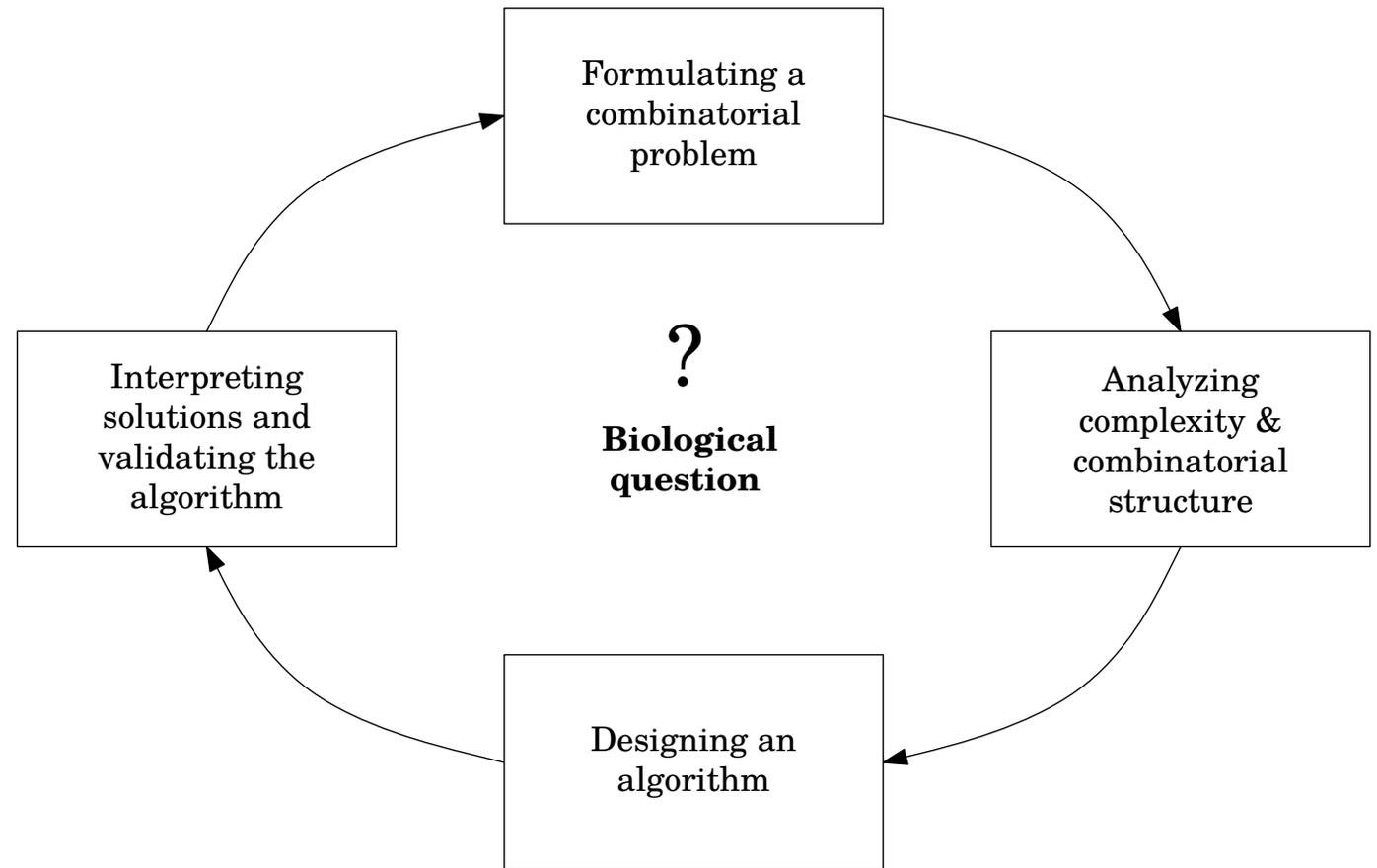
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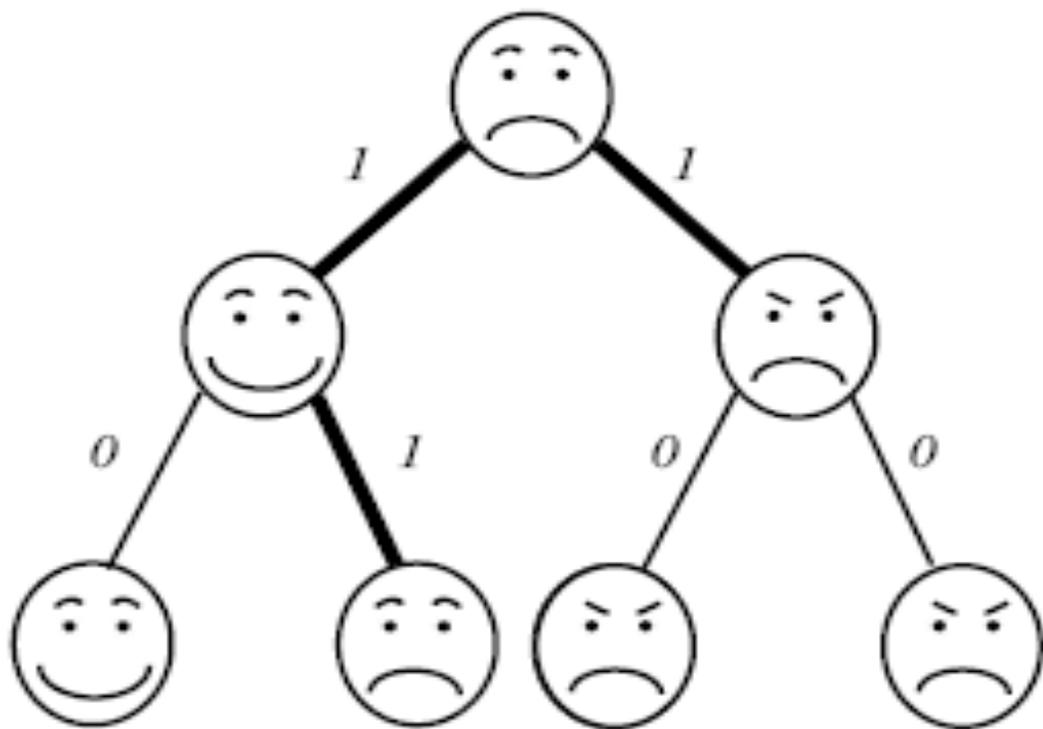
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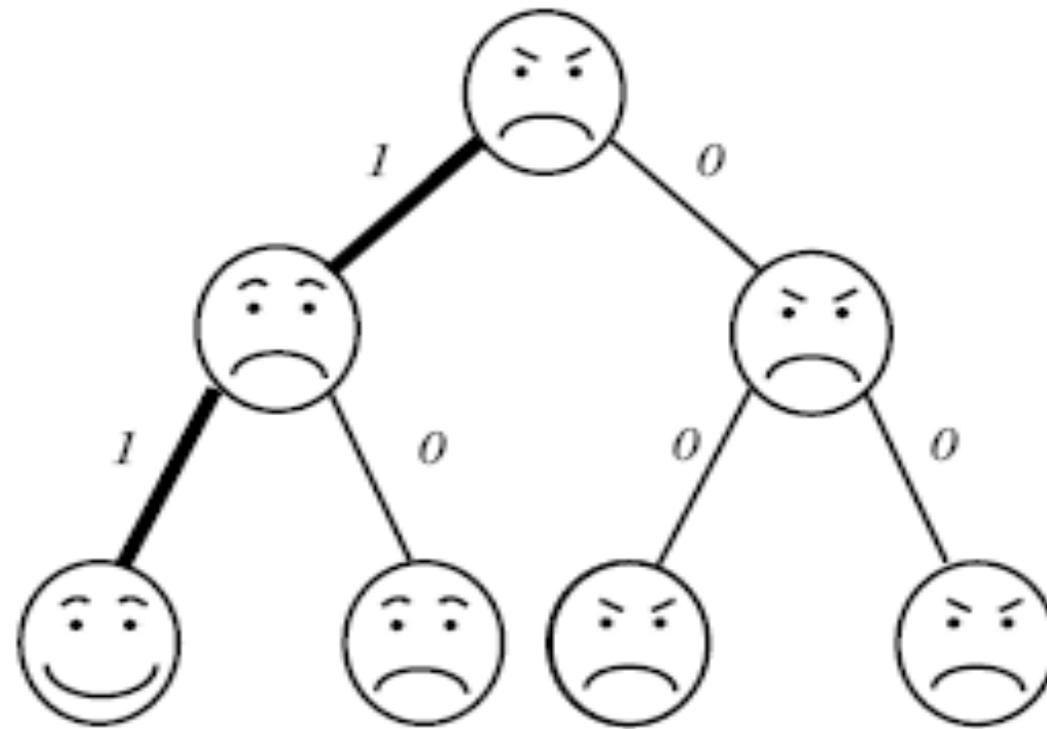
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Character-Based Phylogeny Reconstruction: Criterion



(a) *Parsimony Score=3*



(b) *Parsimony Score=2*

Parsimony: minimize number of changes on edges of tree

A Small and a Large Problem

Small Maximum Parsimony Phylogeny Problem:

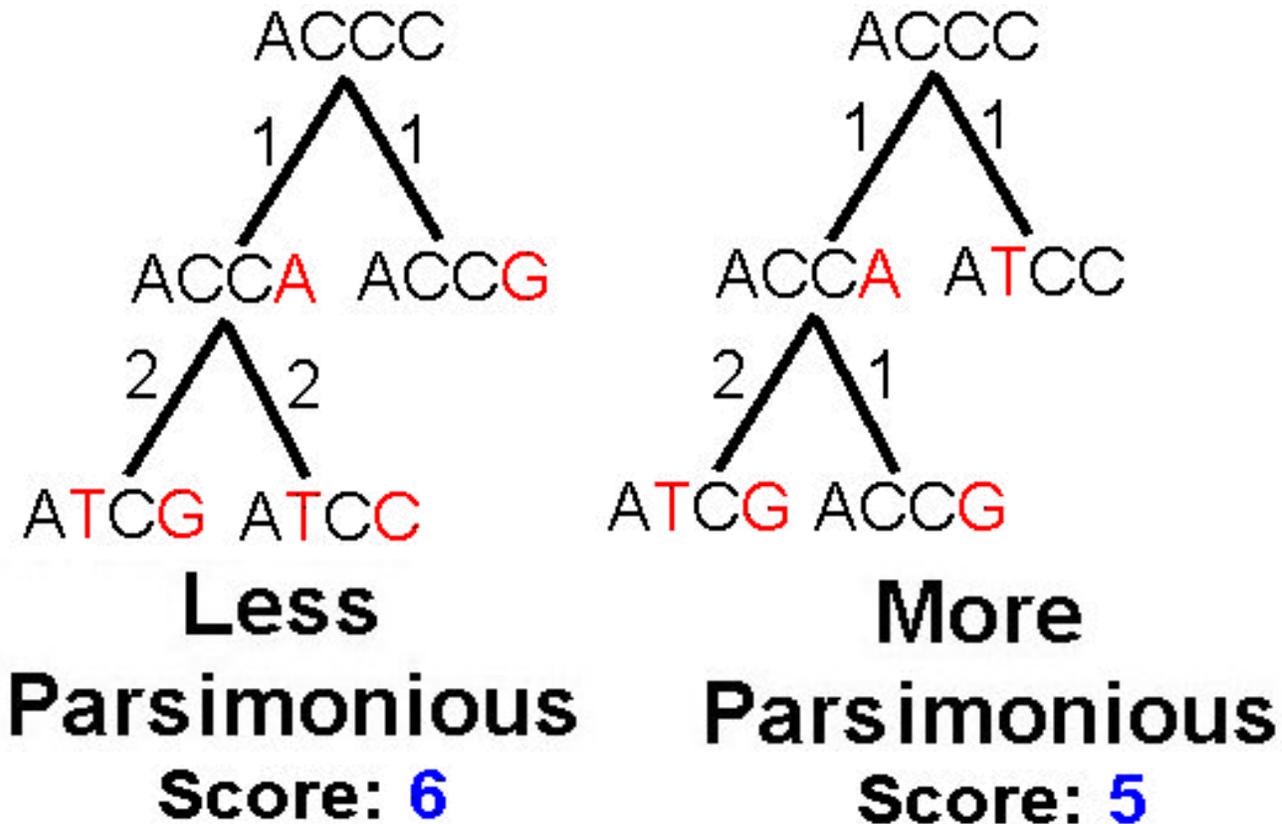
Given $m \times n$ matrix $A = [a_{i,j}]$ and tree T with m leaves, find assignment of character states to each internal vertex of T with minimum parsimony score.

Large Maximum Parsimony Phylogeny Problem:

Given $m \times n$ matrix $A = [a_{i,j}]$, find a tree T with m leaves labeled according to A and an assignment of character states to each internal vertex of T with minimum parsimony score.

Question: Are both problems easy (i.e. in P)?

Small Maximum Parsimony Phylogeny Problem



Key observations: (1) Characters can be solved independently.
(2) Optimal substructure in subtrees.

Recurrence for Small Maximum Parsimony Problem

Small Maximum Parsimony Phylogeny Problem:

Given rooted tree T whose leaves are labeled by $\sigma : L(T) \rightarrow \Sigma$, find assignment of states to each internal vertex of T with minimum parsimony score.

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Let $\mu(v, s)$ be the minimum number of mutations in the subtree rooted at v when assigning state s to v .

Let $\delta(v)$ be the set of children of v .

Recurrence for Small Maximum Parsimony Problem

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Let $\mu(v, s)$ be the minimum number of mutations in the subtree rooted at v when assigning state s to v .

$$c(s, t) = \begin{cases} 0, & \text{if } s = t \\ 1, & \text{if } s \neq t, \end{cases}$$

Let $\delta(v)$ be the set of children of v .

$$\mu(v, s) = \min \begin{cases} \infty, & \text{if } v \in L(T) \text{ and } s \neq \sigma(v), \\ 0, & \text{if } v \in L(T) \text{ and } s = \sigma(v), \\ \sum_{w \in \delta(v)} \min_{t \in \Sigma} \{c(s, t) + \mu(w, t)\}, & \text{if } v \notin L(T). \end{cases}$$

Filling out DP Table and Traceback

Filling out μ $\text{Fill}(T, r(T), \sigma, \Sigma)$

$\text{Fill}(T, v, \sigma, \Sigma)$ $O(m|\Sigma|^2)$

if $v \in \mathcal{L}(T)$ then

For $s \in \Sigma$

if $s = \sigma(v)$ then

$\mu(v, s) = 0$

else $\mu(v, s) = \infty$

else

For $w \in \delta(v)$

$\text{Fill}(T, w, \sigma, \Sigma)$

// children

$\mu(v, s) = 0$

For $w \in \delta(v)$

$\mu(v, s) += \min_{t \in \Sigma} \{c(s, t) + \mu(w, t)\}$

Traceback (T, v, μ)

if $v = r(T)$

$\sigma(r(T)) = \arg \min_{s \in \Sigma} \{ \mu(r(T), s) \}$

else let u be the parent of v and let $s = \sigma(u)$

$\sigma(v) = \arg \min_{t \in \Sigma} \{ c(s, t) + \mu(v, t) \}$

For $w \in \delta(v)$

Traceback (T, w, μ)

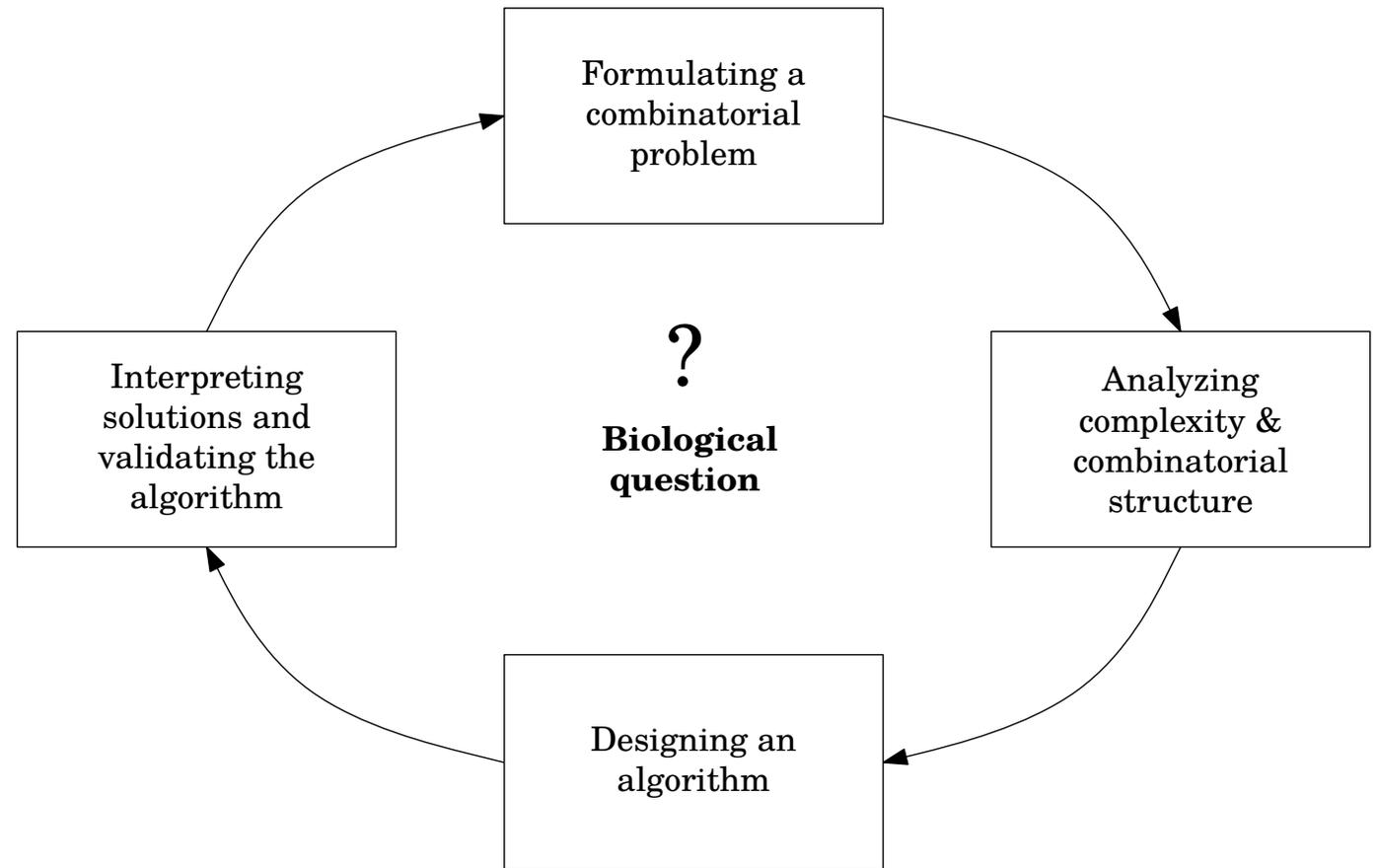
Let $r(T)$ be the root vertex

Outline

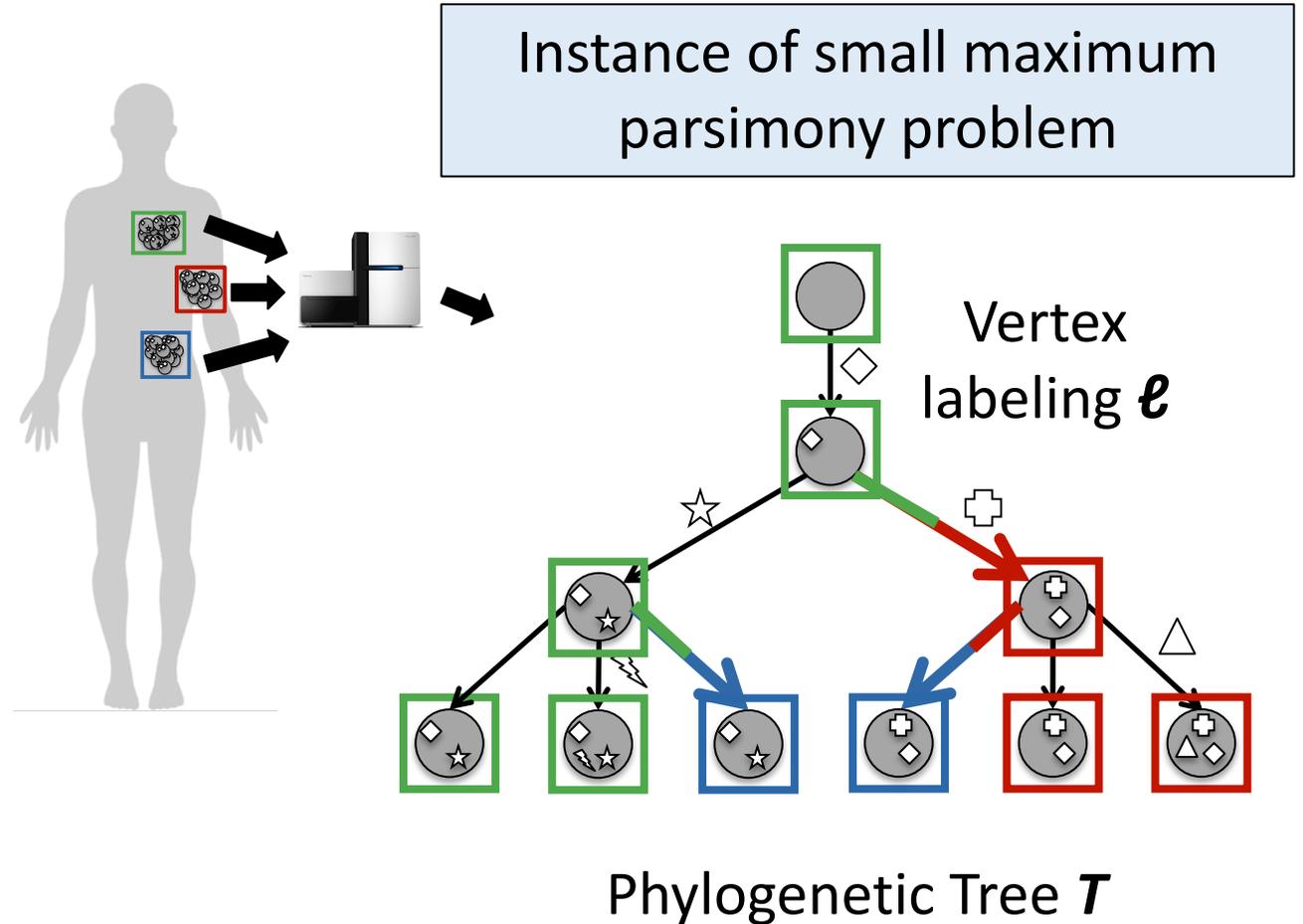
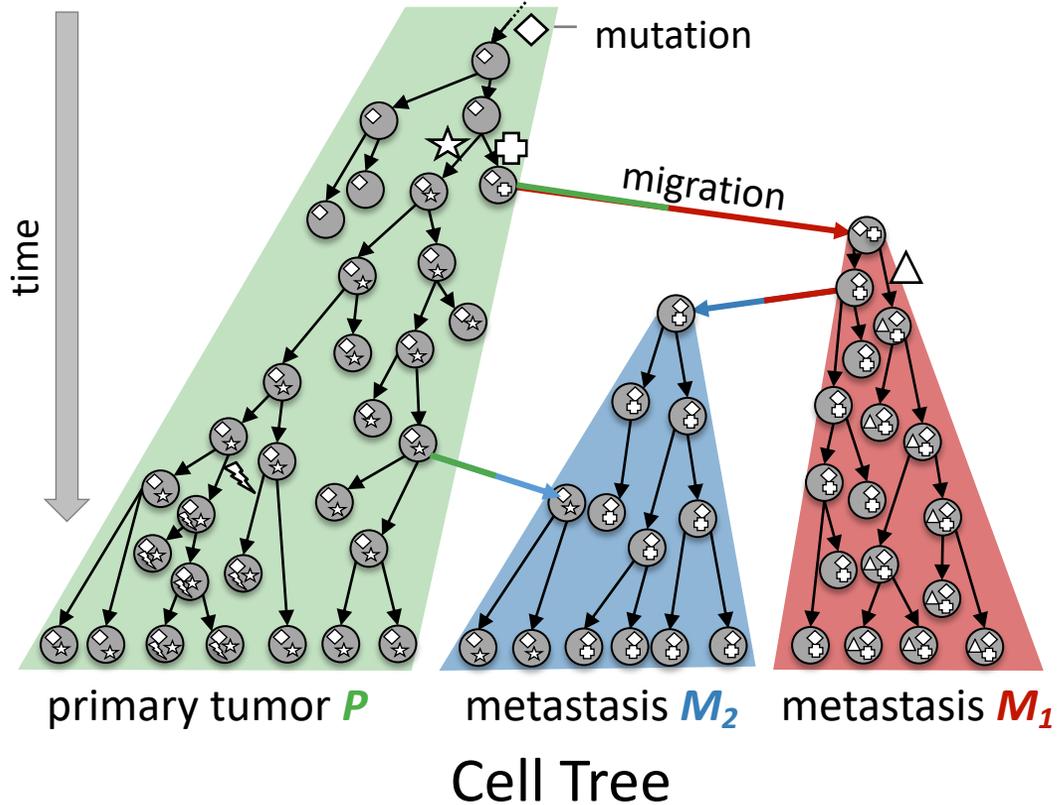
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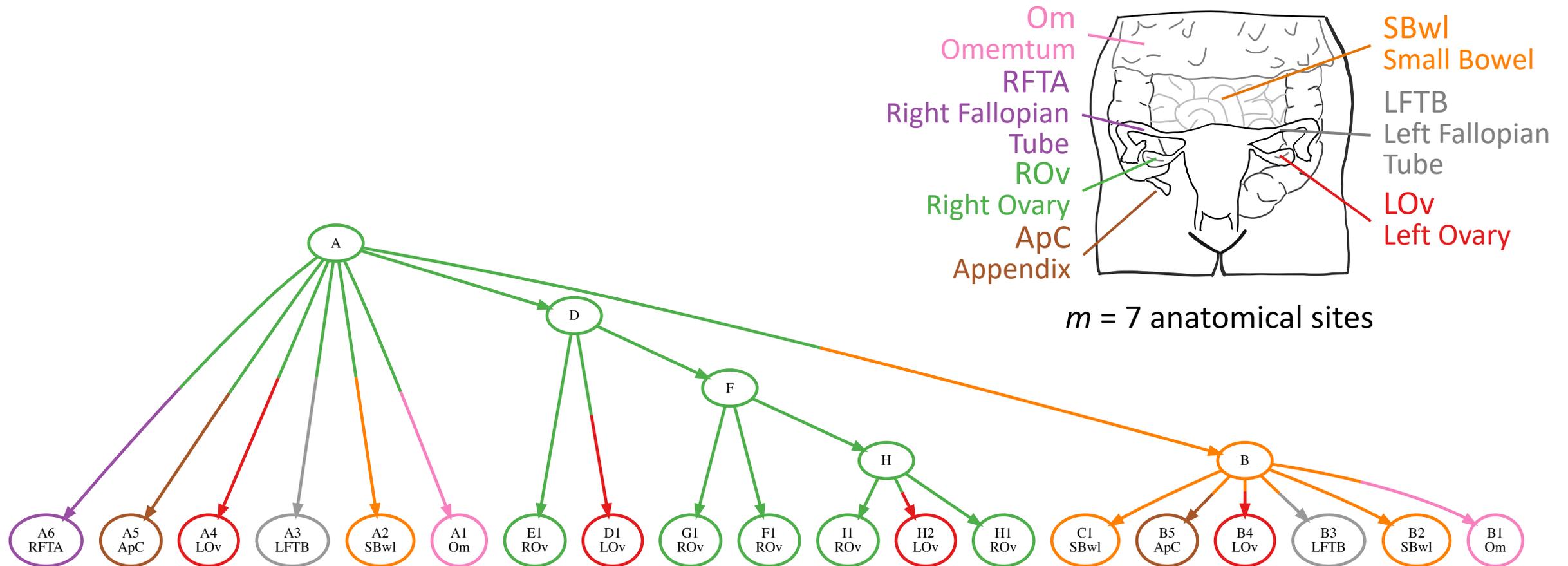


Goal: Given phylogenetic tree T , find *parsimonious* vertex labeling ℓ with fewest migrations

Minimum Migration Analysis in Ovarian Cancer

McPherson et al. (2016). Divergent modes of clonal spread and intraperitoneal mixing in high-grade serous ovarian cancer. *Nature Genetics*.

- Instance of the maximum parsimony small phylogeny problem [Fitch, 1971; Sankoff, 1975]

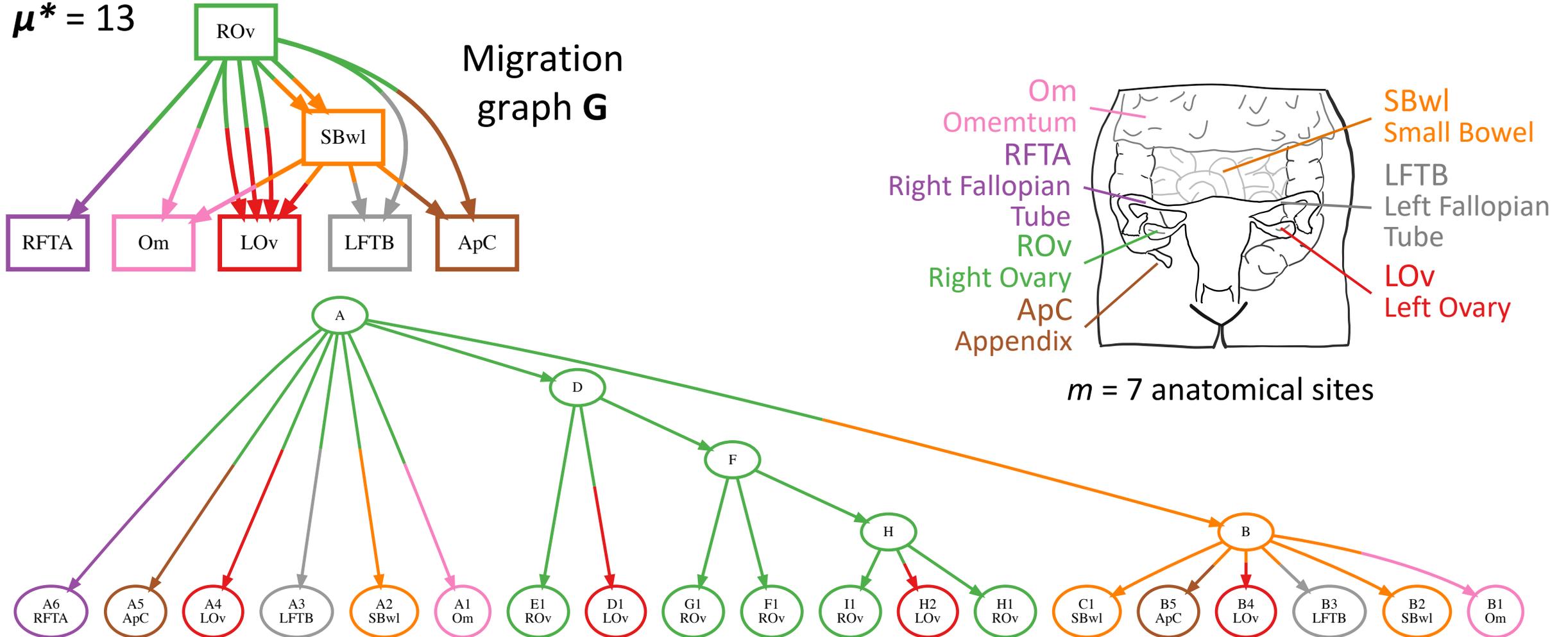


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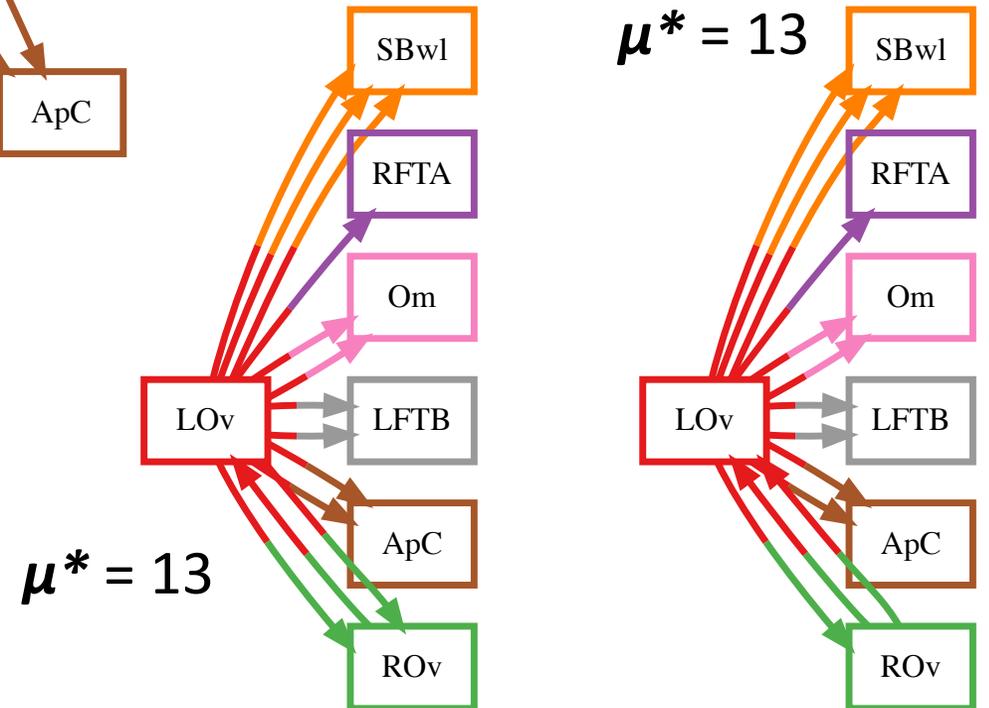
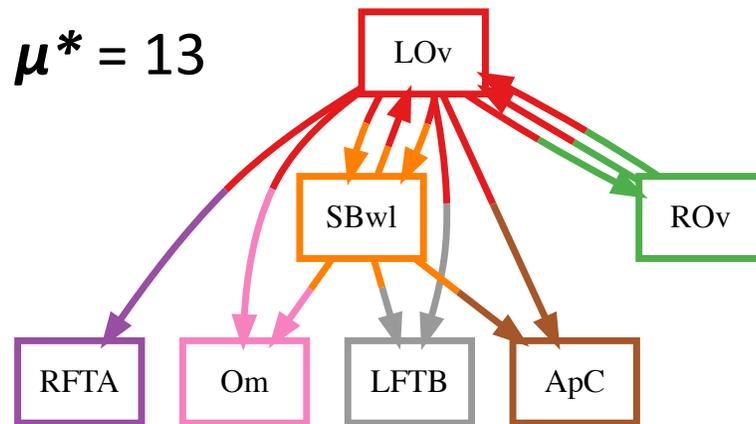
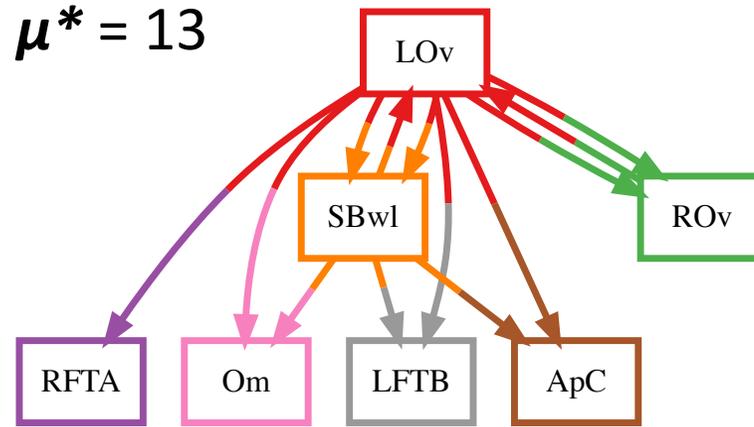
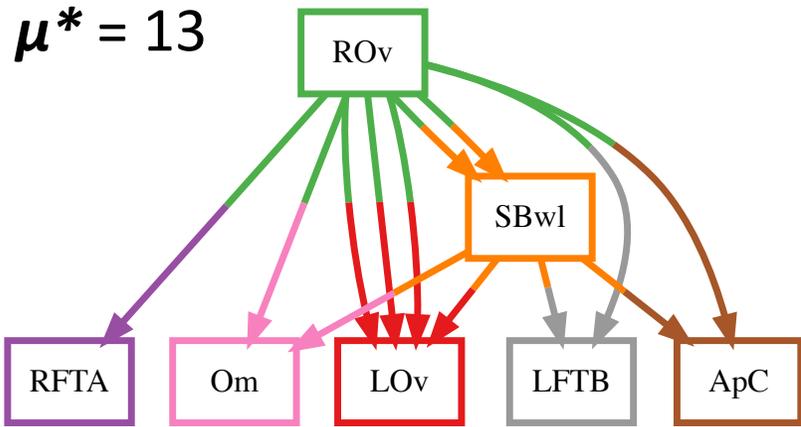
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$\mu^* = 13$



Minimum Migration History is *Not* Unique

- Enumerate all minimum-migration vertex labelings in the backtrace step

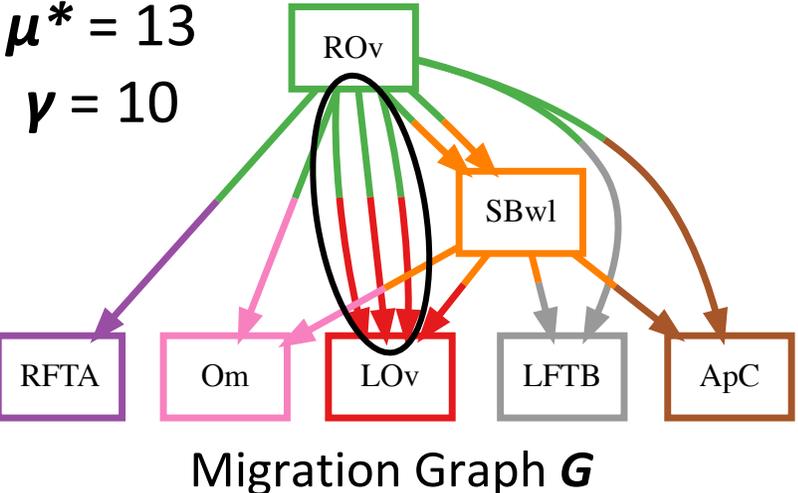


ApC	Appendix
LFTB	Left Fallopian Tube
LOv	Left Ovary
RFTA	Right Fallopian Tube
ROv	Right Ovary
SBwl	Small Bowel
Om	Omentum

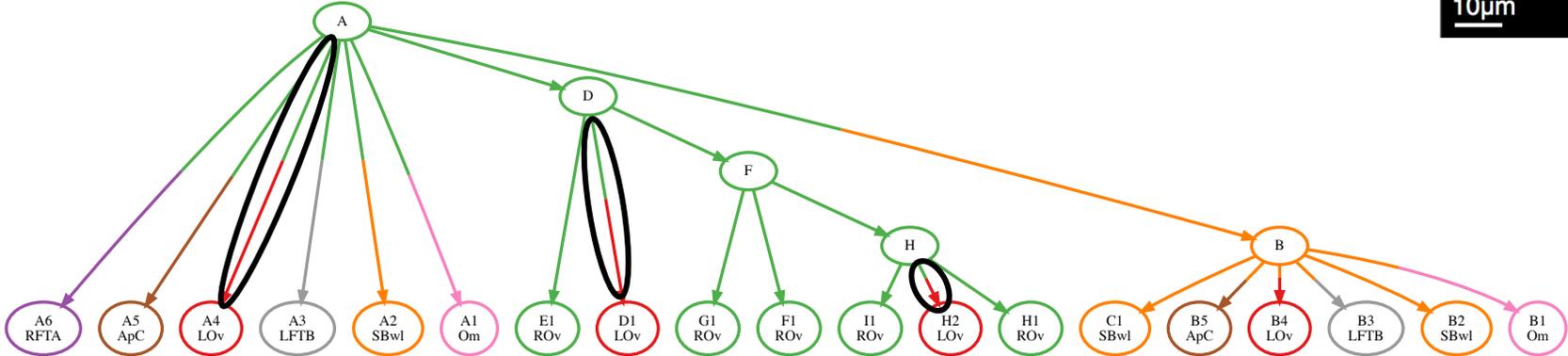
Comigrations: Simultaneous Migrations of Multiple Clones

- Multiple tumor cells migrate simultaneously through the blood stream [Cheung et al., 2016]
- Second objective: number ν of **comigrations** is the number of multi-edges in migration graph G^\dagger

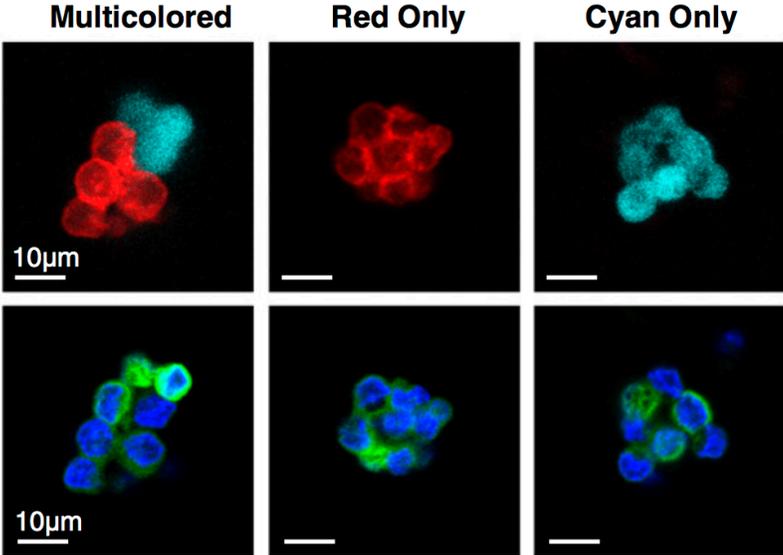
† Not necessarily true in the case of directed cycles



Clone Tree T



F Circulating tumor cell clusters

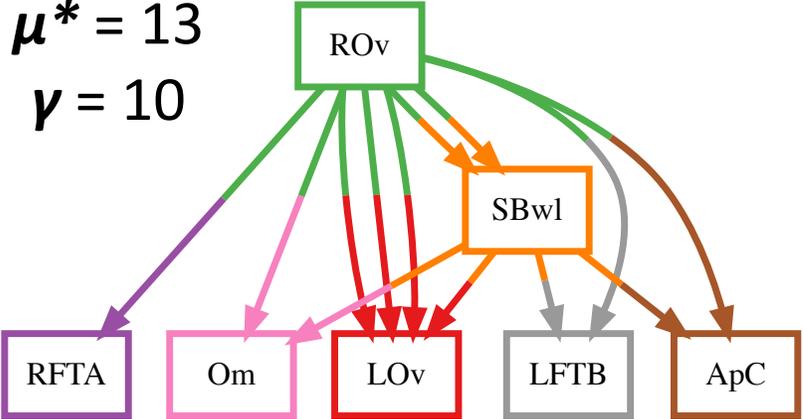


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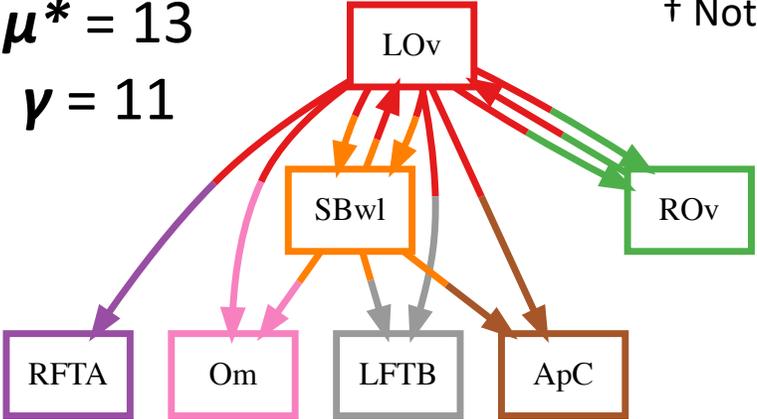
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$\mu^* = 13$
 $\gamma = 10$

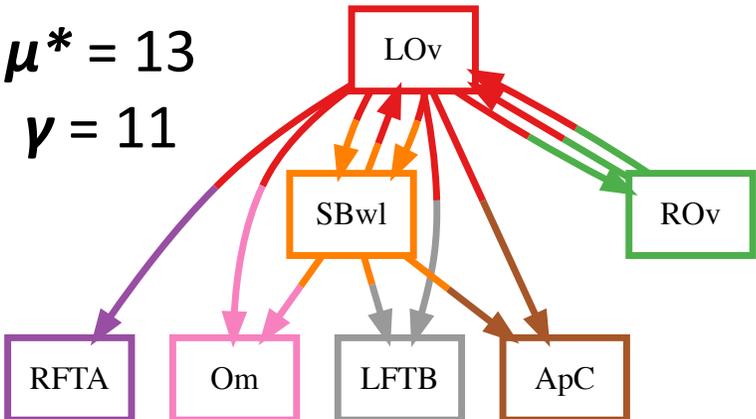


$\mu^* = 13$
 $\gamma = 11$

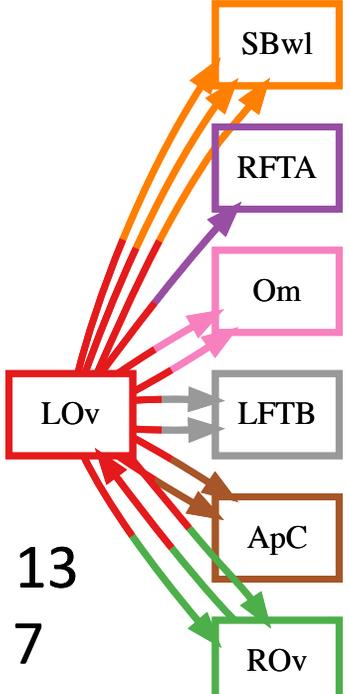


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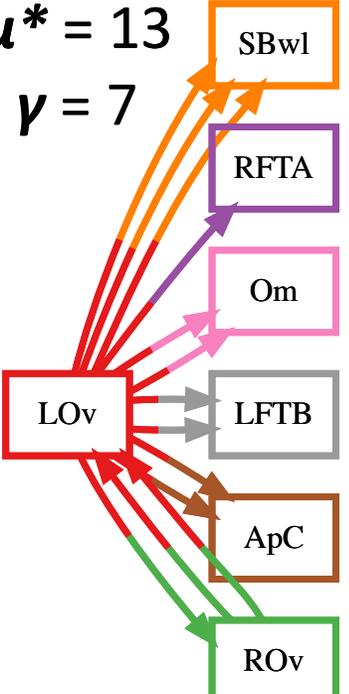
$\mu^* = 13$
 $\gamma = 11$



$\mu^* = 13$
 $\gamma = 7$

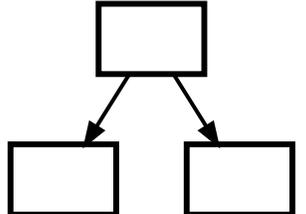
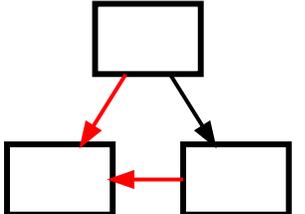
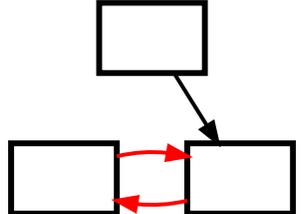
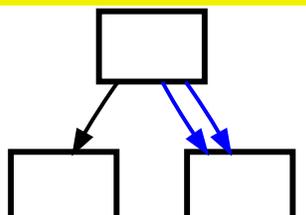
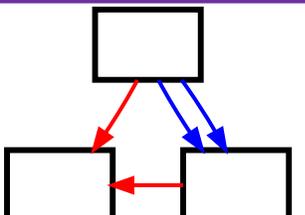
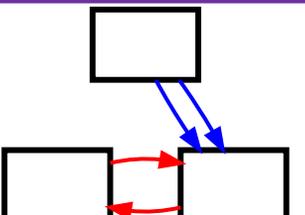


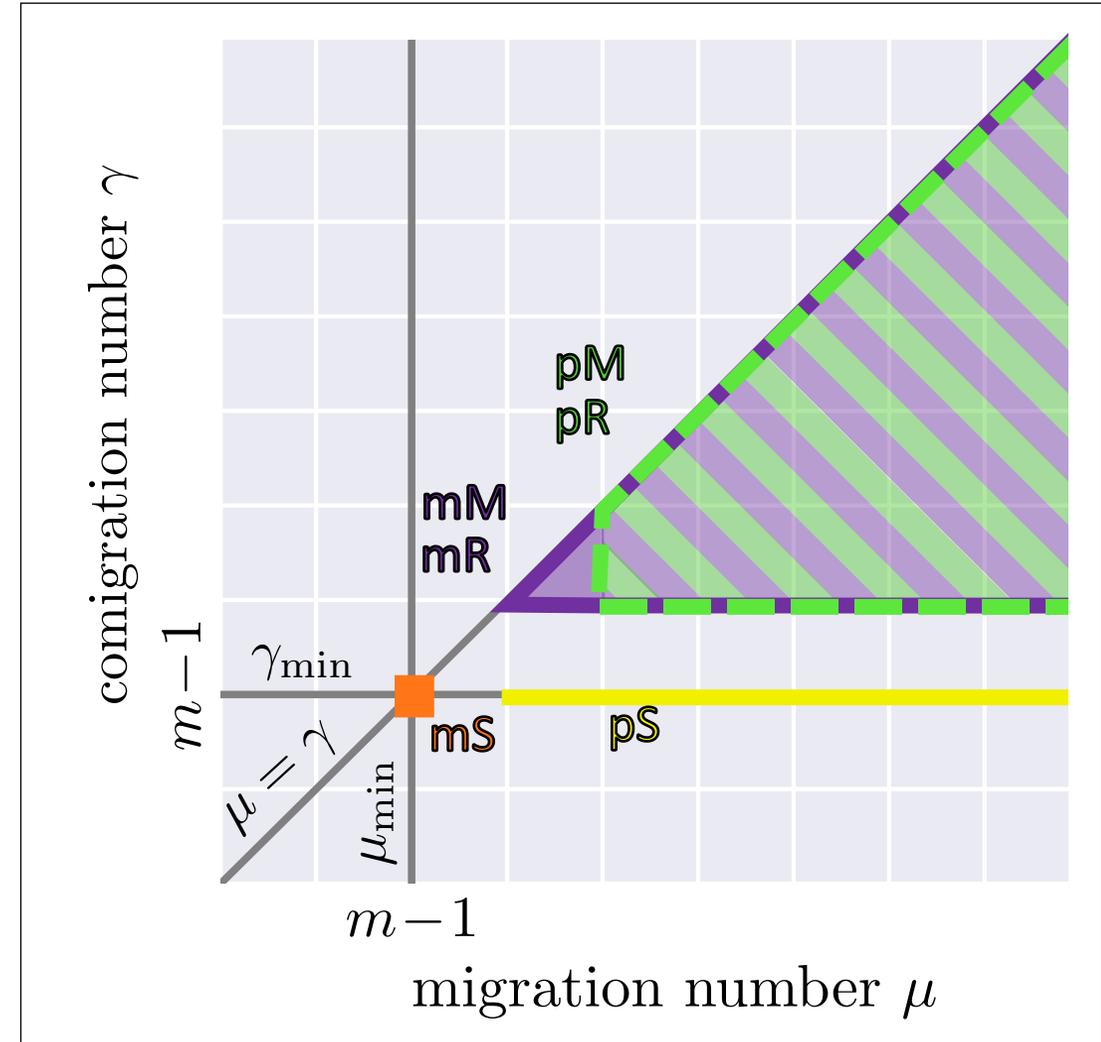
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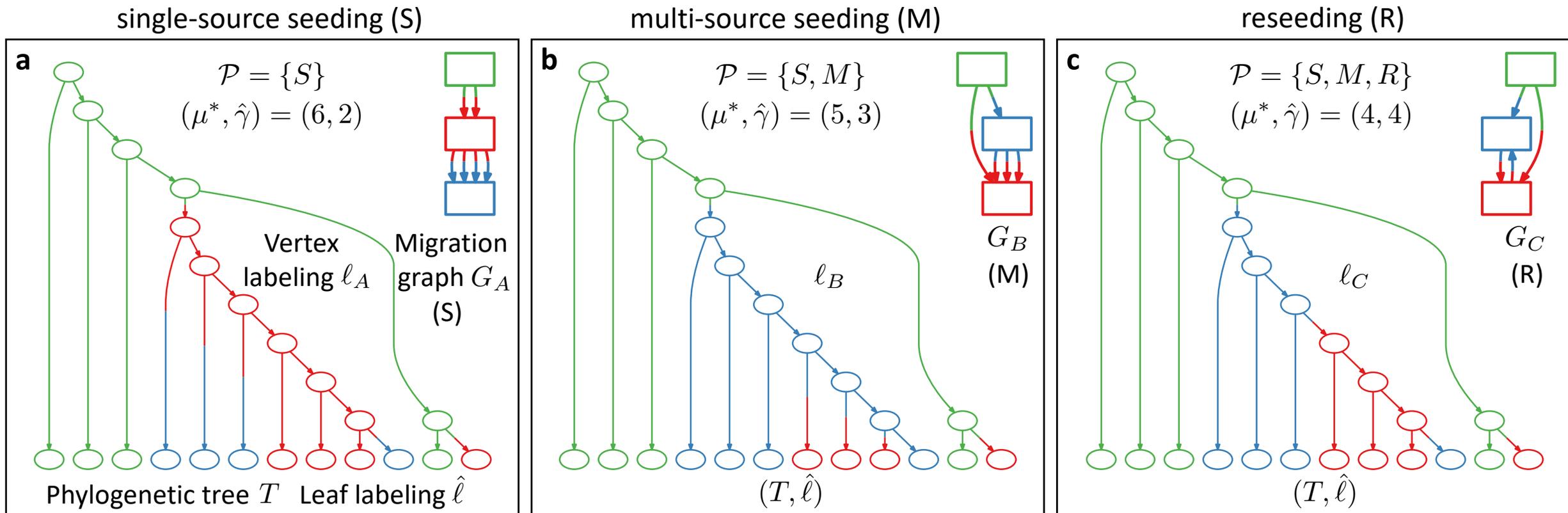
Tradeoffs between Migrations, Comigrations and Migration Pattern

	single-source seeding (S)	multi-source seeding (M)	reseeding (R)
monoclonal (m)	 <p>tree</p>	 <p>directed acyclic graph</p>	 <p>directed graph</p>
polyclonal (p)	 <p>multi-tree</p>	 <p>directed acyclic multi-graph</p>	 <p>directed multi-graph</p>



Constrained Multi-objective Optimization Problem

Parsimonious Migration History (PMH): Given a phylogenetic tree T and a set $\mathcal{P} \subseteq \{S, M, R\}$ of allowed migration patterns, find vertex labeling ℓ with minimum migration number $\mu^*(T)$ and smallest comigration number $\hat{\gamma}(T)$.

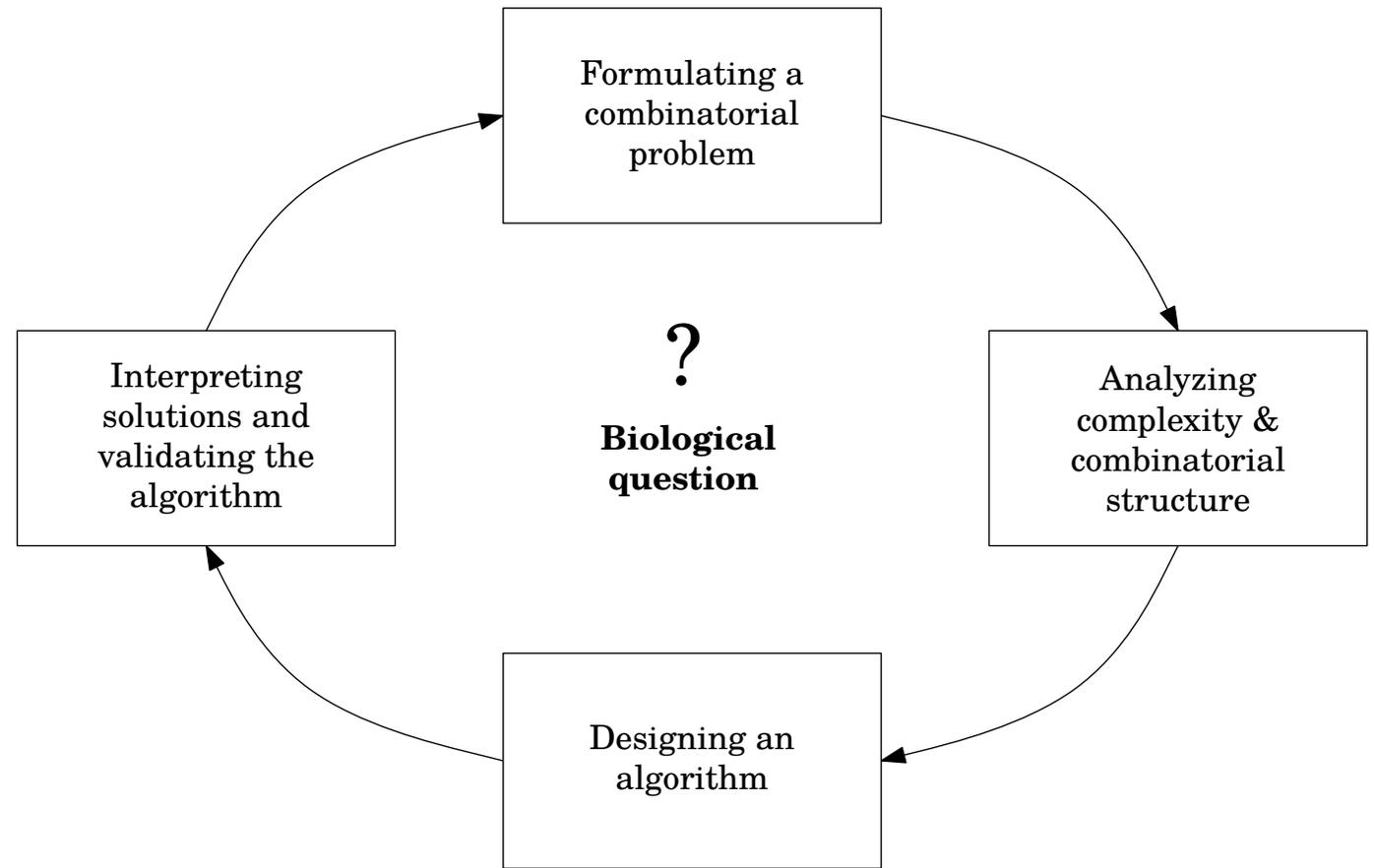


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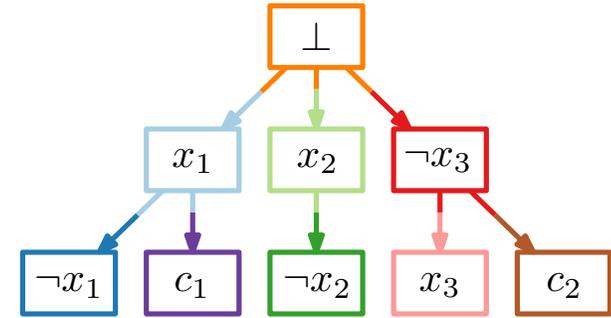
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PMH is NP-hard when $\mathcal{P} = \{S\}$

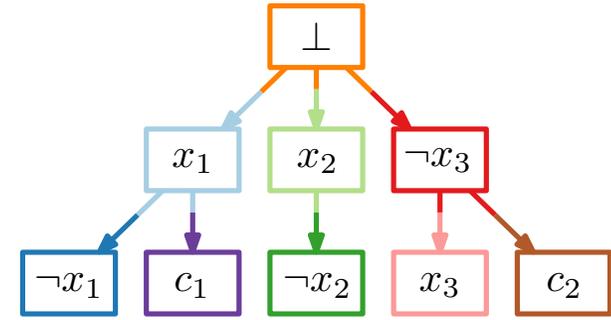
3-SAT: Given $\varphi = \bigwedge_{i=1}^k (y_{i,1} \vee y_{i,2} \vee y_{i,3})$ with variables $\{x_1, \dots, x_n\}$ and k clauses, find $\phi : [n] \rightarrow \{0,1\}$ satisfying φ



$$\Sigma = \{x_1, \dots, x_n, \neg x_1, \dots, \neg x_n, c_1, \dots, c_k, \perp\}$$

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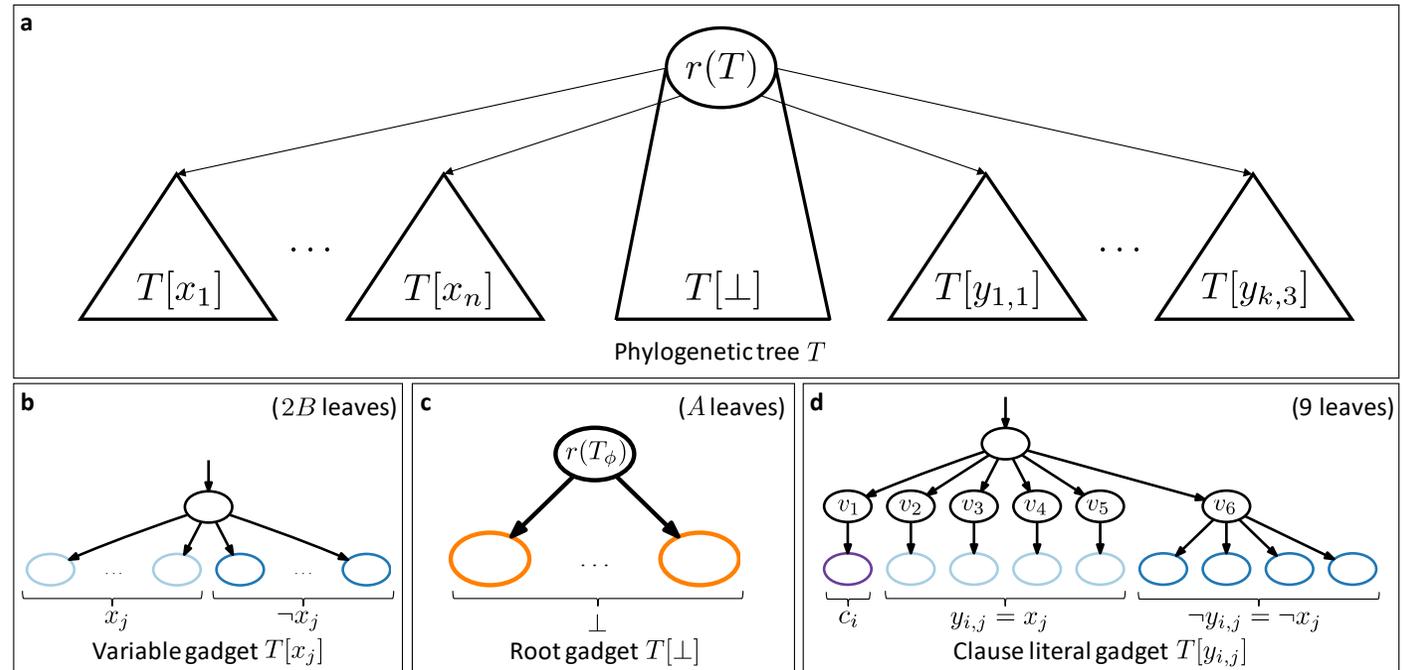
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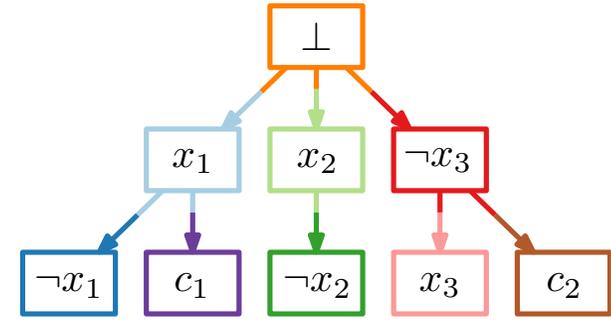
Three ideas:

1. Ensure that $(x, \neg x) \in E(G)$ or $(\neg x, x) \in E(G)$
2. Ensure that $\ell^*(r(T)) = \perp$
3. Ensure that φ is satisfiable if and only if ℓ^* encodes a satisfying truth assignment



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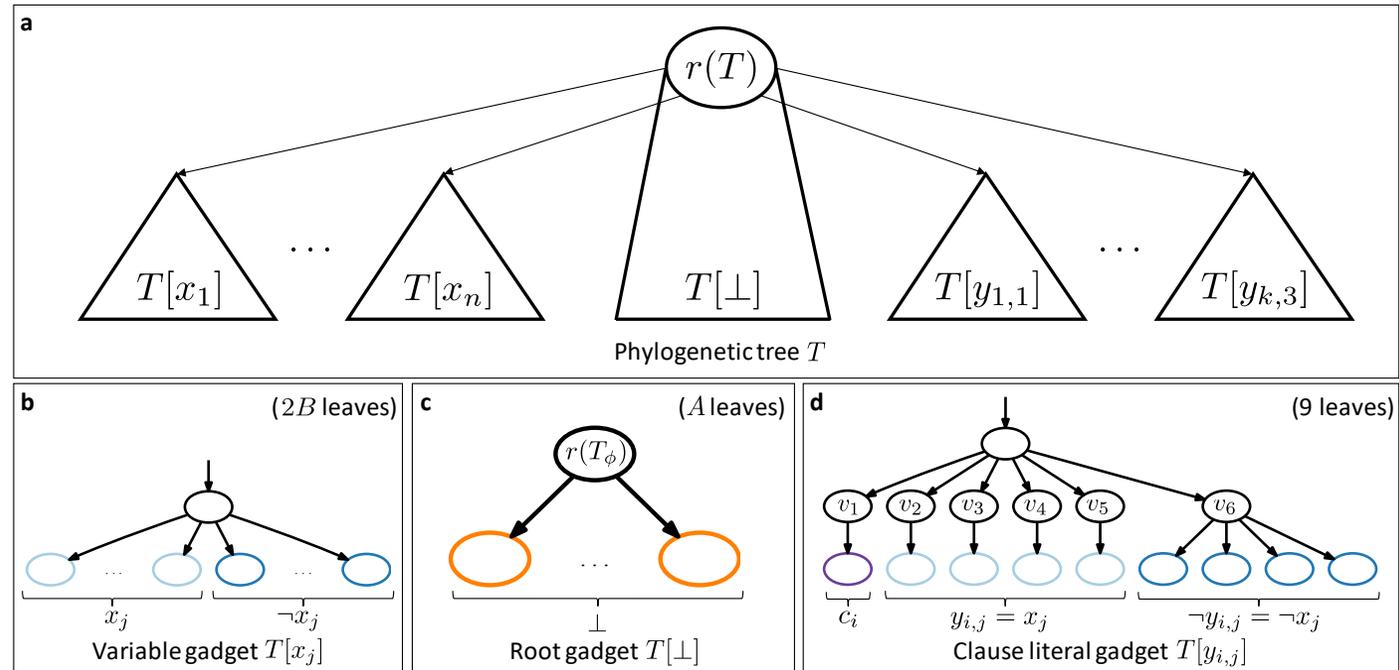
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2. Ensure that $\ell^*(r(T)) = \perp$
3. Ensure that φ is satisfiable if and only if ℓ^* encodes a satisfying truth assignment



Lemma: Let $B > 10k + 1$ and $A > 2Bn + 27k$.

Then, φ is satisfiable if and only if $\mu^*(T) = (B + 1)n + 25k$

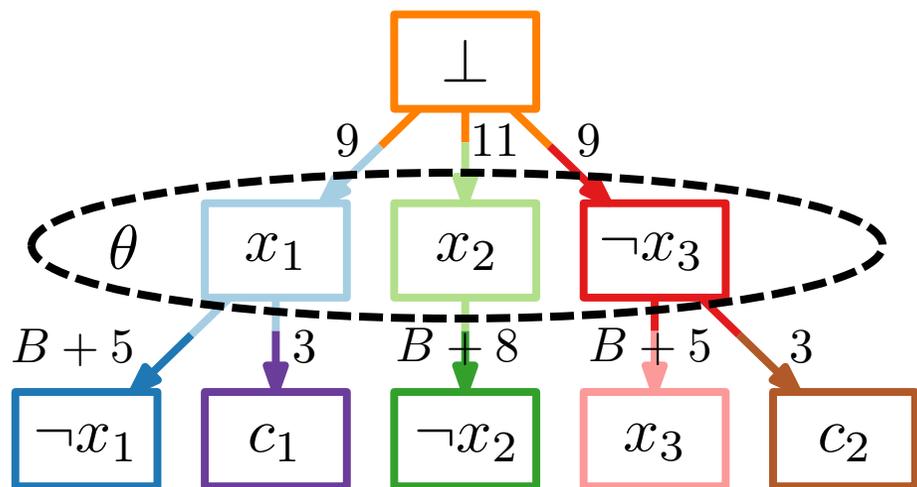
PMH is NP-hard when $\mathcal{P} = \{S\}$

$$\varphi = (x_1 \vee x_2 \vee \neg x_3) \wedge (\neg x_1, \neg x_2, \neg x_3)$$

$$k = 2, n = 3$$

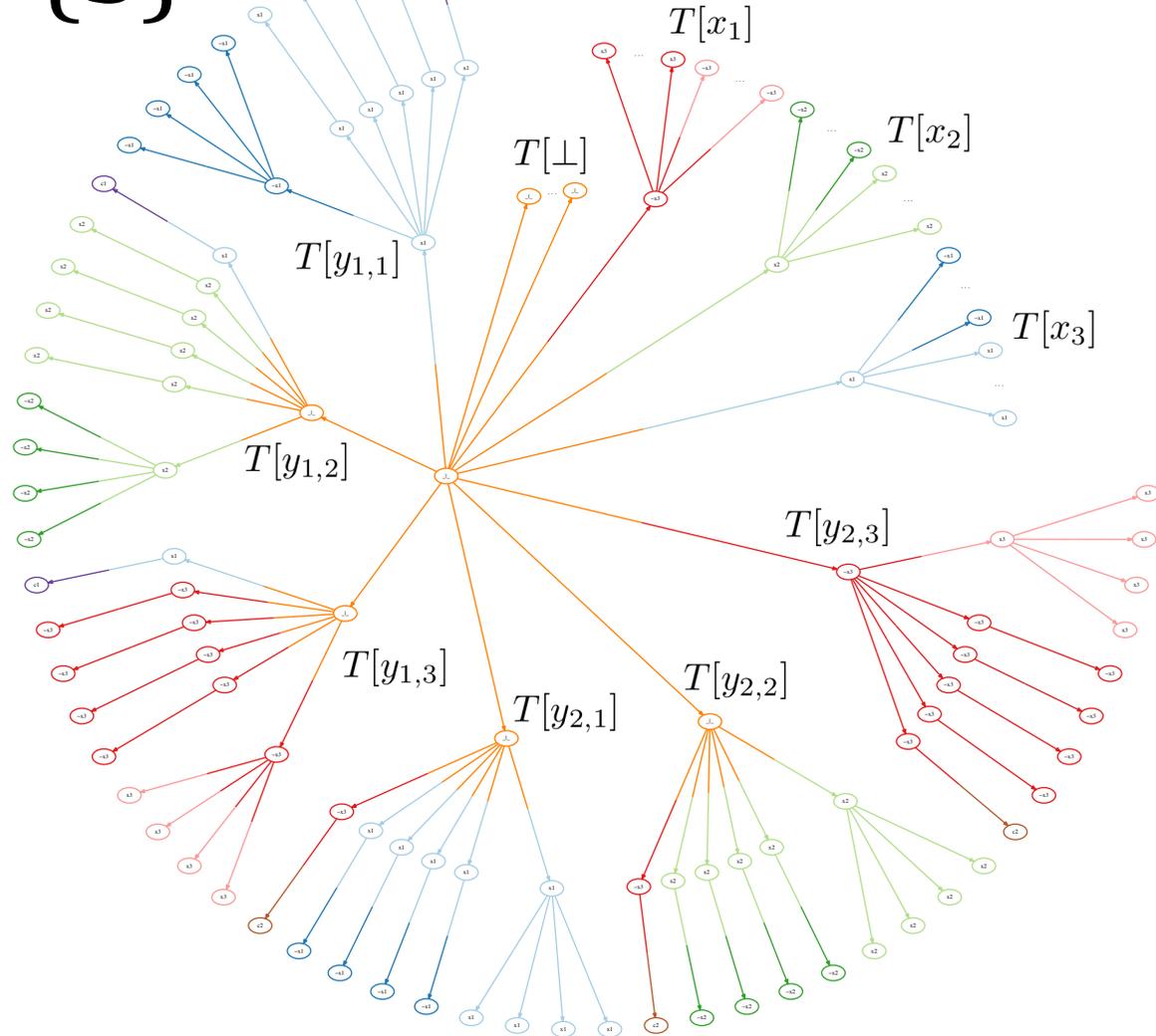
$$B = 10k + 2 = 22$$

$$A = 2Bn + 27k + 1 = 187$$



$$\Sigma = \{x_1, x_2, x_3, \neg x_1, \neg x_2, \neg x_3, c_1, c_2, \perp\}$$

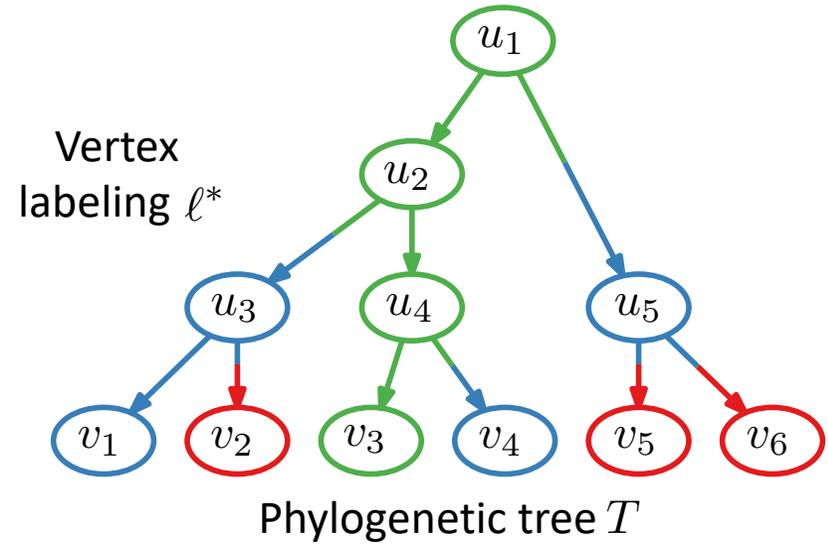
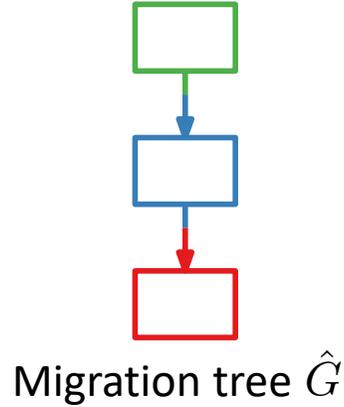
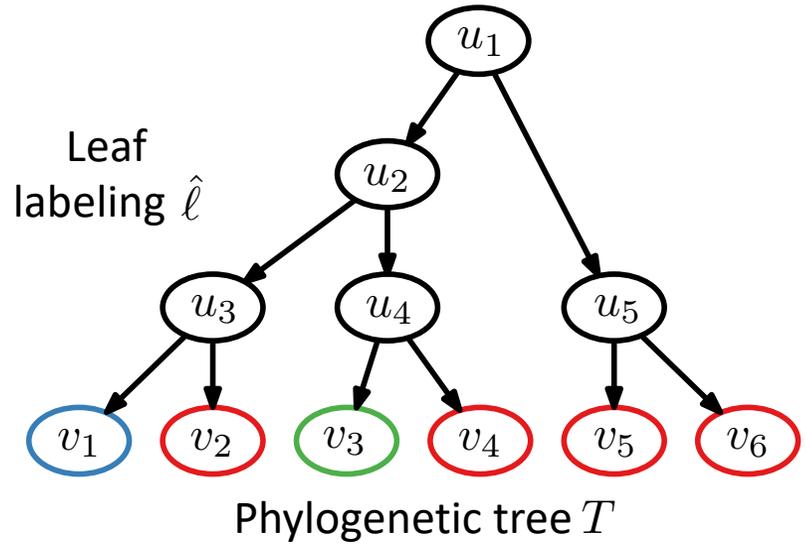
$$\begin{aligned} \mu^*(T) &= (B + 1)n + 25k \\ &= 23 * 3 + 50 * 2 = 119 \end{aligned}$$



Lemma: Let $B > 10k + 1$ and $A > 2Bn + 27k$.

Then, φ is satisfiable if and only if $\mu^*(T) = (B + 1)n + 25k$

PMH is FPT in number m of locations when $\mathcal{P} = \{S\}$



Lemma: If there exists labeling ℓ consistent with \hat{G} then

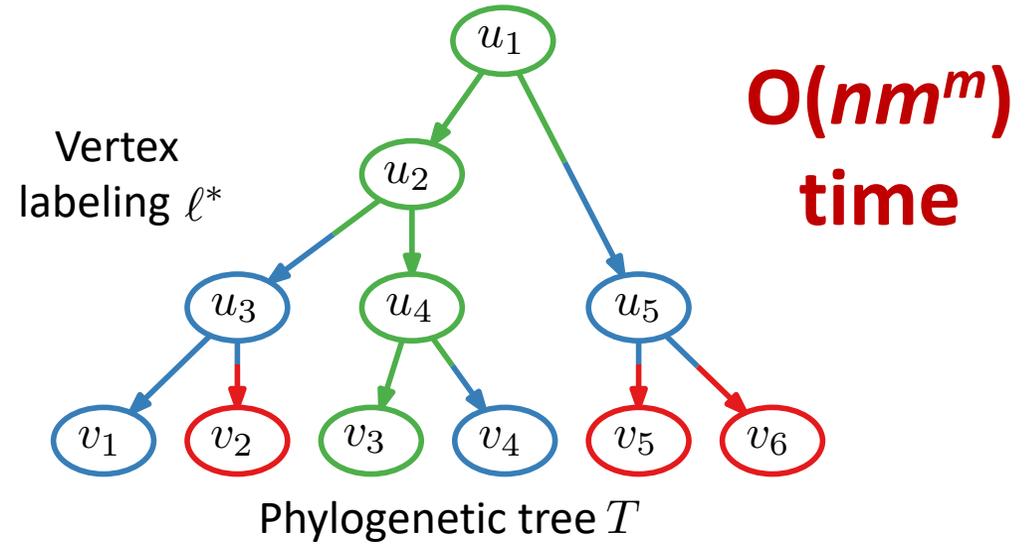
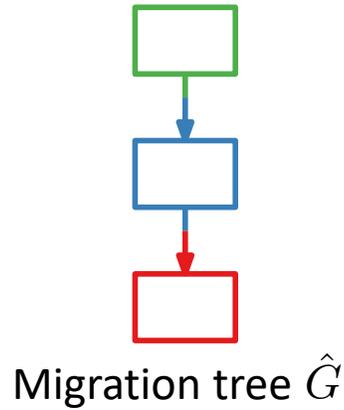
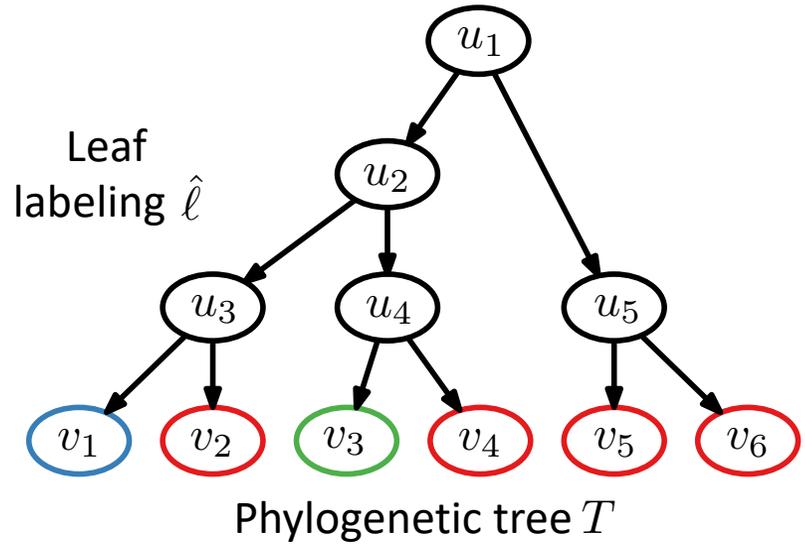
$$d_T(u, v) \geq d_{\hat{G}}(\text{lca}_{\hat{G}}(u), \hat{\ell}(v)) \quad \forall u, v \in V(T) \text{ such that } u \preceq_T v. \quad (1)$$

$$\ell^*(v) = \begin{cases} \text{LCA}_{\hat{G}}(r(T)), & \text{if } v = r(T), \\ \sigma(\ell^*(\pi(v)), \text{LCA}_{\hat{G}}(v)), & \text{if } v \neq r(T), \end{cases}$$

where $\sigma(s, t) = s$ if $s = t$ and otherwise $\sigma(s, t)$ is the unique child of s that lies on the path from s to t in \hat{G} .

Lemma: If (1) holds then ℓ^* is a minimum migration labeling consistent with \hat{G} .

PMH is FPT in number m of locations when $\mathcal{P} = \{S\}$



Lemma: If there exists labeling ℓ consistent with \hat{G} then

$$d_T(u, v) \geq d_{\hat{G}}(\text{lca}_{\hat{G}}(u), \hat{\ell}(v)) \quad \forall u, v \in V(T) \text{ such that } u \preceq_T v. \quad (1)$$

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where $\sigma(s, t) = s$ if $s = t$ and otherwise $\sigma(s, t)$ is the unique child of s that lies on the path from s to t in \hat{G} .

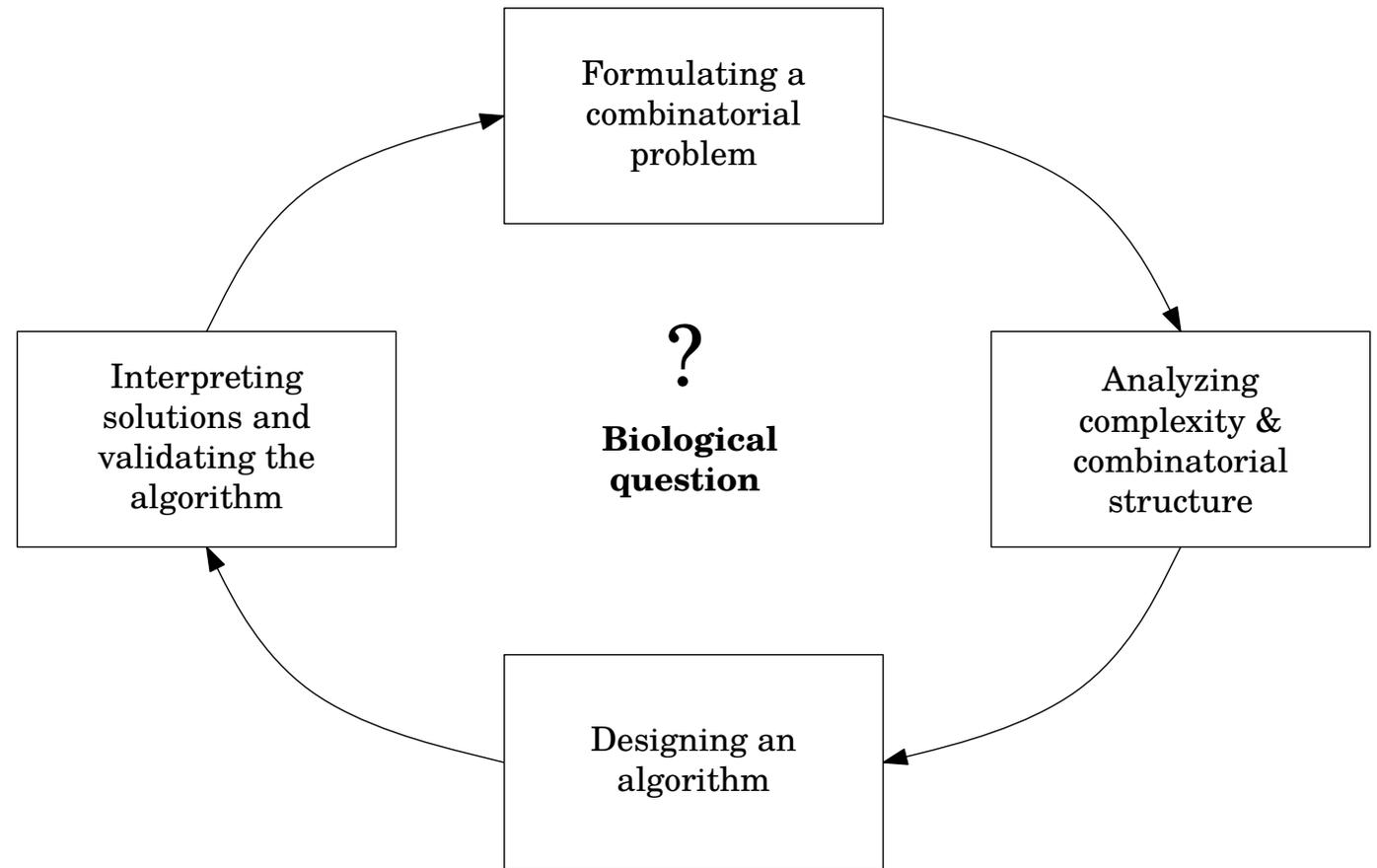
Lemma: If (1) holds then ℓ^* is a minimum migration labeling consistent with \hat{G} .

Outline

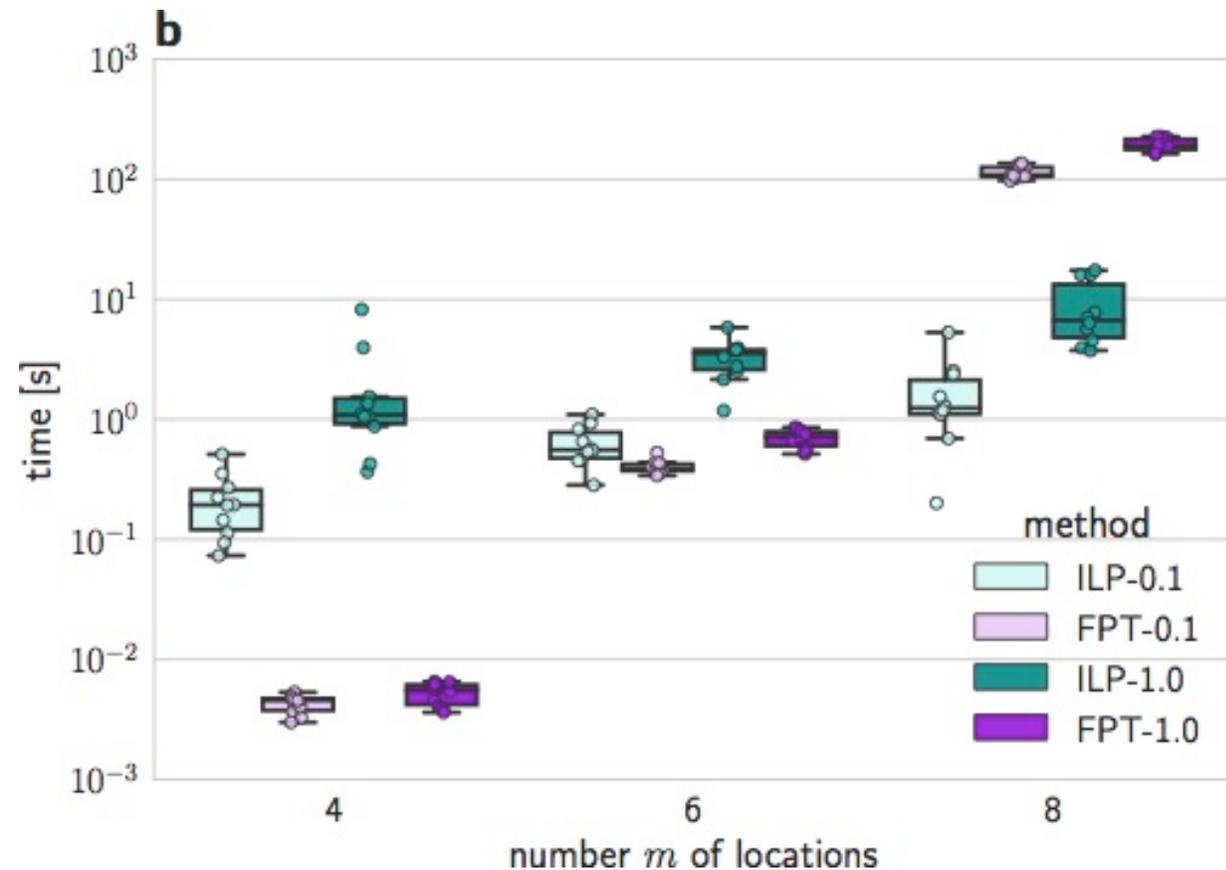
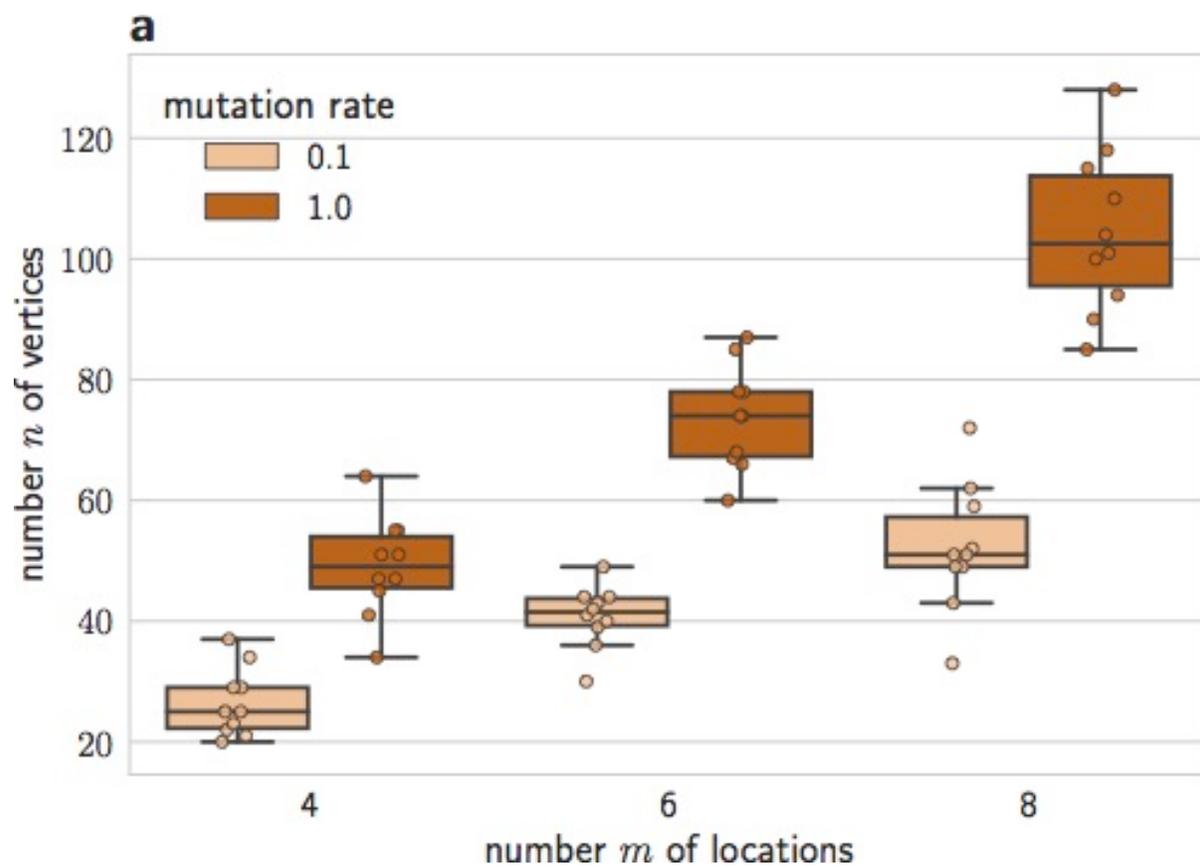
- Metastasis
- Maximum parsimony
- Problem statement
- Complexity
- Algorithm & results
- Problem variants

Reading:

- **M. El-Kebir**, G. Satas and B.J. Raphael. Inferring parsimonious migration histories for metastatic cancers. [Nature Genetics, 50:718-726, 2018.](#)
- **M. El-Kebir**[†]. [Parsimonious Migration History Problem: Complexity and Algorithms.](#) WABI 2018, Helsinki, Finland, August 20-22, 2018.



Simulations



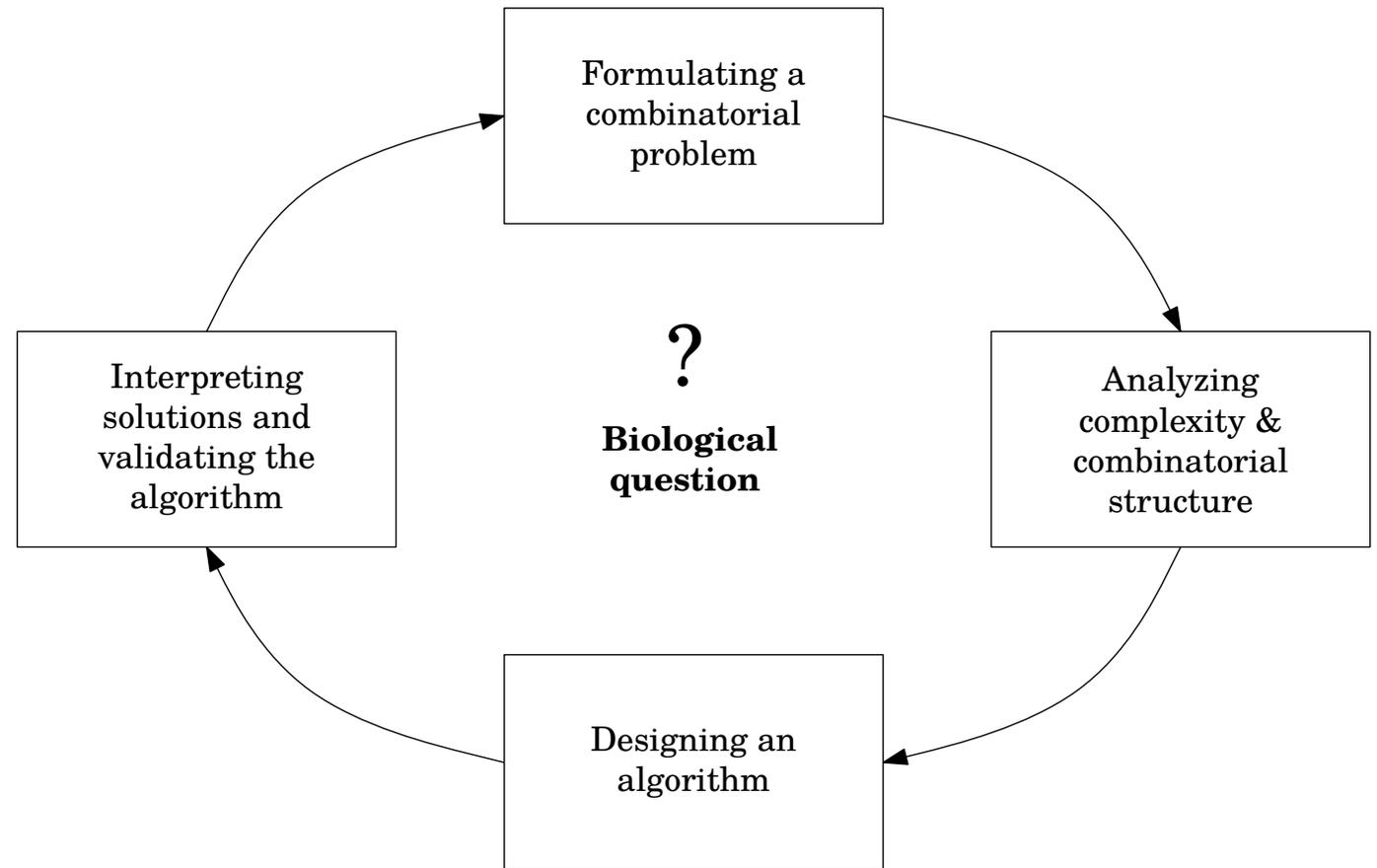
Available on: <https://github.com/elkebir-group/PMH-S>

Outline

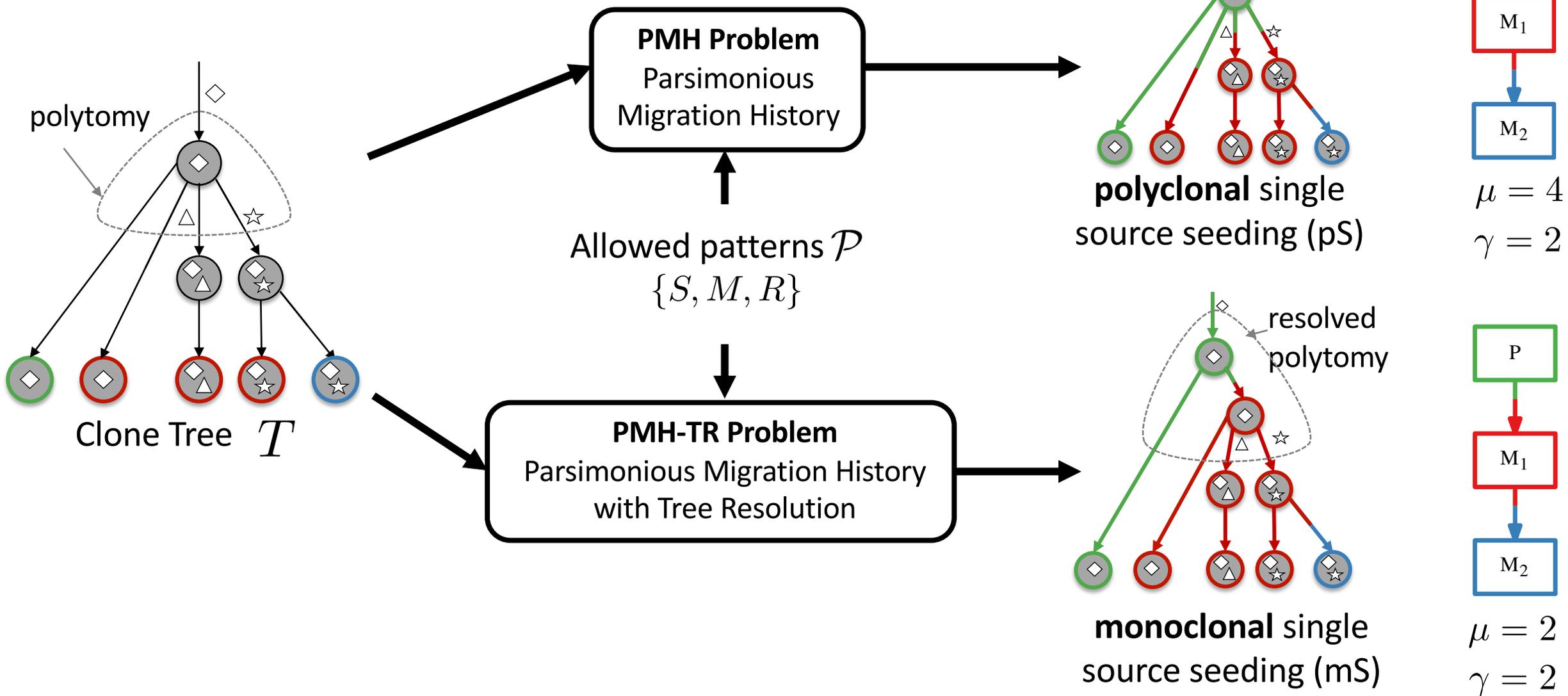
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Reading:

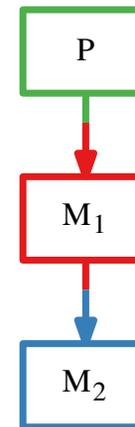
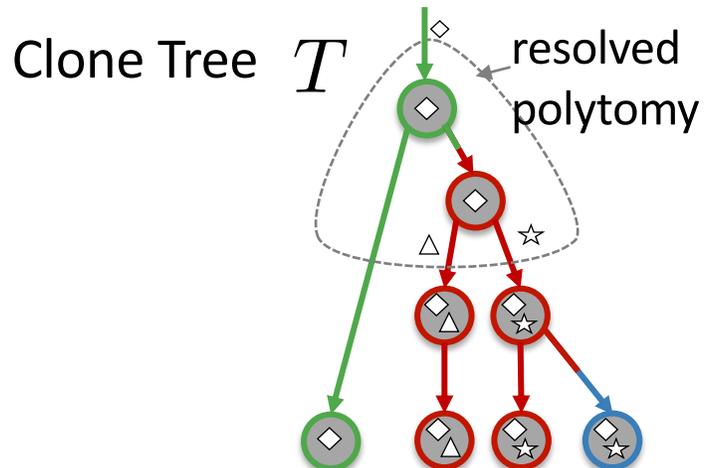
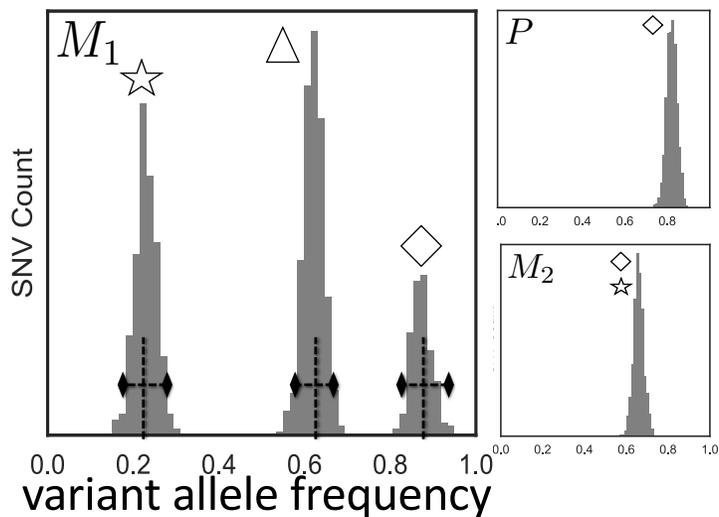
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Resolving Clone Tree Ambiguities



Resolving Clone Tree Ambiguities



$$\mu = 2$$

$$\gamma = 2$$

monoclonal single source seeding (mS)

confidence intervals

$$F^- = \begin{matrix} & \diamond & \triangle & \star \\ P & \begin{bmatrix} 0.75 & 0 & 0 \end{bmatrix} \\ M_1 & \begin{bmatrix} 0.82 & 0.15 & 0.56 \end{bmatrix} \\ M_2 & \begin{bmatrix} 0.66 & 0 & 0.68 \end{bmatrix} \end{matrix}$$

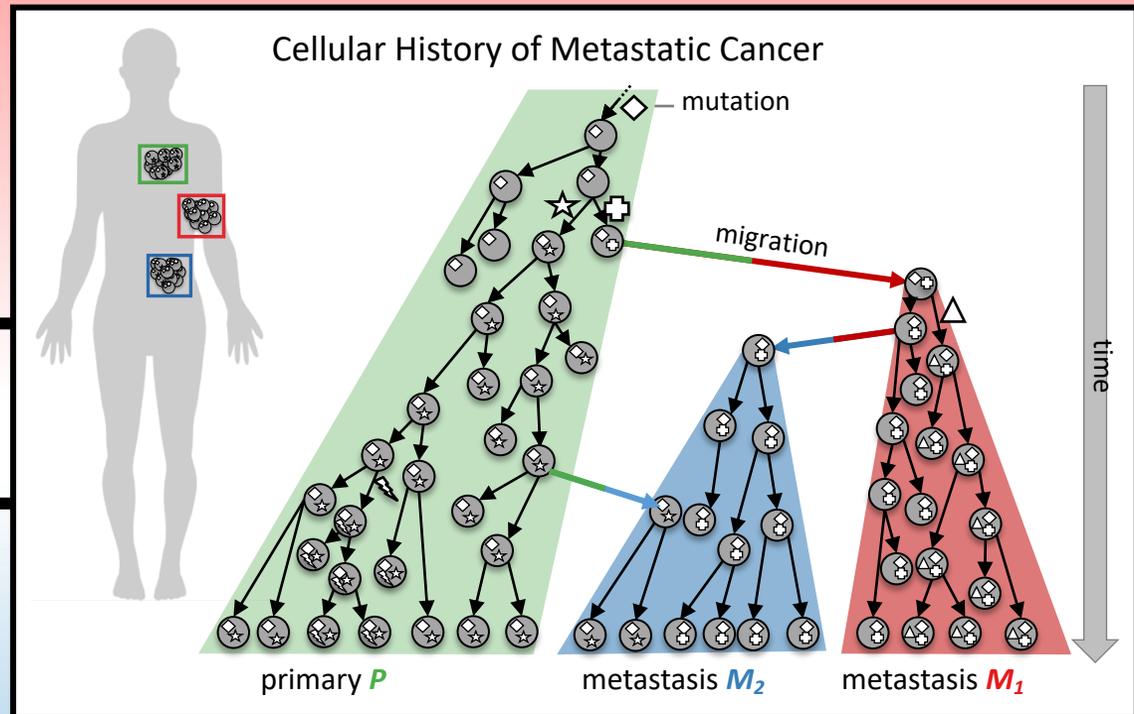
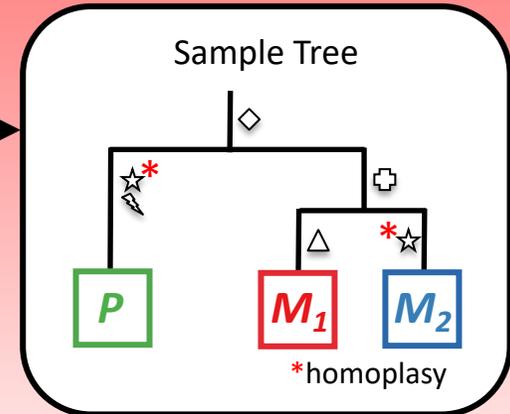
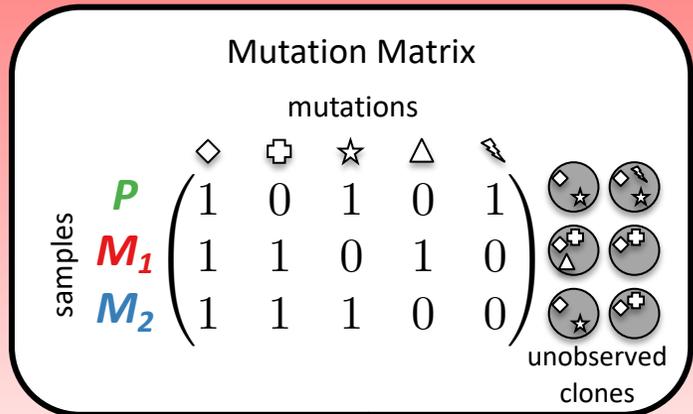
$$F^+ = \begin{matrix} & \diamond & \triangle & \star \\ P & \begin{bmatrix} 0.86 & 0 & 0 \end{bmatrix} \\ M_1 & \begin{bmatrix} 0.92 & 0.30 & 0.68 \end{bmatrix} \\ M_2 & \begin{bmatrix} 0.78 & 0 & 0.77 \end{bmatrix} \end{matrix}$$

$$\underbrace{\begin{matrix} & \diamond & \triangle & \star \\ P & \begin{bmatrix} 0.82 & 0 & 0 \end{bmatrix} \\ M_1 & \begin{bmatrix} 0.88 & 0.24 & 0.64 \end{bmatrix} \\ M_2 & \begin{bmatrix} 0.73 & 0 & 0.73 \end{bmatrix} \end{matrix}}_{\hat{F}} = \underbrace{\begin{matrix} & \diamond & \triangle & \star \\ P & \begin{bmatrix} 0.82 & 0 & 0 \end{bmatrix} \\ M_1 & \begin{bmatrix} 0 & 0.24 & 0.64 \end{bmatrix} \\ M_2 & \begin{bmatrix} 0 & 0 & 0.73 \end{bmatrix} \end{matrix}}_{U} \underbrace{\begin{matrix} & \diamond & \triangle & \star \\ & \begin{bmatrix} 1 & 0 & 0 \end{bmatrix} \\ & \begin{bmatrix} 1 & 1 & 0 \end{bmatrix} \\ & \begin{bmatrix} 1 & 0 & 1 \end{bmatrix} \end{matrix}}_{B}$$

PMH-TI Problem
Parsimonious Migration History with Tree Inference

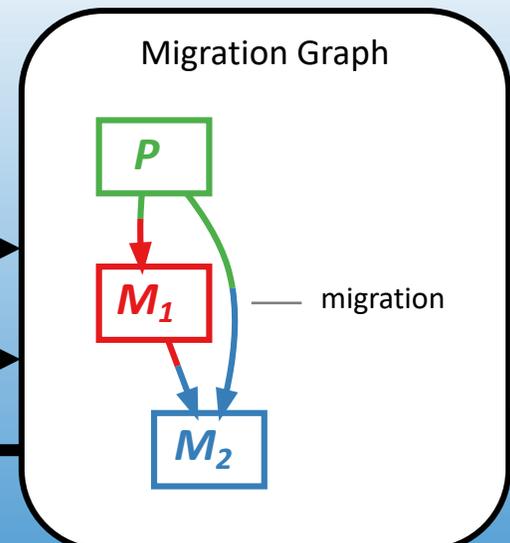
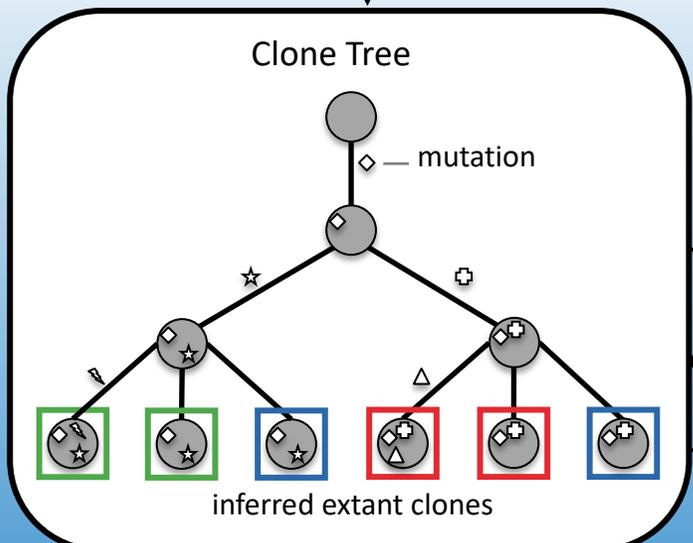
\mathcal{P}

Standard Phylogenetic Techniques



Sequencing and Mutation Calling

Tumor Phylogenetic Techniques



MACHINA

Cell Division and Mutation History

Cell Migration History

Label ancestral vertices by anatomical sites

Resolve clone tree ambiguities

MACHINA accurately infers clone trees and migration histories on simulated data

