

# CS 598MEB

# Computational Cancer Genomics

## Lecture 2

Mohammed El-Kebir

January 17, 2019



# Lecture Outline

- Recap
- Maximum Parsimony
- Two-state Perfect Phylogeny
- Two-state Perfect Phylogeny Mixtures

## **Reading**

- Lecture notes

# Hallmarks of Cancer



Inter-tumor heterogeneity:  
Every tumor is different!

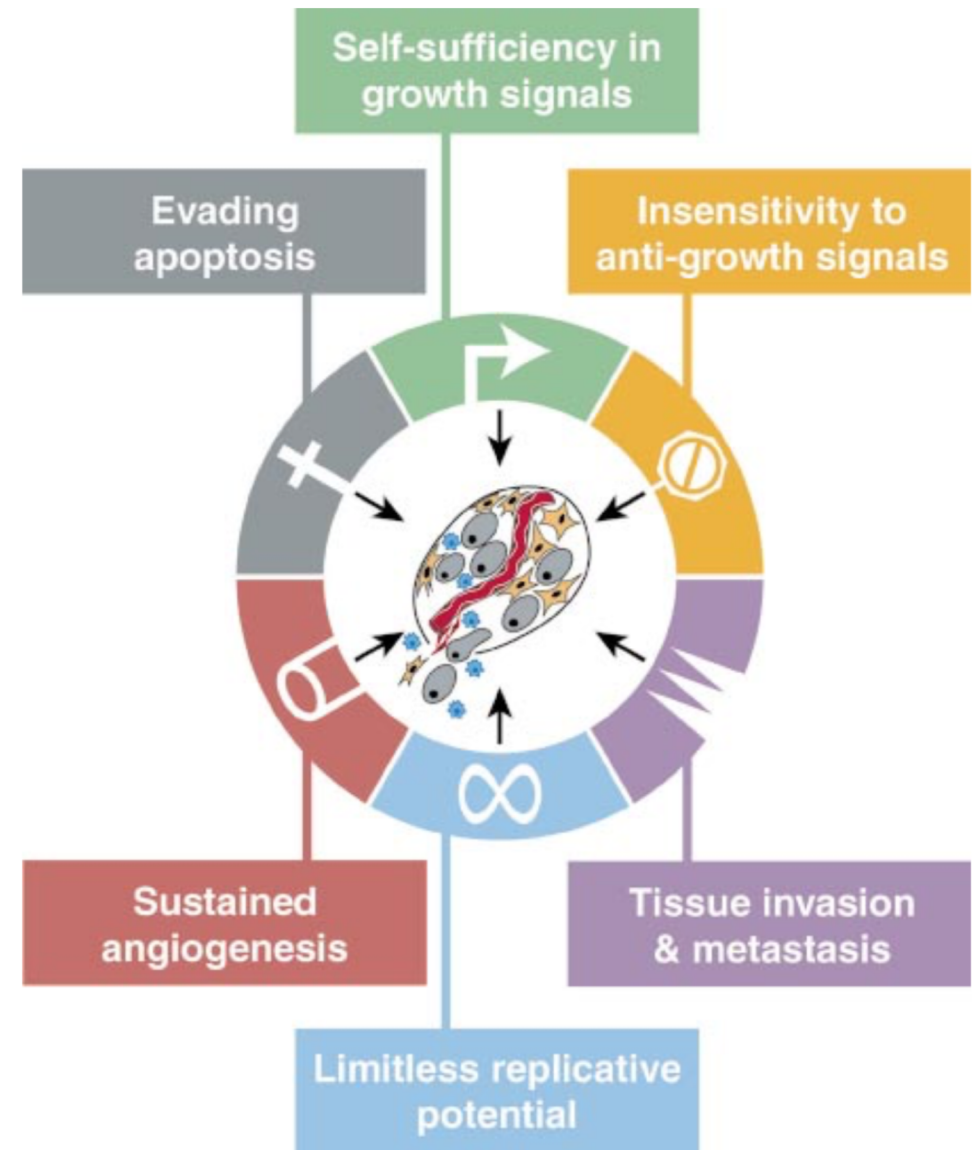
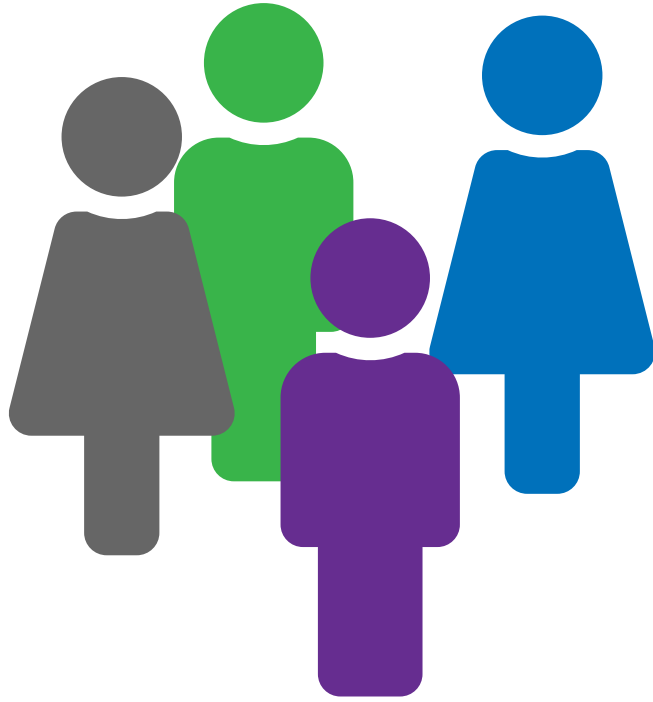


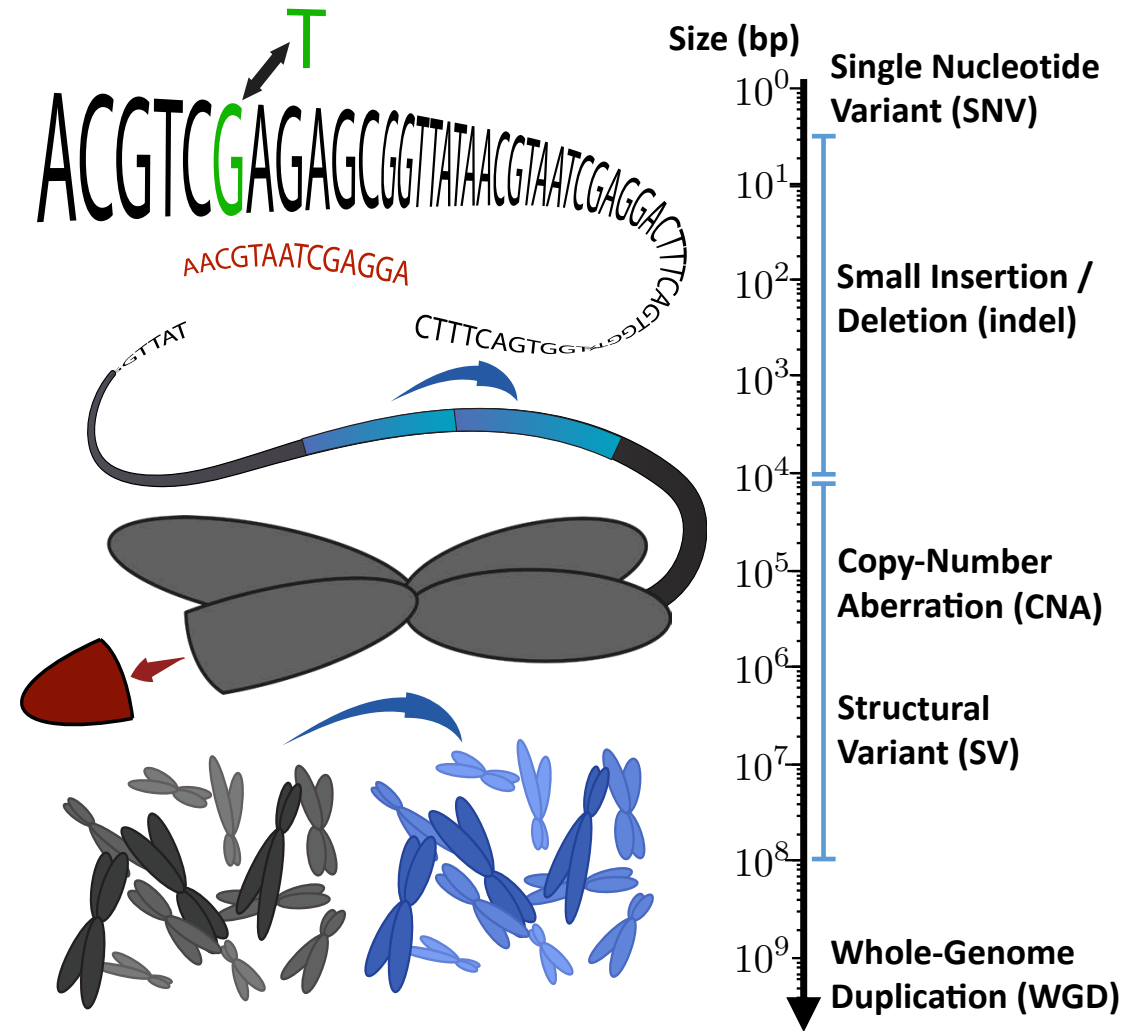
Figure 1. Acquired Capabilities of Cancer

We suggest that most if not all cancers have acquired the same set of functional capabilities during their development, albeit through various mechanistic strategies.

# Cancer is Caused by Somatic Mutations



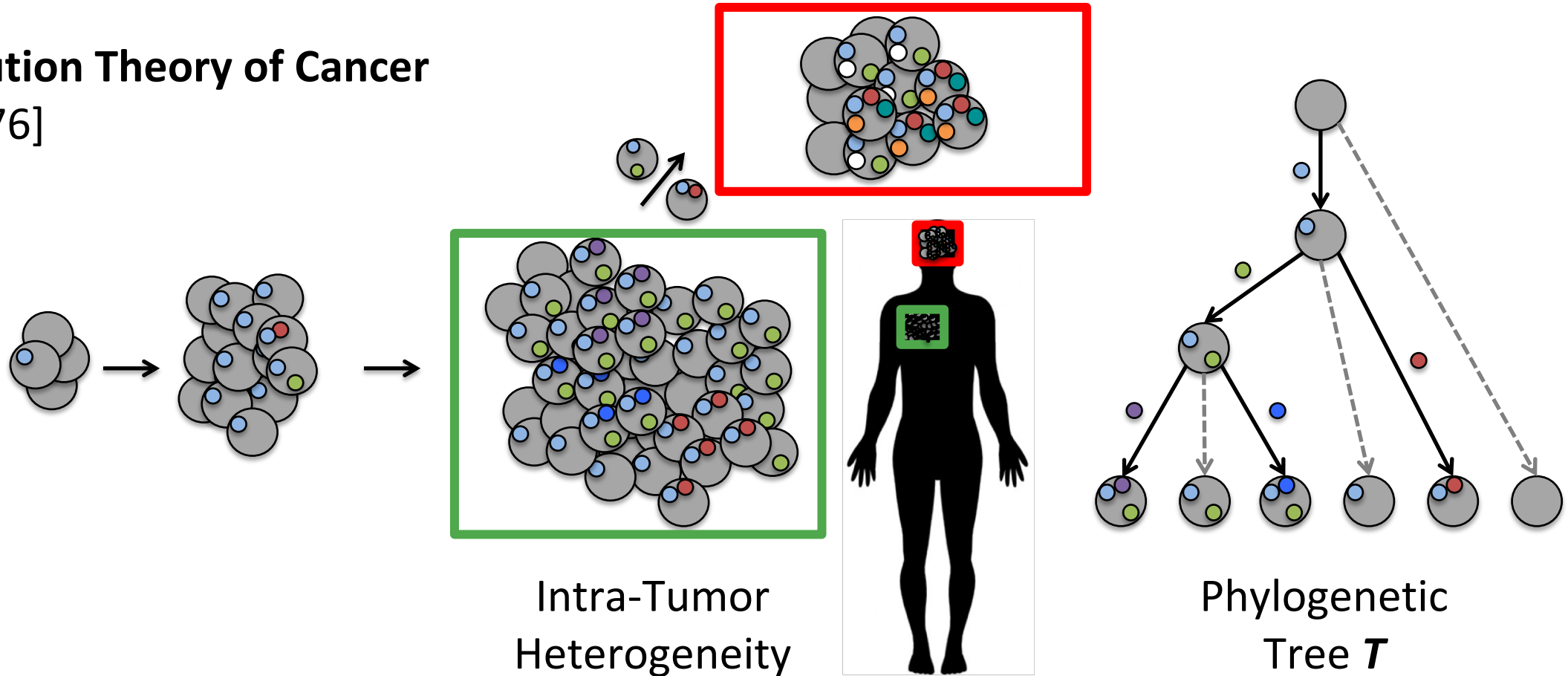
Inter-tumor heterogeneity:  
Every tumor is different!





# Tumorigenesis: Cell Mutation, Division & Migration

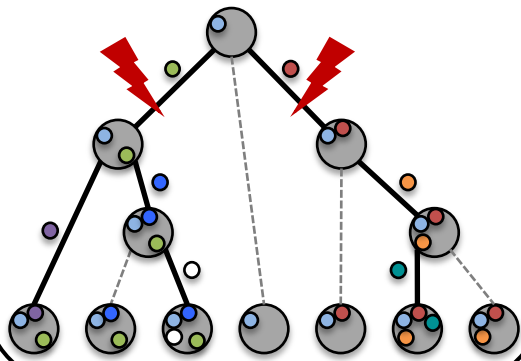
## Clonal Evolution Theory of Cancer [Nowell, 1976]



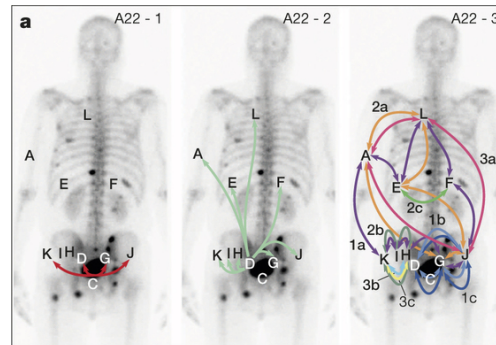
**Intra-tumor heterogeneity:** Every tumor cell is different

# Phylogenies are Key to Understanding Cancer

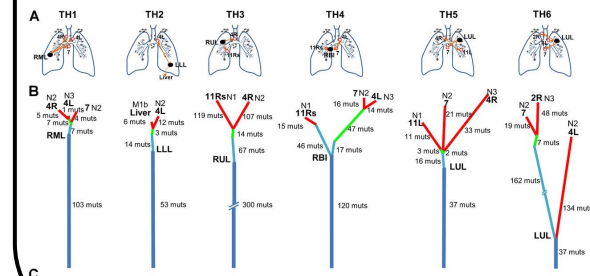
## Identify targets for treatment



## Understand metastatic development



## Recognize common patterns of tumor evolution across patients



These downstream analyses **critically rely** on accurate tumor phylogeny inference

**Key challenge in phylogenetics:**  
Accurate phylogeny inference from data at present time

# Lecture Outline

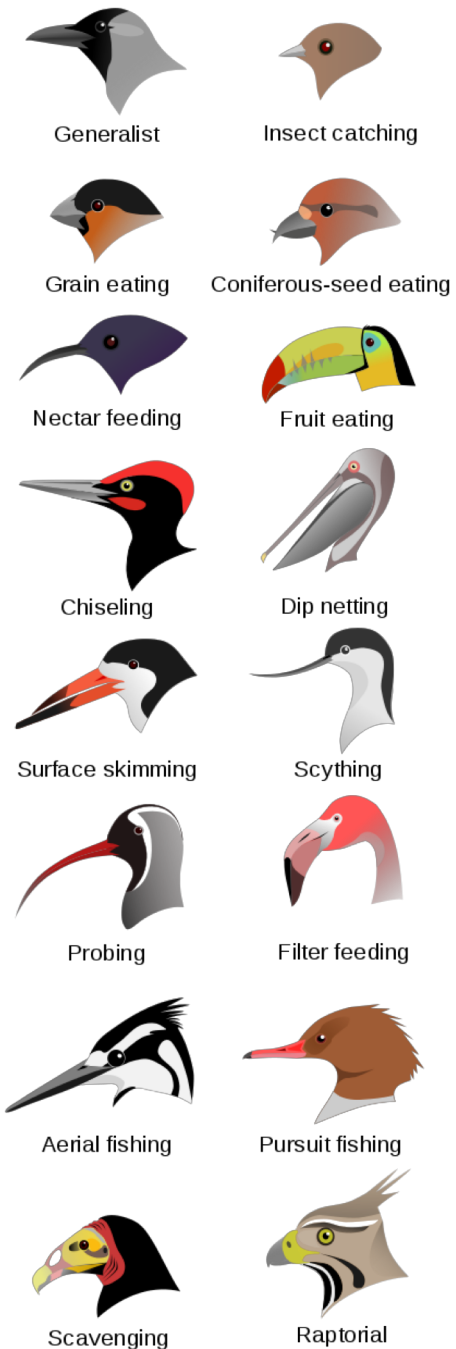
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## Reading

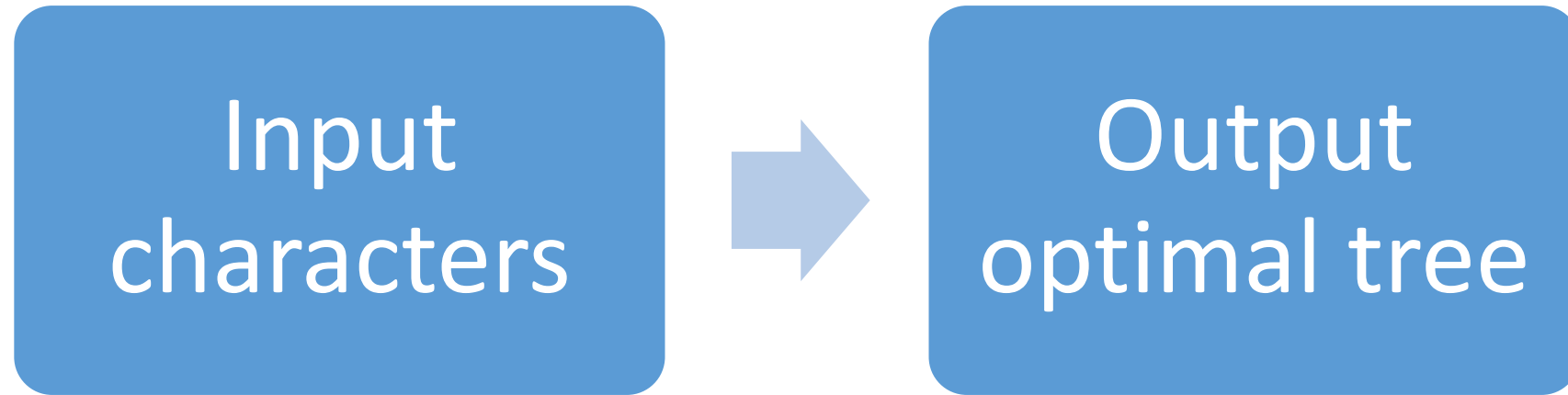
- Lecture notes

# Character-Based Tree Reconstruction

- Characters may be morphological features
  - Shape of beak {generalist, insect catching, ...}
  - Number of legs {2,3,4, ..}
  - Hibernation {yes, no}
- Character may be nucleotides/amino acids
  - {A, T, C, G}
  - 20 amino acids
- Values of a character are called states
  - We assume discrete states

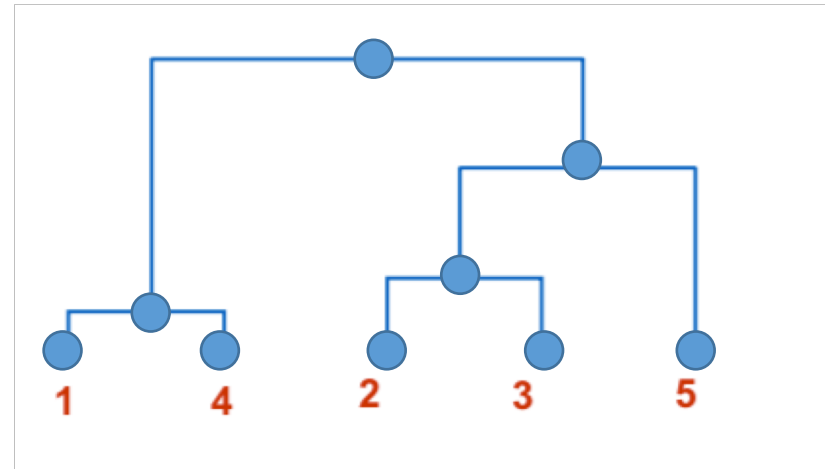


# Character-Based Phylogeny Reconstruction



**Question:** What is optimal?

**Want:** Optimization criterion



# Character-Based Phylogeny Reconstruction

Input  
characters



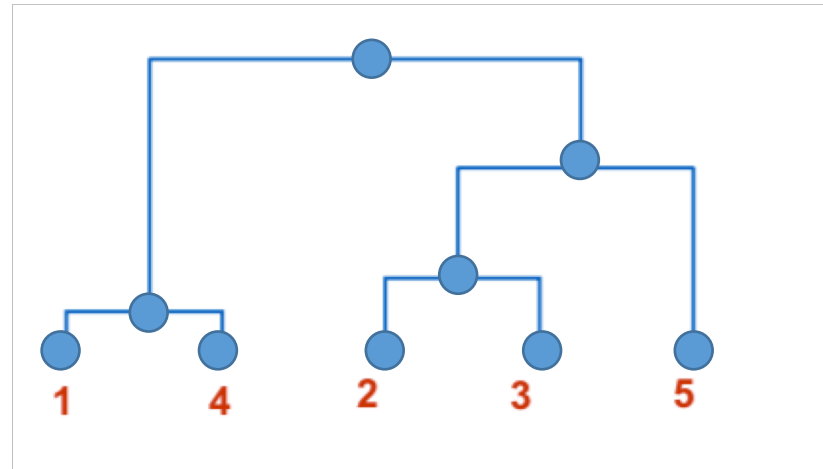
Output  
optimal tree

**Question:** What is optimal?

**Want:** Optimization criterion

**Question:** How to optimize this criterion?

**Want:** Algorithm

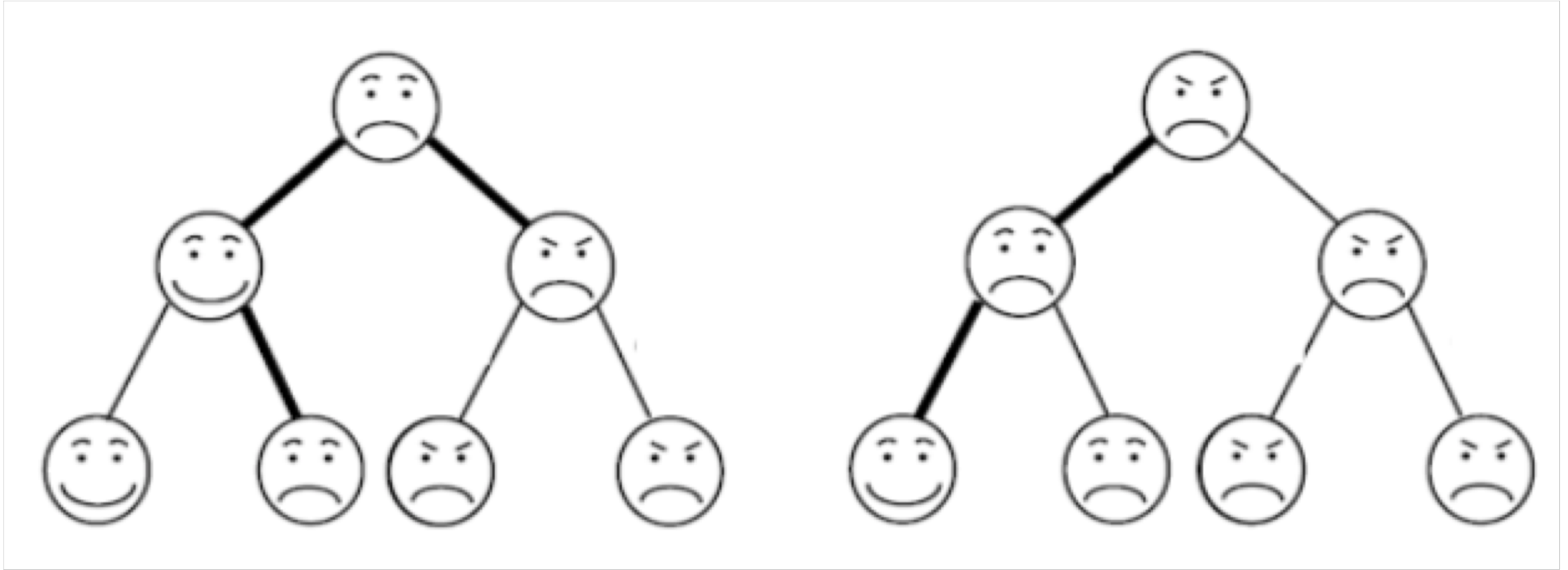


# Character-Based Phylogeny Reconstruction: Input

Characters / states	State 1	State 2
Mouth	Smile	Frown
Eyebrows	Normal	Pointed



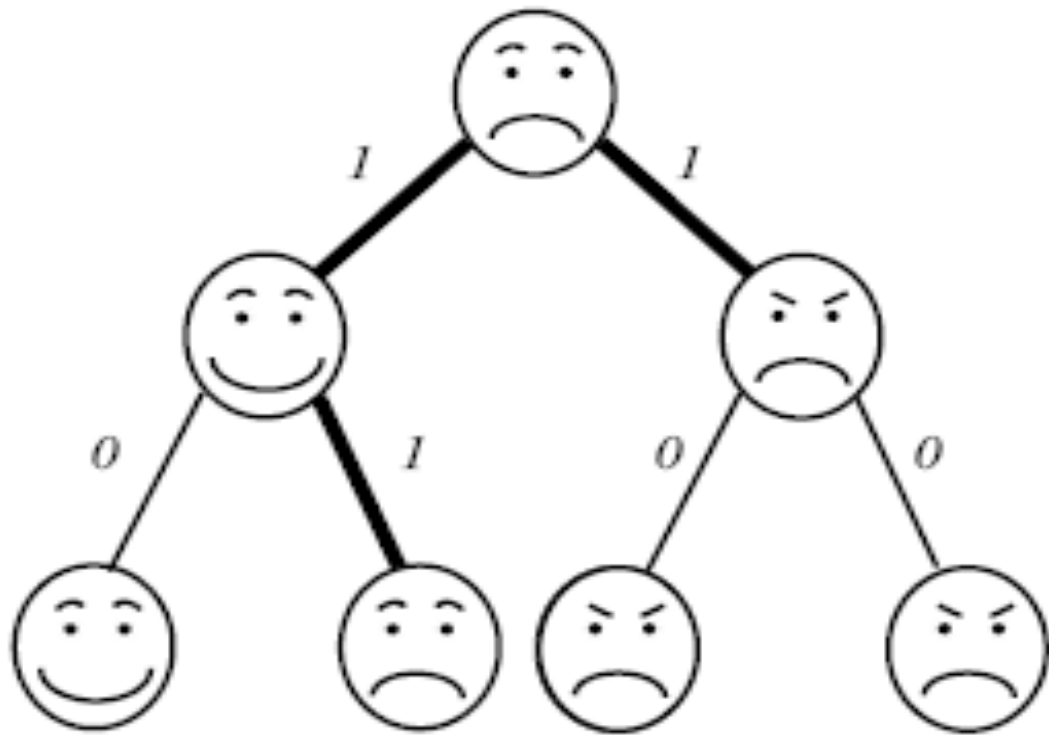
# Character-Based Phylogeny Reconstruction: Criterion



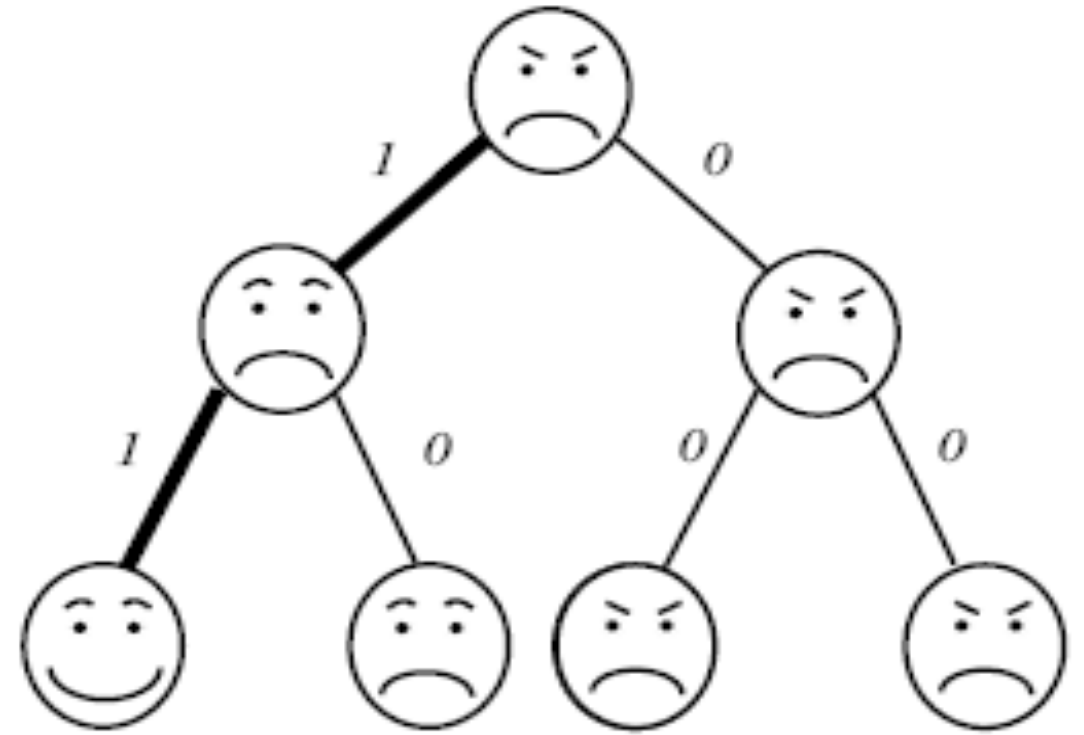
**Question:** Which tree is better?



# Character-Based Phylogeny Reconstruction: Criterion



(a) *Parsimony Score=3*

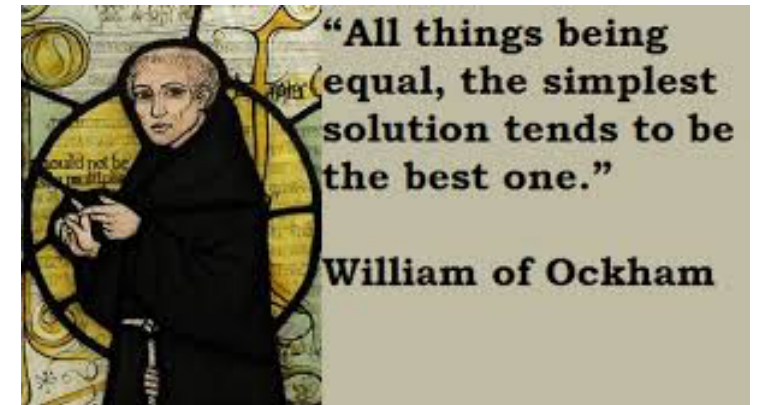


(b) *Parsimony Score=2*

**Parsimony:** minimize number of changes on edges of tree

# Why Parsimony?

- Ockham's razor: “simplest” explanation for data
- Assumes that observed character differences resulted from the fewest possible mutations
- Seeks tree with the lowest **parsimony score**, i.e. the sum of all (costs of) mutations in the tree.



# Binary Characters

		Characters				
		1	2	3	4	5
Species	A	0	1	1	0	0
	B	0	0	1	1	0
	C	1	1	1	1	0
	D	1	1	0	1	1

Characters only have two possible states

Possible Encoding:  
0 : not-mutated  
1 : mutated

Possible Encoding:  
0 : no wings  
1 : wings

# A Small and a Large Problem

## **Small Maximum Parsimony Phylogeny Problem:**

Given  $m \times n$  matrix  $A = [a_{i,j}]$  and tree  $T$  with  $m$  leaves, find assignment of character states to each internal vertex of  $T$  with minimum parsimony score.

## **Large Maximum Parsimony Phylogeny Problem:**

Given  $m \times n$  matrix  $A = [a_{i,j}]$ , find a tree  $T$  with  $m$  leaves labeled according to  $A$  and an assignment of character states to each internal vertex of  $T$  with minimum parsimony score.

# A Small and a Large Problem

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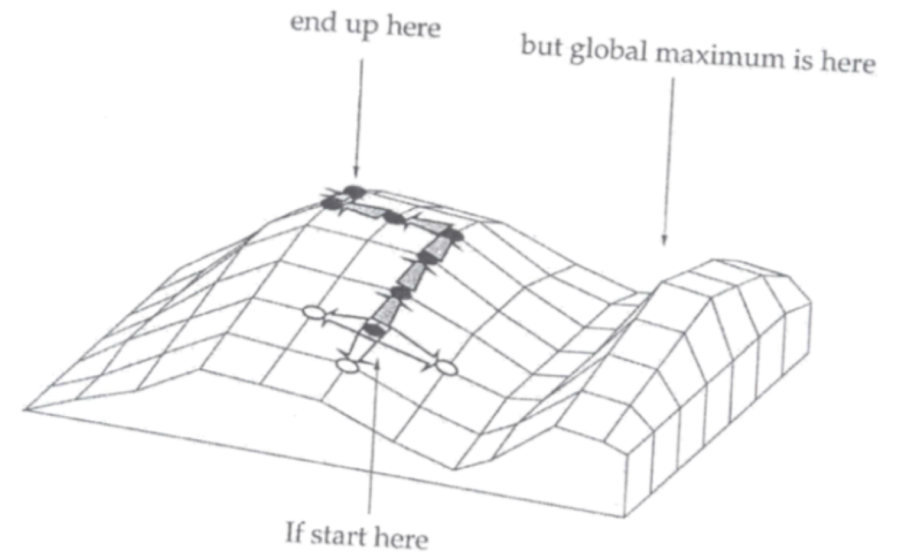
**Question:** Are both problems easy (i.e. in P)?

# Large Maximum Parsimony Phylogeny

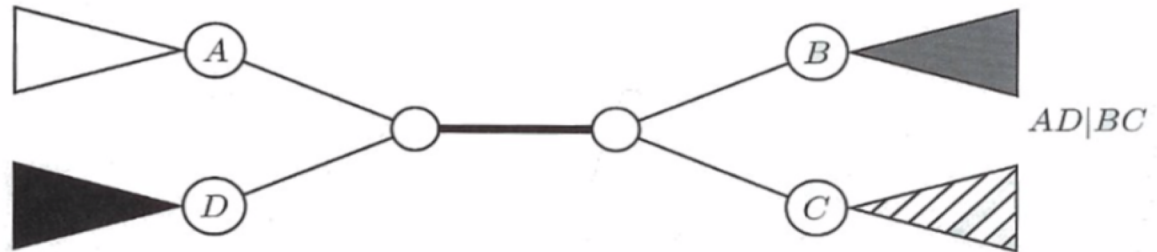
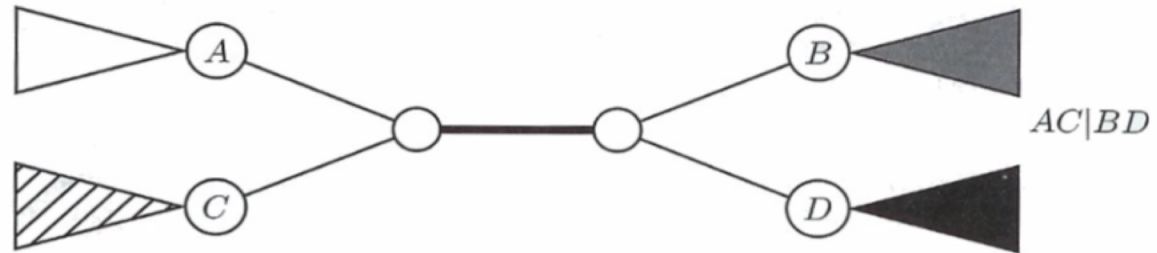
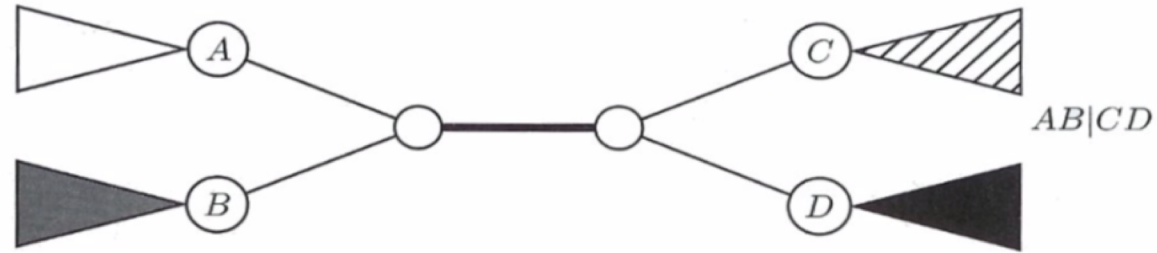
- This problem is NP-hard
- Heuristics using local search (tree moves)

1. Start with an arbitrary tree  $T$ .
2. Check “neighbors” of  $T$ .
3. Move to a neighbor if it provides the best improvement in parsimony/likelihood score.

Caveats:  
Could be stuck in **local** optimum, and not achieve global optimum



# Local Search: Nearest-Neighbor Interchange (NNI)

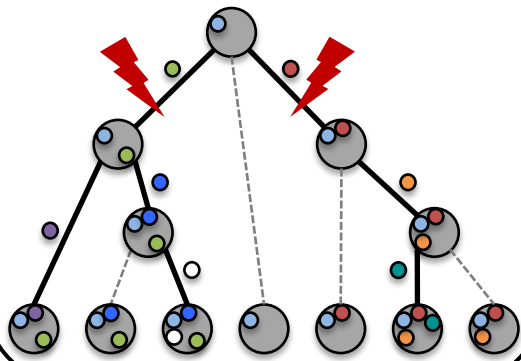


Rearrange four subtrees  
defined by one  
internal edge

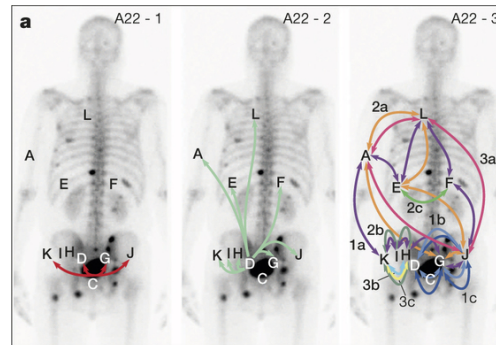
Figure: Jones and Pevzner

# Phylogenies are Key to Understanding Cancer

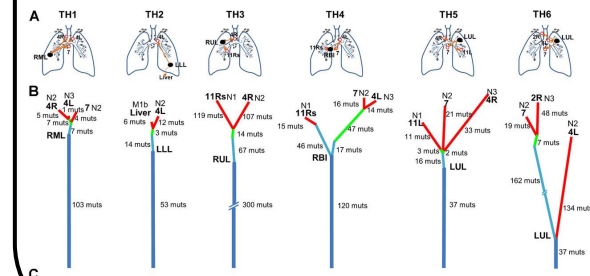
Identify targets for treatment



Understand metastatic development



Recognize common patterns of tumor evolution across patients

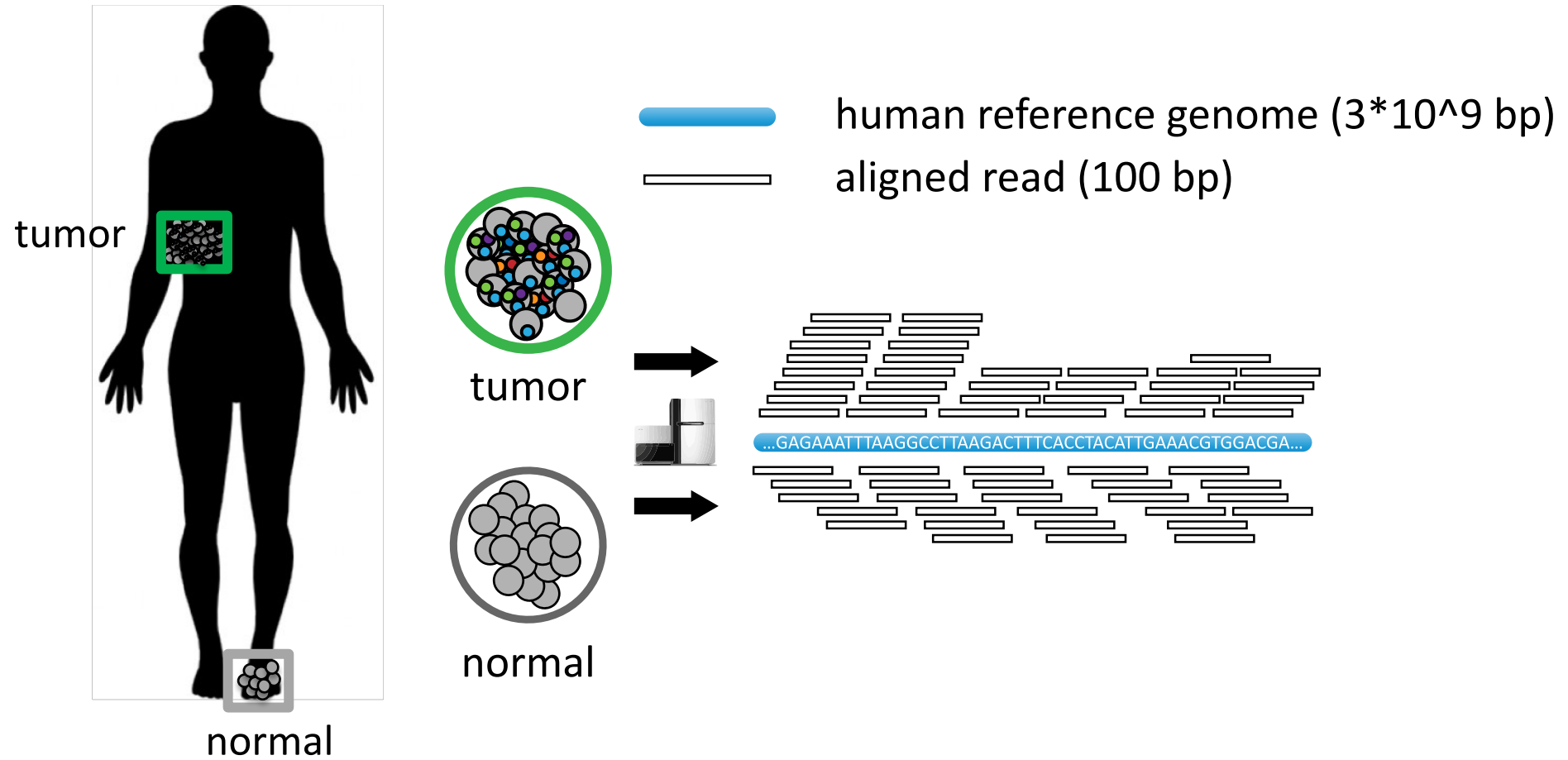


These downstream analyses **critically rely** on accurate tumor phylogeny inference

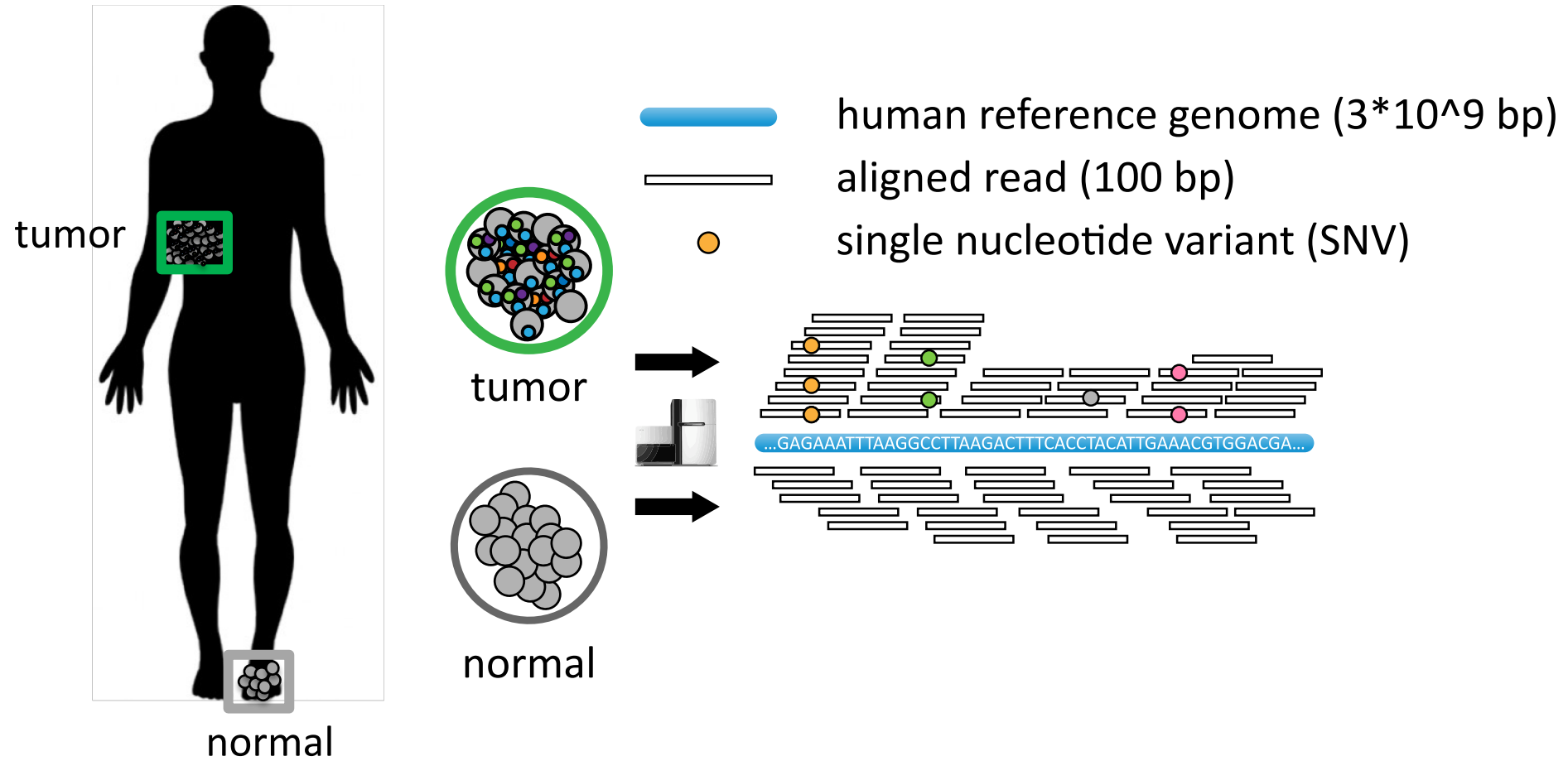
**Key challenge in phylogenetics:**  
Accurate phylogeny inference from data at present time



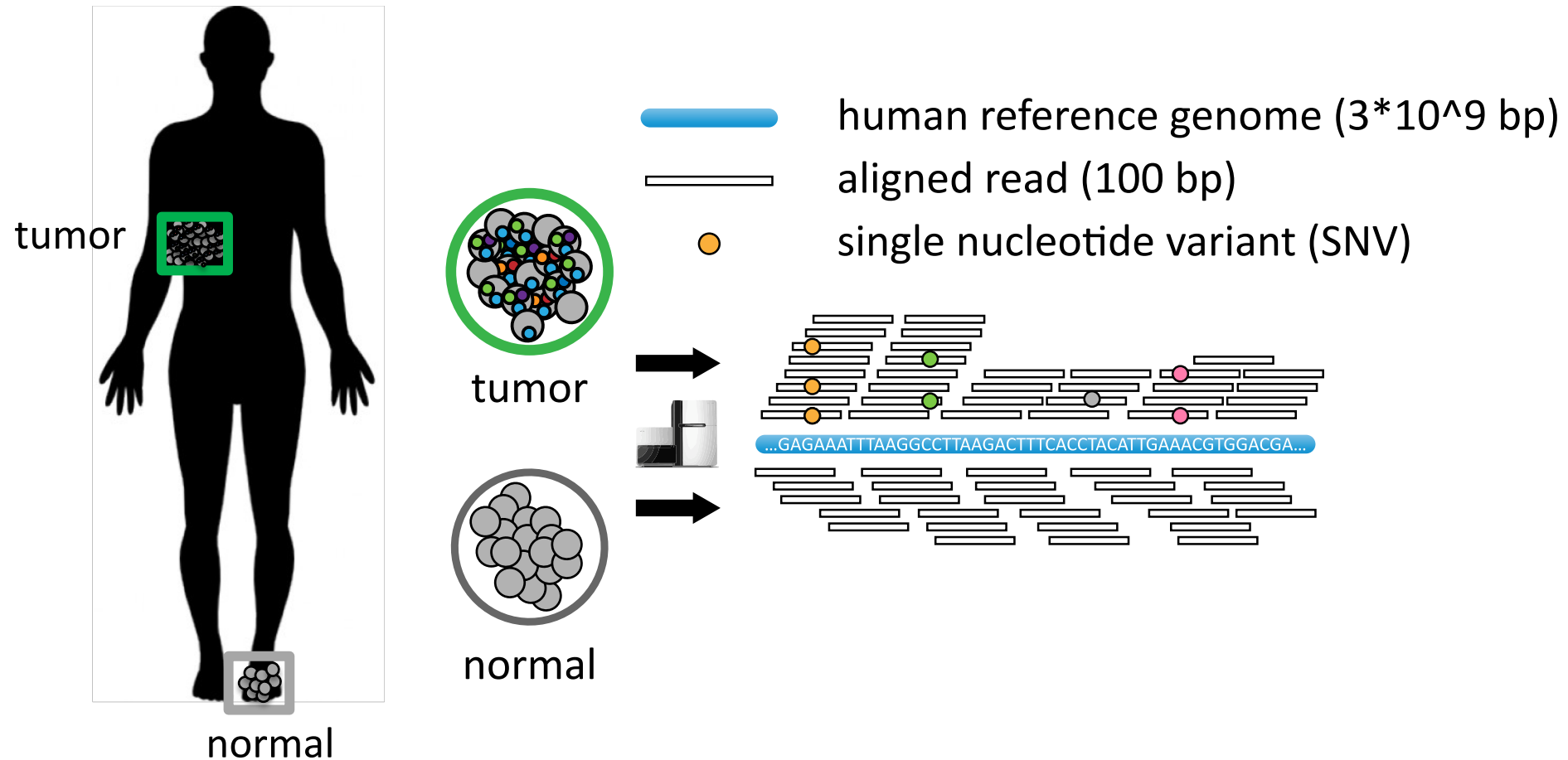
# Additional Challenge in Cancer Phylogenetics



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# Additional Challenge in Cancer Phylogenetics



**Additional challenge in cancer phylogenetics:**  
Phylogeny inference from **mixed bulk samples** at present time







# Tumor Phylogeny Inference

## Metastatic Colorectal Cancer (Patient CRC2)

[Kim et al., *Clin Cancer Res* 21(19), 2015]:

- 5 primary samples (P1-P5)
- 2 metastases (M1-M2)
- 412 single-nucleotide variants (SNVs)

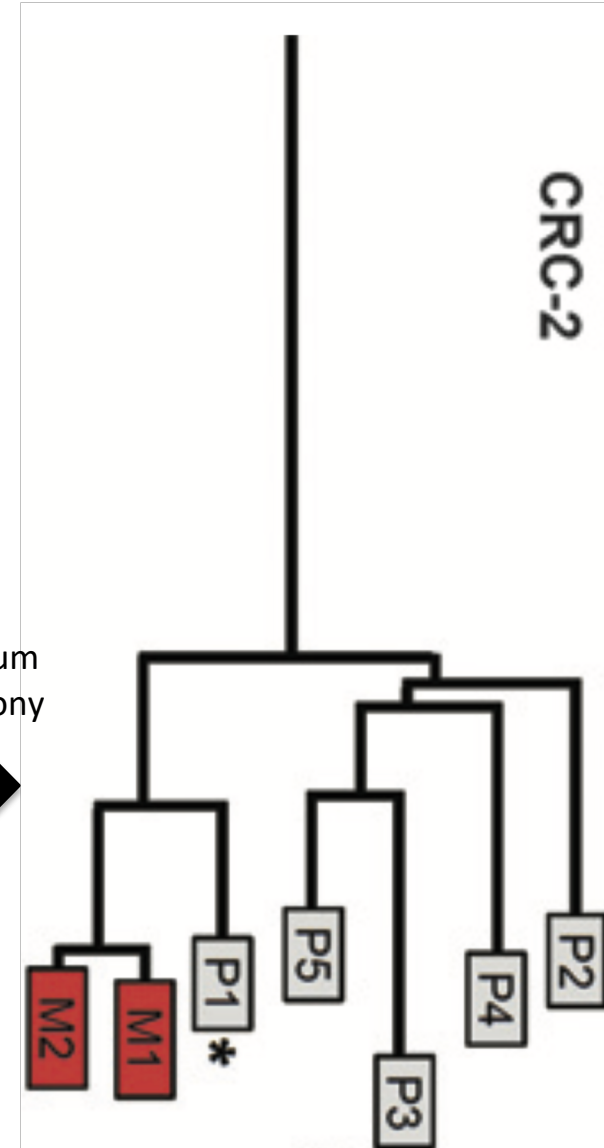
$n$  mutations

						
<b>P1</b>	1	1	1	0	0	0
<b>P2</b>	1	1	0	1	0	0
<b>P3</b>	1	0	1	0	1	1
<b>P4</b>	0	1	1	0	0	0
<b>P5</b>	0	1	0	1	0	1
<b>M1</b>	1	1	0	0	1	0
<b>M2</b>	0	1	1	1	1	1

$m$  samples

Binary Matrix **B**

Maximum  
Parsimony









# Tumor Phylogeny Inference

## Metastatic Colorectal Cancer (Patient CRC2)

[Kim et al., *Clin Cancer Res* 21(19), 2015]:

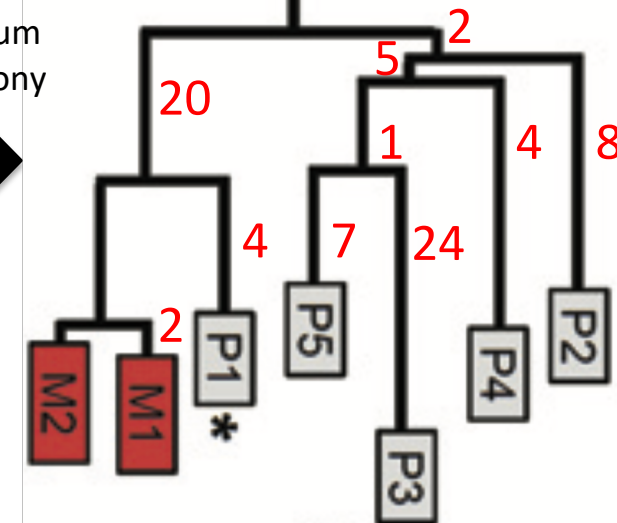
- 5 primary samples (P1-P5)
- 2 metastases (M1-M2)
- 412 single-nucleotide variants (SNVs)
- 41 mutate more than once (**homoplasy**)

$m$  samples

		$n$ mutations					
							
P1	1	1	1	1	0	0	0
P2	1	1	0	1	0	0	0
P3	1	0	1	0	1	1	1
P4	0	1	1	0	0	0	0
P5	0	1	0	1	0	0	1
M1	1	1	0	0	1	0	0
M2	0	1	1	1	1	1	1

Binary Matrix  $B$

Maximum  
Parsimony

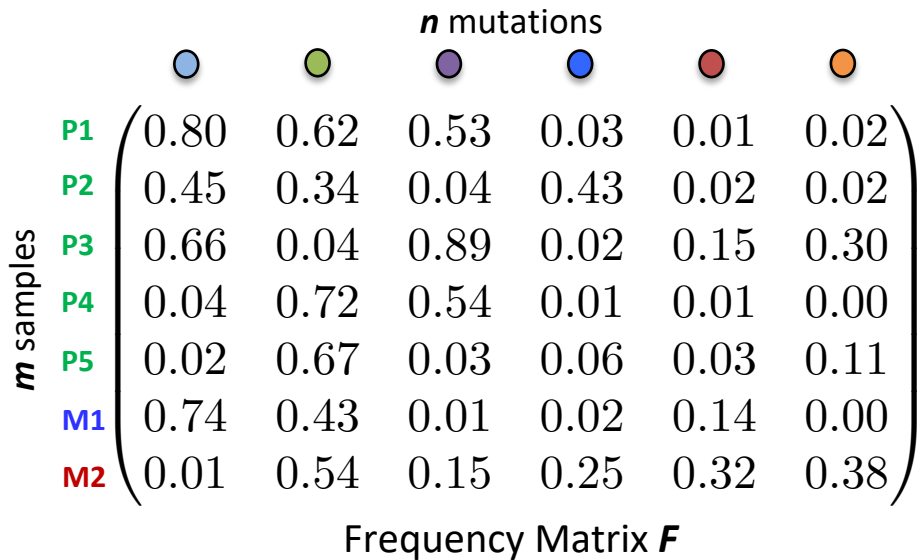


# Heuristic for Tumor Phylogeny Inference

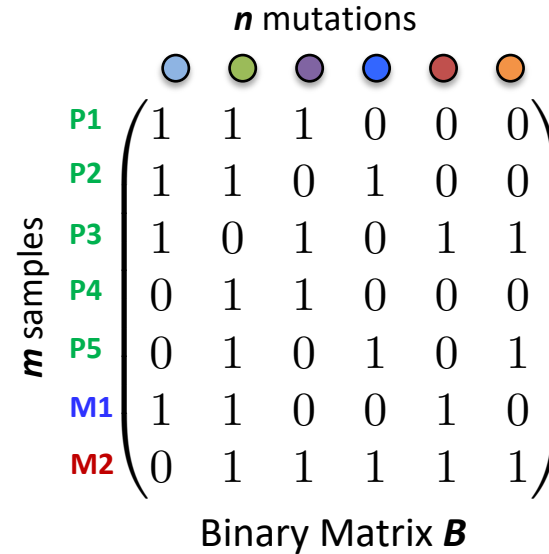
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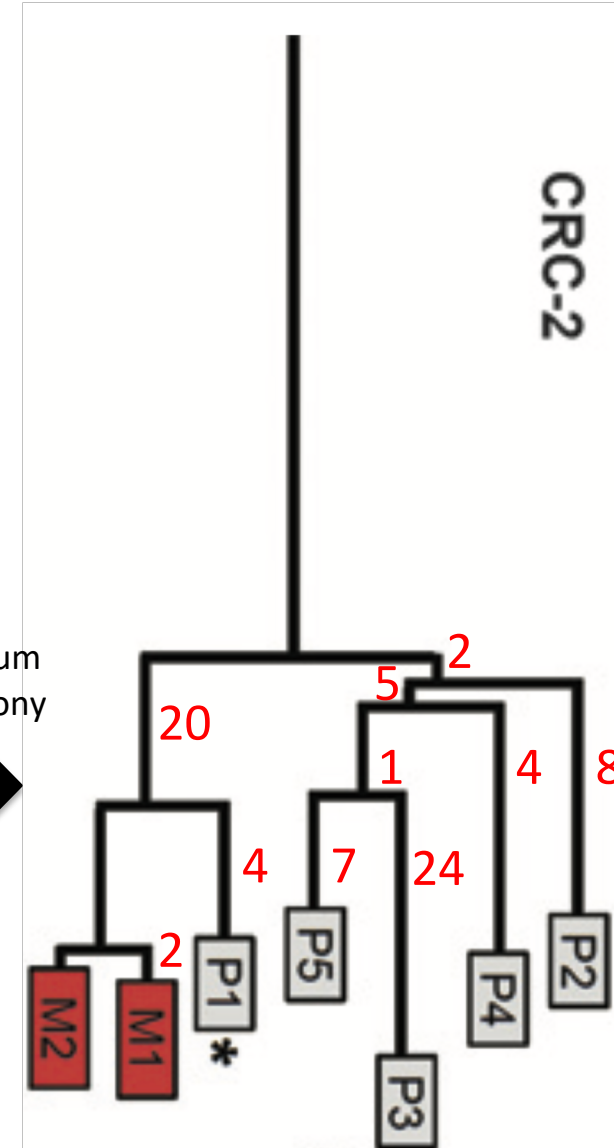
- 5 primary samples (P1-P5)
- 2 metastases (M1-M2)
- 412 single-nucleotide variants (SNVs)
- 41 mutate more than once (**homoplasy**)



## Discretize



## Maximum Parsimony



Resulting **sample tree** is **not** representative of the division/mutation history or the migration history

# Lecture Outline

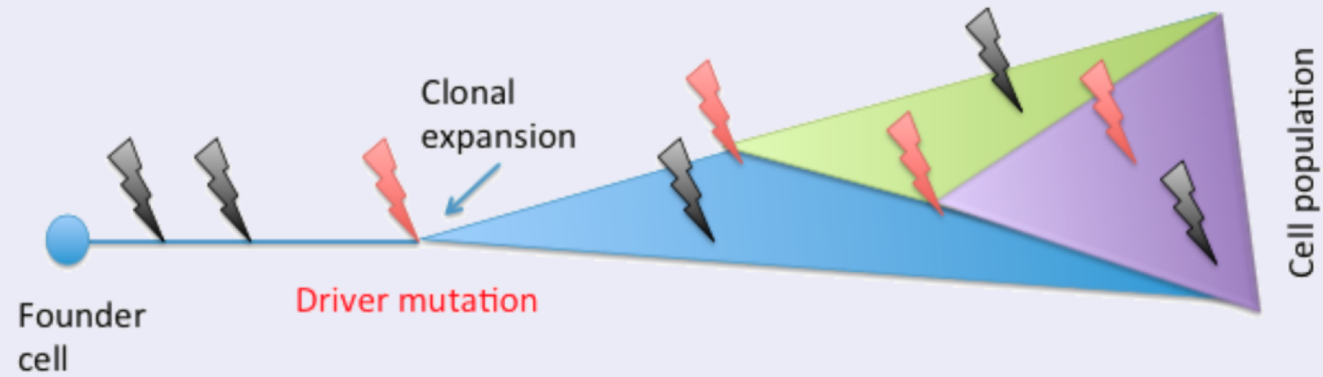
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# Somatic Mutations and Cancer

## Clonal theory of cancer (Nowell, 1976)

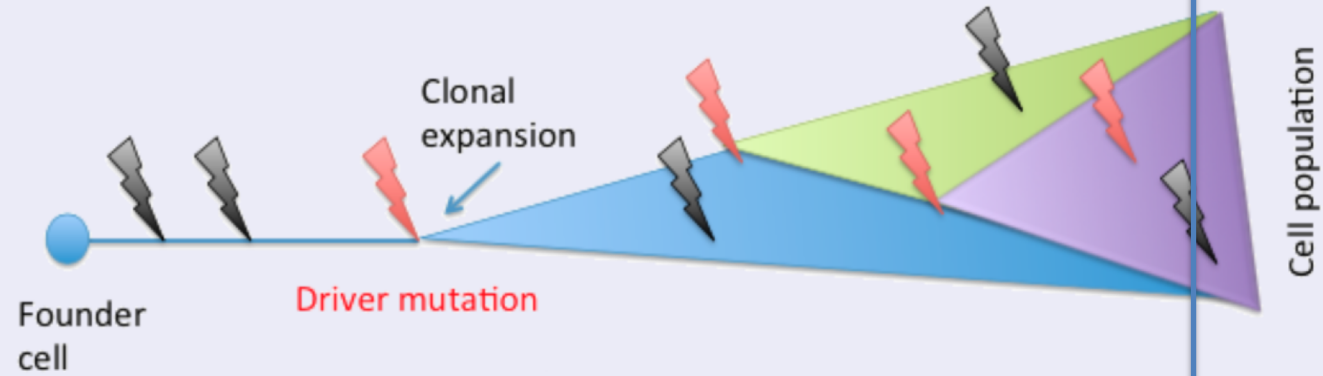


“typical tumor”:     ~10 driver mutations  
                              100’s – 1000’s of passenger mutations



# Somatic Mutations and Cancer

## Clonal theory of cancer (Nowell, 1976)



“typical tumor”: ~10 driver mutations  
100’s – 1000’s of passenger mutations

THE CANCER GENOME ATLAS



International  
Cancer Genome  
Consortium



Sequence genome

# Progression of Somatic Mutations

## Single nucleotide mutation

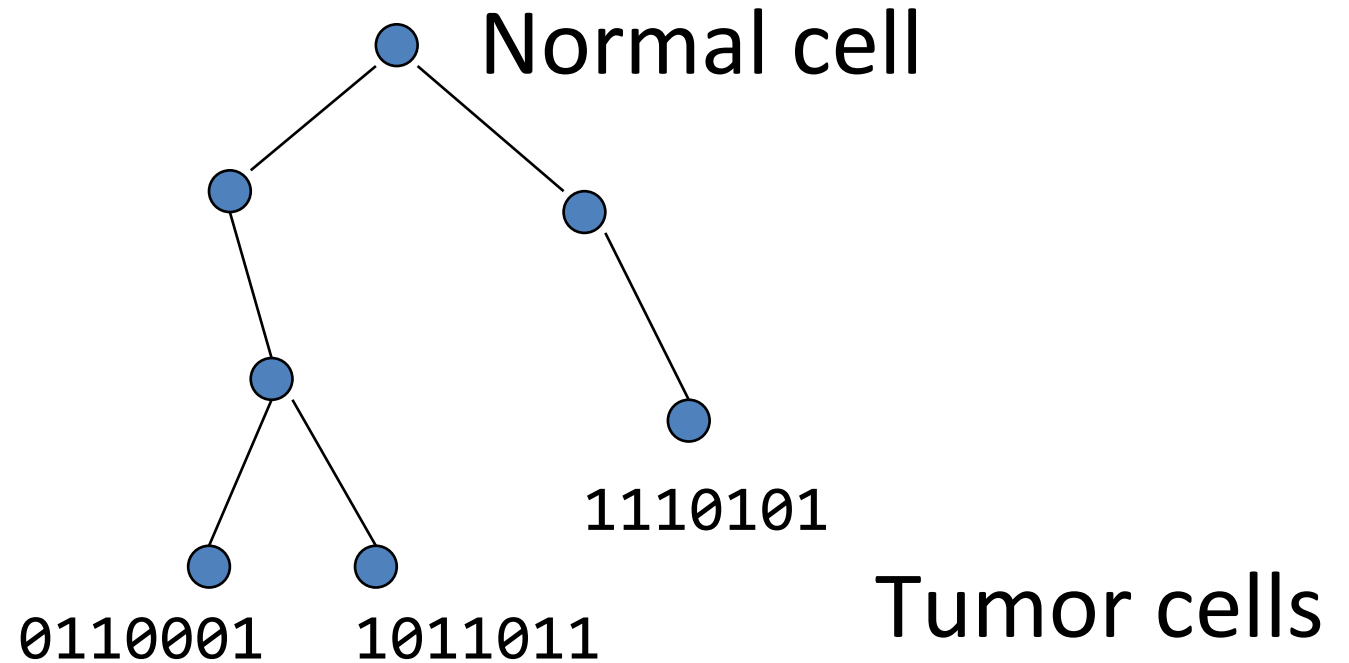
... CGT**A**TTAG ...



... CGT**C**TTAG ...

0 = normal

1 = mutated



Root is the normal, founder cell and leaves are cells in tumor.

# Progression of Somatic Mutations

## Single nucleotide mutation

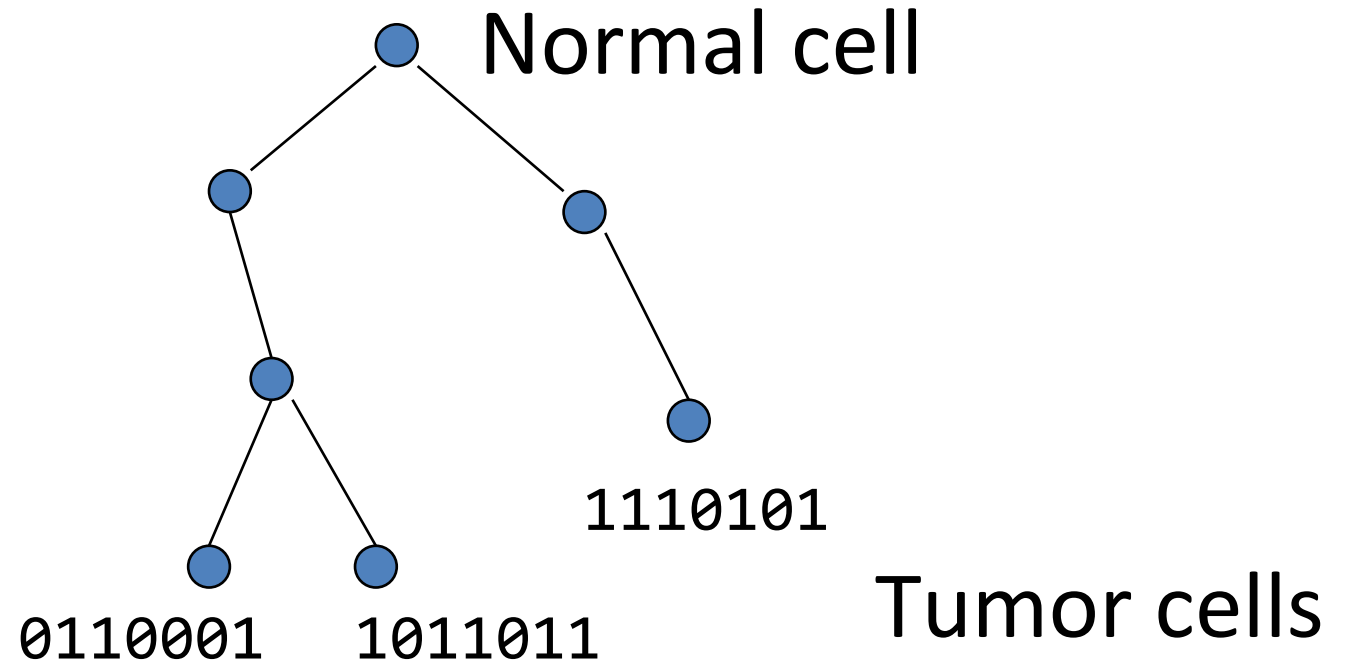
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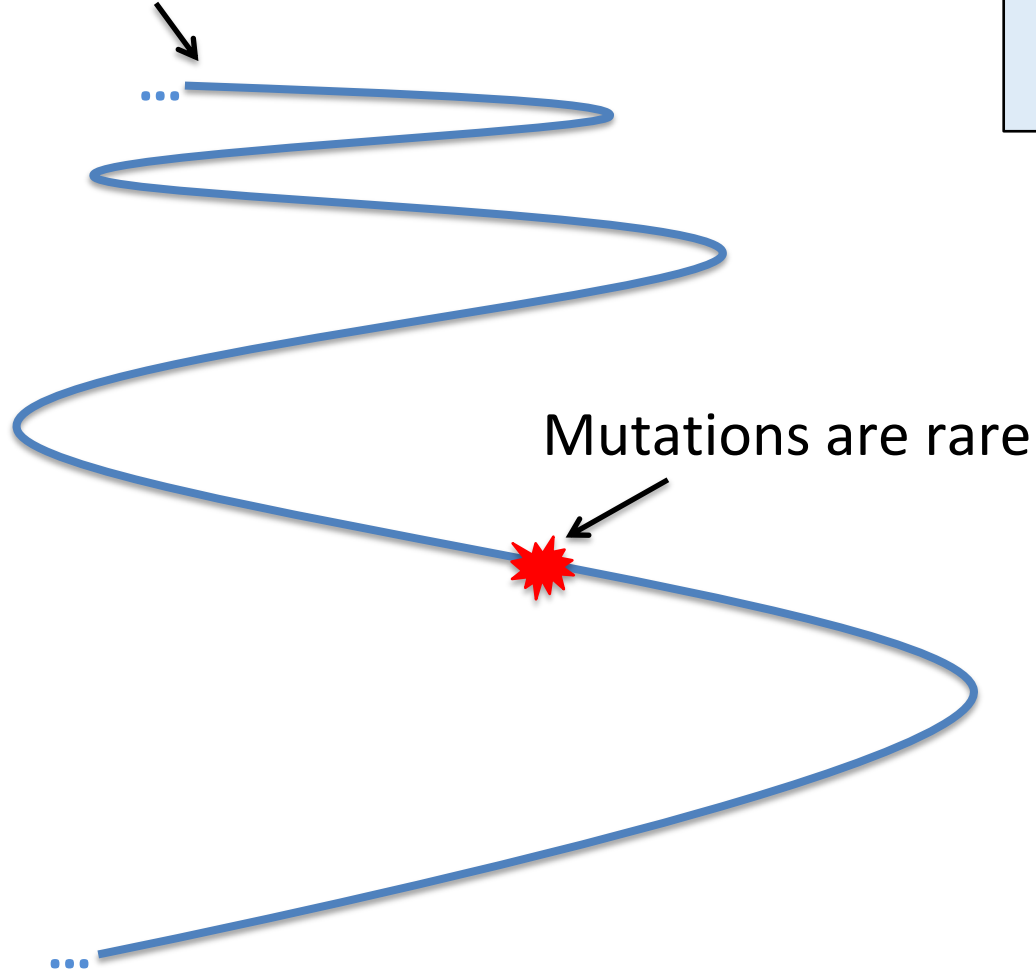


Root is the normal, founder cell and leaves are cells in tumor.

**Infinite sites assumption:** each locus mutates only once.







# Infinite Sites Model

The genome is large



[Kimura, 1969]

**Infinite sites model:** multiple mutations never occur at the same position

		Mutated Loci					
							
Species (cancer cells)	A	0	0	0	0	1	1
	B	0	0	0	1	1	1
	C	0	0	1	0	1	0
	D	1	0	0	0	0	0
	E	1	1	0	0	0	0

1: mutated

0: not

All sites are bi-allelic: mutated or not.

# Two-state Perfect Phylogeny

Matrix  $M \in \{0, 1\}^{n \times m}$  has  $n$  **taxa** and  $m$  **characters**

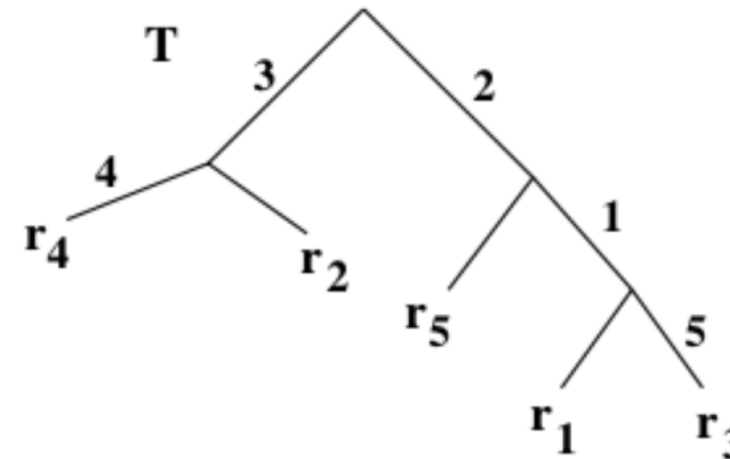
- Taxon  $f$  has state 1 for character  $c$   
 $\Leftrightarrow f$  **possesses** character  $c$

	$c_1$	$c_2$	$c_3$	$c_4$	$c_5$
$r_1$	1	1	0	0	0
$r_2$	0	0	1	0	0
$r_3$	1	1	0	0	1
$r_4$	0	0	1	1	0
$r_5$	0	1	0	0	0

## Definition

A **perfect phylogeny** for  $M$  is a rooted tree  $T$  with  $n$  leaves such that:

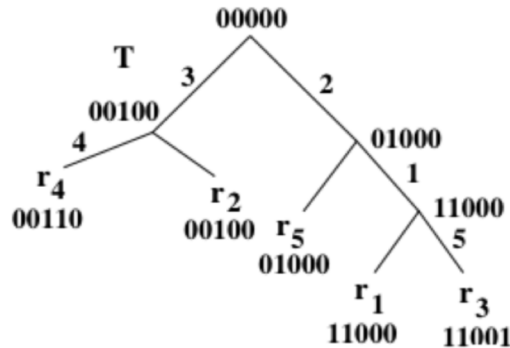
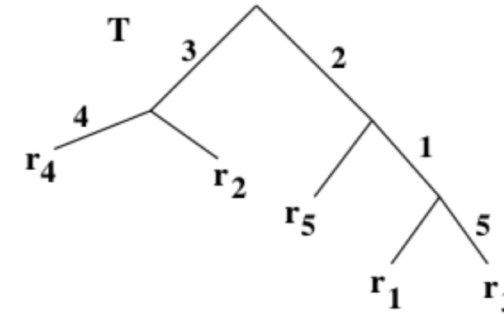
- 1 Each taxon labels only one leaf
- 2 Each character labels only one edge
- 3 Character possessed by a taxon are on unique path to root



Root node is all zero ancestor

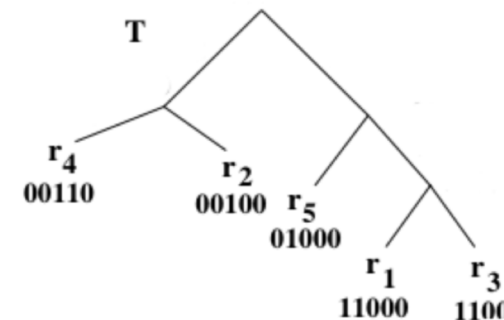
# Two-state Perfect Phylogeny – Alternative Definitions

- ① Each taxon labels exactly one leaf
- ② Each character labels exactly one edge
- ③ Character possessed by a taxon are on unique path to root



- ① Each taxon labels exactly one leaf
- ② Each node is labeled by  $\{0, 1\}^m$
- ③ Nodes labeled with state  $i$  for character  $c$  form a connected subtree

- ① Each taxon labels exactly one leaf
- ②  $T_c(i)$  is smallest subtree connecting all leaves labeled with state  $i$  for character  $c$
- ③  $T_c(0)$  and  $T_c(1)$  are disjoint for all  $c$



# Two-state Perfect Phylogeny Problem

## Input:

Matrix  $M \in \{0, 1\}^{n \times m}$  has  $n$  **taxa** and  $m$  **characters**

- Taxon  $f$  has state 1 for character  $c$   
 $\Leftrightarrow f$  **possesses** character  $c$

	$c_1$	$c_2$	$c_3$	$c_4$	$c_5$
$r_1$	1	1	0	0	0
$r_2$	0	0	1	0	0
$r_3$	1	1	0	0	1
$r_4$	0	0	1	1	0
$r_5$	0	1	0	0	0

## Problem

*Given  $M \in \{0, 1\}^{n \times m}$  does  $M$  have a perfect phylogeny?*

# Try it yourself!

Only one of these matrices can be used to build a perfect phylogeny.

- (1) As a group, **decide on an approach** to try to determine which one is which.
- (2) Try out your approach to see if you can construct the tree.
- (3) What did you learn from your attempt?

$M_1 =$

Species	Characters				
	$C_1$	$C_2$	$C_3$	$C_4$	$C_5$
A	0	1	0	0	0
B	0	0	1	0	0
C	1	1	0	0	0
D	0	0	1	1	0
E	1	1	0	0	1

$M_2 =$

Species	Characters				
	$C_1$	$C_2$	$C_3$	$C_4$	$C_5$
A	0	0	1	1	0
B	0	0	1	0	1
C	1	1	0	0	1
D	1	1	0	0	0
E	0	1	0	0	1





# The Perfect Phylogeny Problem – Preliminaries

## Problem

Given  $M \in \{0, 1\}^{n \times m}$  does  $M$  have a perfect phylogeny?

## Definition

$I(c)$  is the set of taxa that possess character  $c$ ; and  $\sigma(f)$  is the set of characters possessed by taxon  $f$ .

	$c_1$	$c_2$	$c_3$	$c_4$	$c_5$		$c_1$ (2)	$c_2$ (1)	$c_3$ (3)	$c_4$ (5)	$c_5$ (4)
$r_1$	1	1	0	0	0	$\Rightarrow$	1	1	0	0	0
$r_2$	0	0	1	0	0		0	0	1	0	0
$r_3$	1	1	0	0	1		1	1	0	1	0
$r_4$	0	0	1	1	0		0	0	1	0	1
$r_5$	0	1	0	0	0		1	0	0	0	0

$$I(c_1) = \{r_1, r_3\}$$

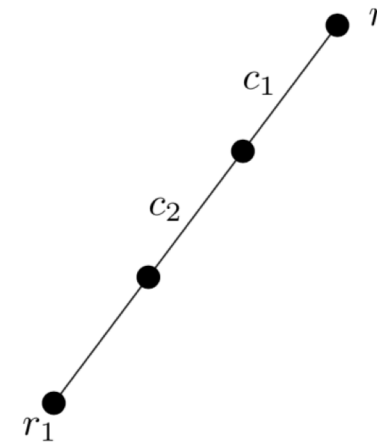
$$\sigma(r_1) = \{c_1, c_2\}$$

Sort columns of  $M$  s.t.  $c < d$  iff  $|I(c)| \geq |I(d)|$ . Break ties arbitrarily.

- Consider rows of  $M$  iteratively
  - $T_i$  is tree of first  $i$  rows of  $M$
- $T_1$  is a path graph
  - Terminal nodes  $r$  and 1
  - $|\sigma(1)| + 1$  edges labeled by  $\sigma(1)$

$$c < d \text{ iff } |I(c)| \geq |I(d)|$$

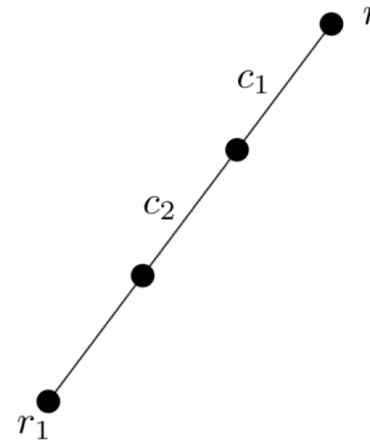
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  - $T_i$  is tree of first  $i$  rows of  $M$
- $T_1$  is a path graph
  - Terminal nodes  $r$  and 1
  - $|\sigma(1)| + 1$  edges labeled by  $\sigma(1)$
- $T_{i+1}$  is a supertree of  $T_i$ 
  - Let  $v$  be last node on walk from  $r$  matching characters  $\sigma(i+1)$ 
    - ★ Character  $d$  is the last match
    - ★ Unmatched characters  $\tau(i+1)$

$$c < d \text{ iff } |I(c)| \geq |I(d)|$$

	$c_1$	$c_2$	$c_3$	$c_4$	$c_5$
$r_1$	1	1	0	0	0
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$r_4$	0	0	1	0	1
$r_5$	1	0	0	0	0



- Consider rows of  $M$  iteratively

- $T_i$  is tree of first  $i$  rows of  $M$

- $T_1$  is a path graph

- Terminal nodes  $r$  and 1

- $|\sigma(1)| + 1$  edges labeled by  $\sigma(1)$

- $T_{i+1}$  is a supertree of  $T_i$

- Let  $v$  be last node on walk from  $r$  matching characters  $\sigma(i+1)$

- ★ Character  $d$  is the last match

- ★ Unmatched characters  $\tau(i+1)$

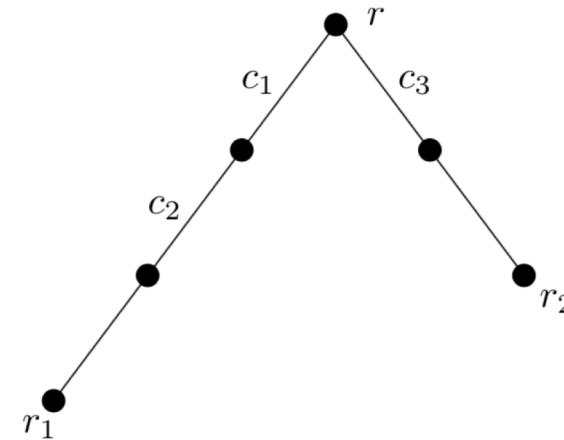
- Extend  $T_i$  with path  $\Pi$

- ★  $\Pi$  has terminals  $v$  and  $i+1$

- ★  $\Pi$  has  $|\tau(i+1)| + 1$  edges labeled by  $\tau(i+1)$

$$c < d \text{ iff } |I(c)| \geq |I(d)|$$

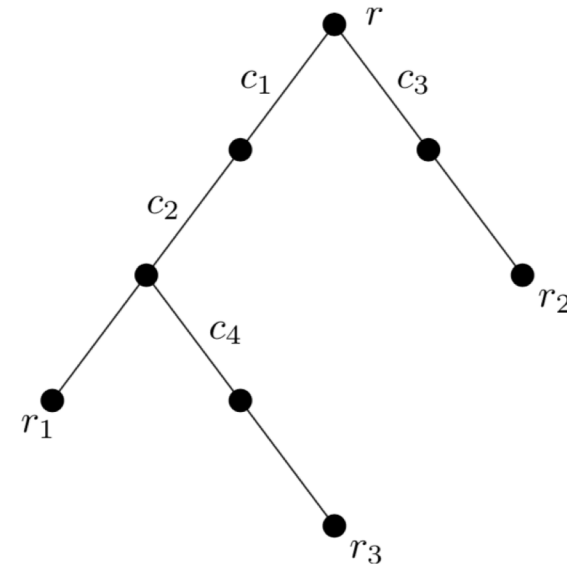
	$c_1$	$c_2$	$c_3$	$c_4$	$c_5$
$r_1$	1	1	0	0	0
$r_2$	0	0	1	0	0
$r_3$	1	1	0	1	0
$r_4$	0	0	1	0	1
$r_5$	1	0	0	0	0



- Consider rows of  $M$  iteratively
  - $T_i$  is tree of first  $i$  rows of  $M$
- $T_1$  is a path graph
  - Terminal nodes  $r$  and 1
  - $|\sigma(1)| + 1$  edges labeled by  $\sigma(1)$
- $T_{i+1}$  is a supertree of  $T_i$ 
  - Let  $v$  be last node on walk from  $r$  matching characters  $\sigma(i + 1)$ 
    - ★ Character  $d$  is the last match
    - ★ Unmatched characters  $\tau(i + 1)$
  - Extend  $T_i$  with path  $\Pi$ 
    - ★  $\Pi$  has terminals  $v$  and  $i + 1$
    - ★  $\Pi$  has  $|\tau(i + 1)| + 1$  edges labeled by  $\tau(i + 1)$

$$c < d \text{ iff } |I(c)| \geq |I(d)|$$

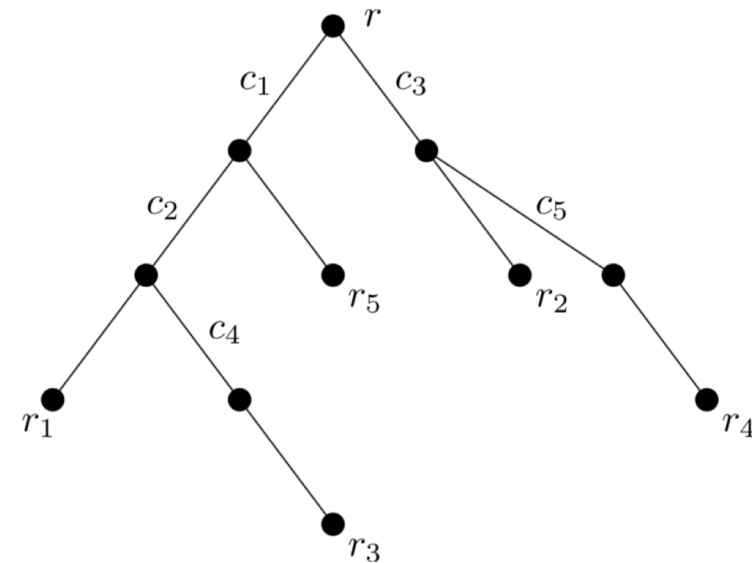
	$c_1$	$c_2$	$c_3$	$c_4$	$c_5$
$r_1$	1	1	0	0	0
$r_2$	0	0	1	0	0
$r_3$	1	1	0	1	0
$r_4$	0	0	1	0	1
$r_5$	1	0	0	0	0



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	$c_1$	$c_2$	$c_3$	$c_4$	$c_5$
$r_1$	1	1	0	0	0
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$r_3$	1	1	0	1	0
$r_4$	0	0	1	0	1
$r_5$	1	0	0	0	0



## Lemma

Let  $M_i \in \{0, 1\}^{i \times m}$  be a submatrix of  $M$ . If  $M$  is conflict-free then  $T_i$  is a perfect phylogeny for  $M_i$ .

# Lecture Outline

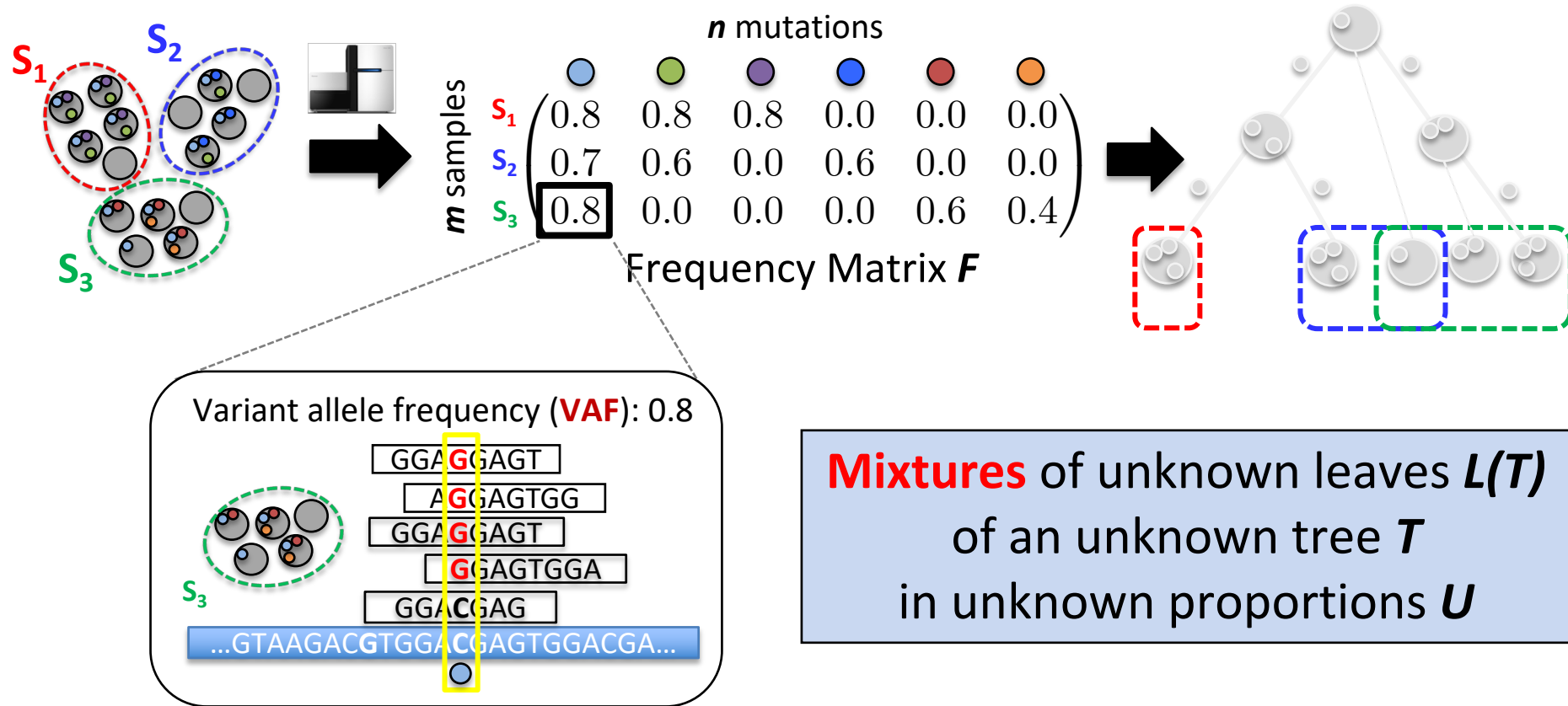
- Recap
- Maximum Parsimony
- Two-state Perfect Phylogeny
- Two-state Perfect Phylogeny Mixtures

## Reading

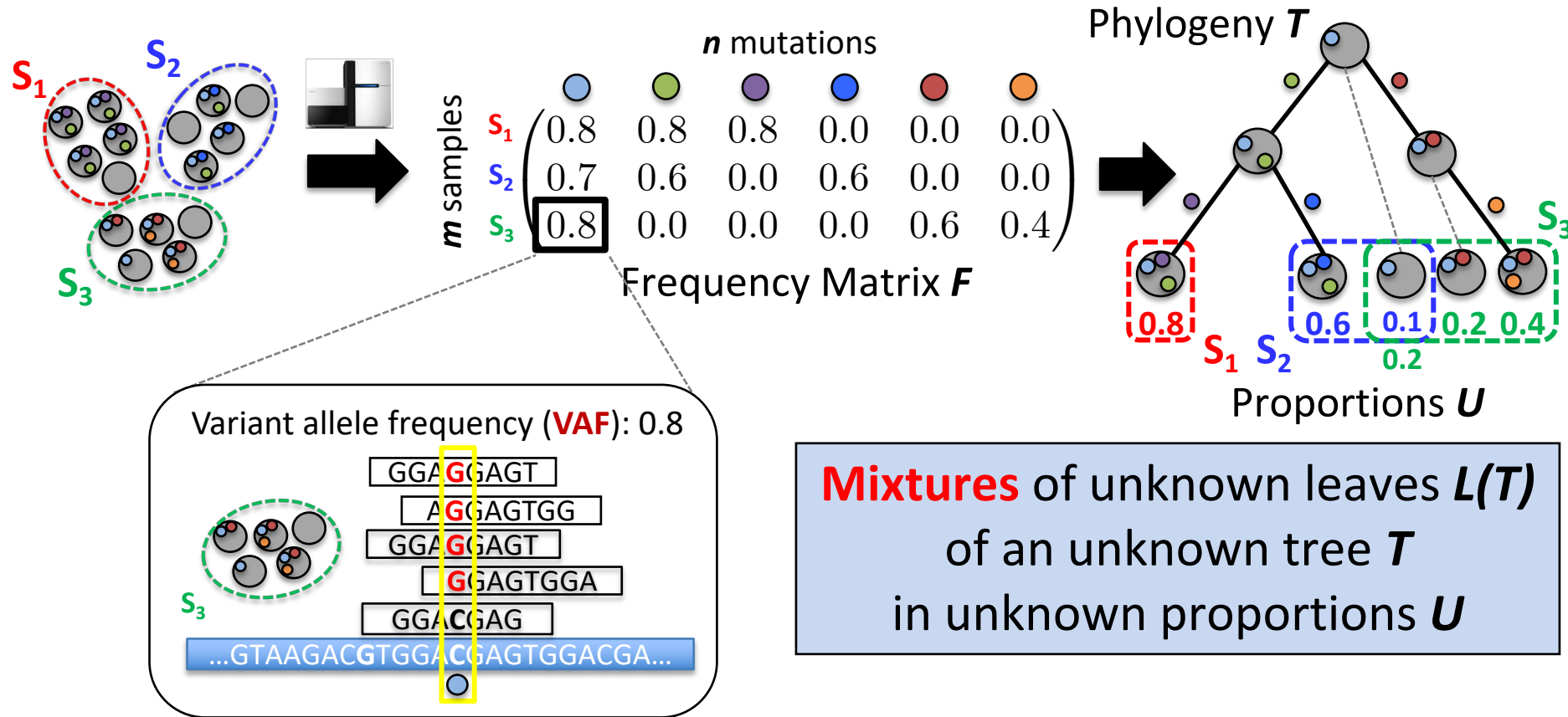
- Lecture notes



# Sequencing and Tumor Phylogeny Inference

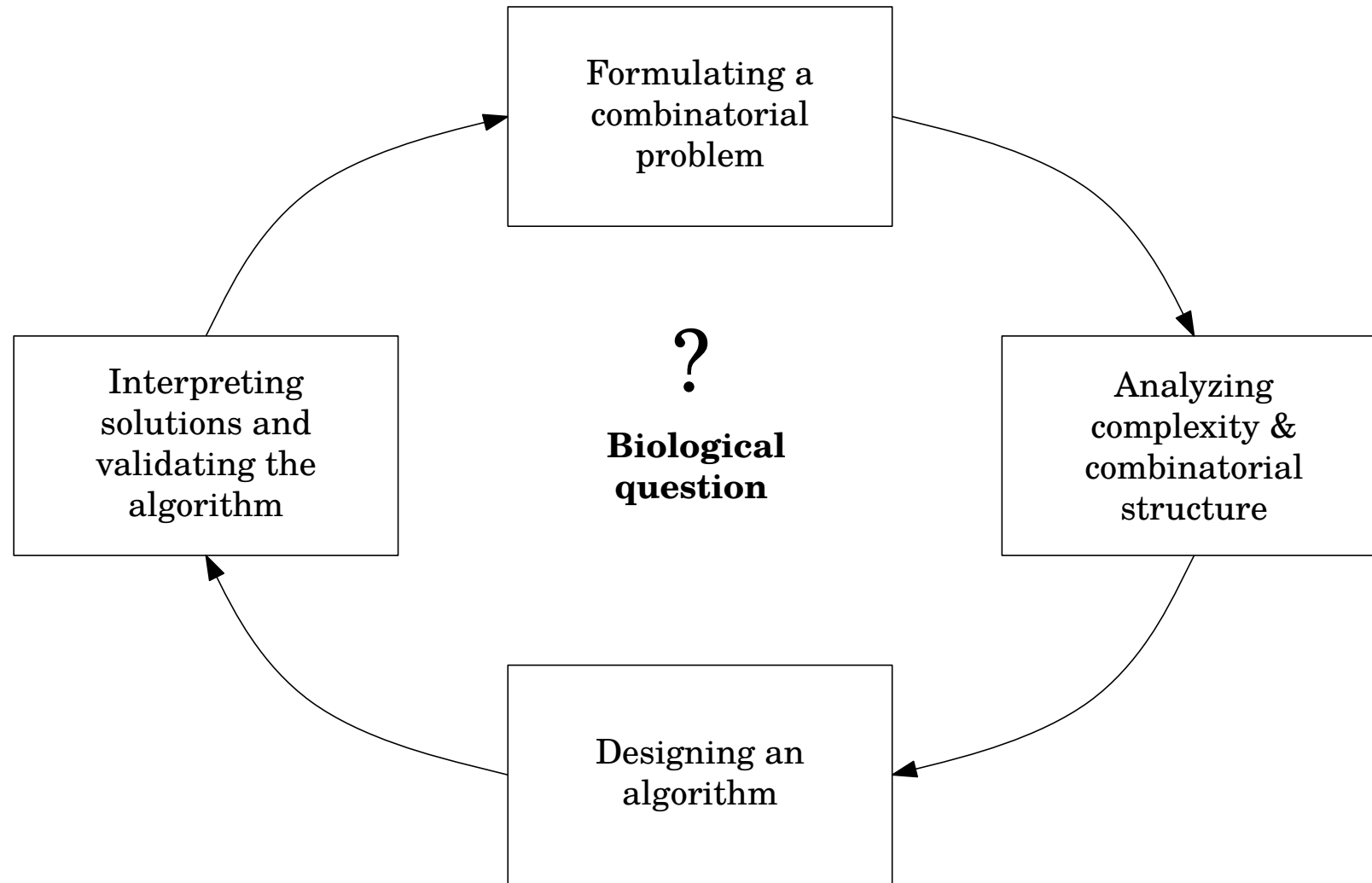


# Sequencing and Tumor Phylogeny Inference



**Tumor Phylogeny Inference:** Given frequencies  $F$ , find phylogeny  $T$  and proportions  $U$

# Key Challenge in Computational Biology

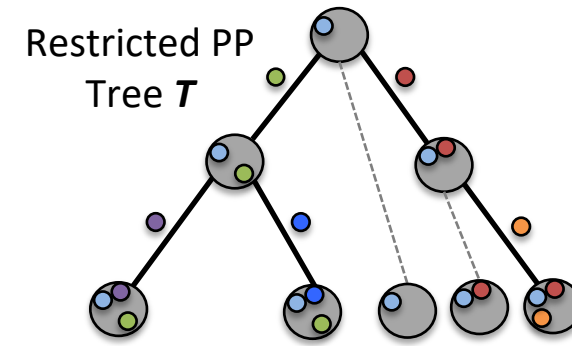


Translating a biological problem into a computational biology

# Perfect Phylogeny Mixture

## Assumptions:

- Infinite sites assumption: a character changes state once
- Error-free data



$n$  mutations

$m$  samples

Frequency Matrix  $F$

$$F = \begin{matrix} & \begin{matrix} \text{blue} & \text{green} & \text{purple} & \text{blue} & \text{red} & \text{orange} \end{matrix} \\ \begin{matrix} S_1 \\ S_2 \\ S_3 \end{matrix} & \begin{pmatrix} 0.8 & 0.8 & 0.8 & 0.0 & 0.0 & 0.0 \\ 0.7 & 0.6 & 0.0 & 0.6 & 0.0 & 0.0 \\ 0.8 & 0.0 & 0.0 & 0.0 & 0.6 & 0.4 \end{pmatrix} \end{matrix}$$

$=$

$m$  samples

Mixture Matrix  $U$

$$U = \begin{matrix} & \begin{matrix} \text{clone 1} & \text{clone 2} & \text{clone 3} & \text{clone 4} & \text{clone 5} & \text{clone 6} \end{matrix} \\ \begin{matrix} S_1 \\ S_2 \\ S_3 \end{matrix} & \begin{pmatrix} 0.0 & 0.0 & 0.8 & 0.0 & 0.0 & 0.0 \\ 0.1 & 0.0 & 0.0 & 0.6 & 0.0 & 0.0 \\ 0.2 & 0.0 & 0.0 & 0.0 & 0.2 & 0.4 \end{pmatrix} \end{matrix}$$

$n$  mutations

Restricted PP Matrix  $B$

$$B = \begin{matrix} & \begin{matrix} \text{blue} & \text{green} & \text{purple} & \text{blue} & \text{red} & \text{orange} \end{matrix} \\ \begin{matrix} \text{clone 1} \\ \text{clone 2} \\ \text{clone 3} \\ \text{clone 4} \\ \text{clone 5} \\ \text{clone 6} \end{matrix} & \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 1 & 1 \end{pmatrix} \end{matrix}$$

1-1 Equivalent

Rows of  $U$  are proportions:

$$u_{pj} \geq 0 \text{ and } \sum_j u_{pj} \leq 1$$

Perfect Phylogeny Theorem

[Estabrook, 1971]

[Gusfield, 1991]

**Perfect Phylogeny Mixture:** [El-Kebir\*, Oesper\* et al., 2015]

Given  $F$ , find  $U$  and  $B$  such that  $F = U B$

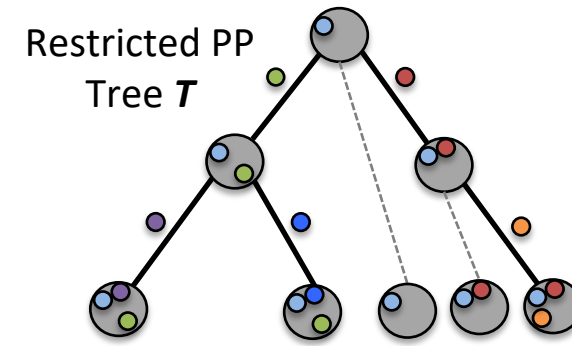
# Previous Work

## Variant of PPM:

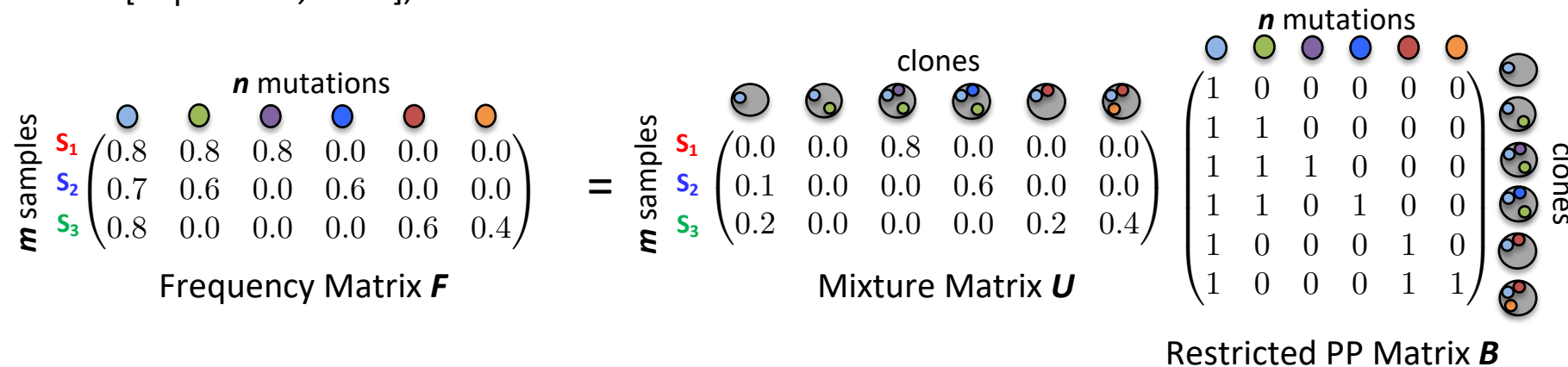
TrAp [Strino *et al.*, 2013], PhyloSub [Jiao *et al.*, 2014]

CITUP [Malikic *et al.*, 2015], BitPhylogeny [Yuan *et al.*, 2015]

LICHeE [Popic *et al.*, 2015], ...



1-1  $\updownarrow$  Equivalent



Rows of  $U$  are proportions:

$$u_{pj} \geq 0 \text{ and } \sum_j u_{pj} \leq 1$$

Perfect Phylogeny Theorem

[Estabrook, 1971]

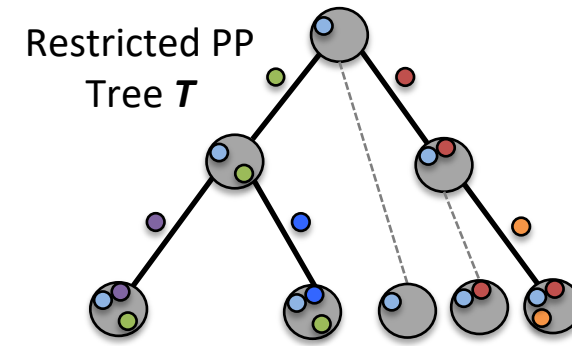
[Gusfield, 1991]

**Perfect Phylogeny Mixture:** [El-Kebir\*, Oesper\* *et al.*, 2015]

Given  $F$ , find  $U$  and  $B$  such that  $F = U B$

# Combinatorial Characterization

- Frequency  $f_{p,i}$  is mass of subtree rooted at node that introduced  $i$
- Usage  $u_{p,i}$  is mass of node that introduced  $i$



1-1  $\updownarrow$  Equivalent

$m$  samples

$n$  mutations

$$\begin{matrix} s_1 \\ s_2 \\ s_3 \end{matrix} \begin{pmatrix} 0.8 & 0.8 & 0.8 & 0.0 & 0.0 & 0.0 \\ 0.7 & 0.6 & 0.0 & 0.6 & 0.0 & 0.0 \\ 0.8 & 0.0 & 0.0 & 0.0 & 0.6 & 0.4 \end{pmatrix}$$

Frequency Matrix  $F$

=

$m$  samples

clones

$$\begin{matrix} s_1 \\ s_2 \\ s_3 \end{matrix} \begin{pmatrix} 0.0 & 0.0 & 0.8 & 0.0 & 0.0 & 0.0 \\ 0.1 & 0.0 & 0.0 & 0.6 & 0.0 & 0.0 \\ 0.2 & 0.0 & 0.0 & 0.0 & 0.2 & 0.4 \end{pmatrix}$$

Mixture Matrix  $U$

$n$  mutations

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 1 & 1 \end{pmatrix}$$

clones

Restricted PP Matrix  $B$

Rows of  $U$  are proportions:

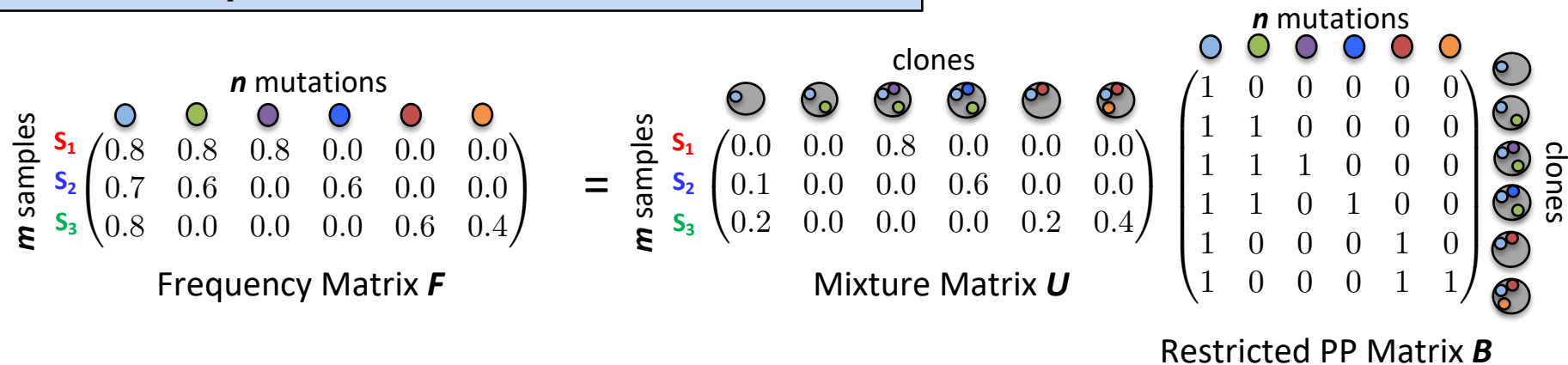
$$u_{pj} \geq 0 \text{ and } \sum_j u_{pj} \leq 1$$

**Perfect Phylogeny Mixture:** [El-Kebir\*, Oesper\* et al., 2015]

Given  $F$ , find  $U$  and  $B$  such that  $F = UB$

# Combinatorial Characterization

- Frequency  $f_{p,i}$  is mass of subtree rooted at node that introduced  $i$
- Usage  $u_{p,i}$  is mass of node that introduced  $i$



## Theorem 1:

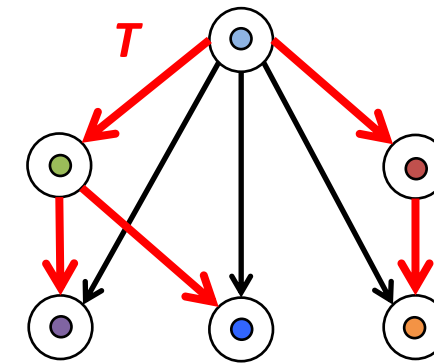
$T$  is a solution to the PPM if and only if  $T$  is a spanning tree of  $G$  satisfying the sum condition

**Perfect Phylogeny Mixture:** [El-Kebir\*, Oesper\* et al., 2015]

Given  $F$ , find  $U$  and  $B$  such that  $F = U B$

# Combinatorial Characterization

- Frequency  $f_{p,i}$  is mass of subtree rooted at node that introduced  $i$
- Usage  $u_{p,i}$  is mass of node that introduced  $i$



**G**

$m$  samples  $\begin{matrix} \text{S}_1 \\ \text{S}_2 \\ \text{S}_3 \end{matrix}$   $\begin{pmatrix} 0.8 & 0.8 & 0.8 & 0.0 & 0.0 & 0.0 \\ 0.7 & 0.6 & 0.0 & 0.6 & 0.0 & 0.0 \\ 0.8 & 0.0 & 0.0 & 0.0 & 0.6 & 0.4 \end{pmatrix}$   $n$  mutations

Frequency Matrix **F**

$m$  samples  $\begin{matrix} \text{S}_1 \\ \text{S}_2 \\ \text{S}_3 \end{matrix}$   $\begin{pmatrix} 0.0 & 0.0 & 0.8 & 0.0 & 0.0 & 0.0 \\ 0.1 & 0.0 & 0.0 & 0.6 & 0.0 & 0.0 \\ 0.2 & 0.0 & 0.0 & 0.0 & 0.2 & 0.4 \end{pmatrix}$  clones

Mixture Matrix **U**

$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 1 & 1 \end{pmatrix}$   $n$  mutations

Restricted PP Matrix **B**

## Theorem 1:

$T$  is a solution to the PPM if and only if  $T$  is a spanning tree of **G** satisfying the sum condition

## Theorem 2:

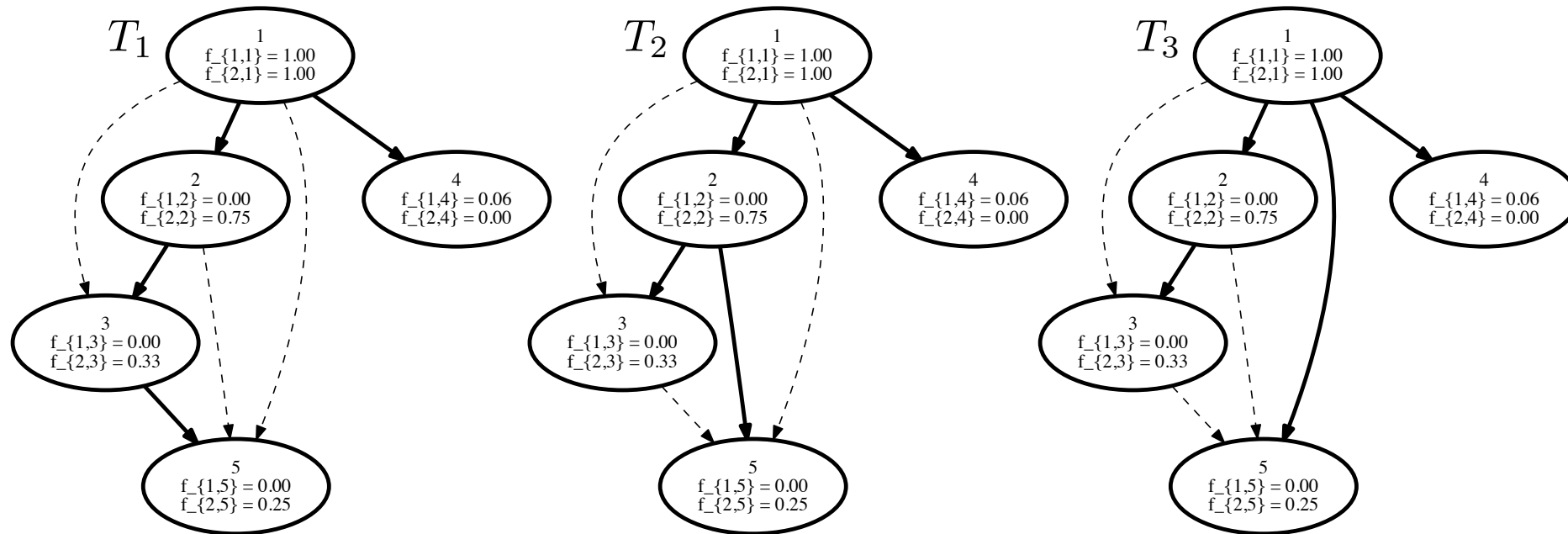
PPM is NP-complete even for  $m=2$

**Perfect Phylogeny Mixture:** [El-Kebir\*, Oesper\* et al., 2015]

Given **F**, find **U** and **B** such that  $F = U B$



# Non-uniqueness of Solutions to PPM



$$F = \begin{pmatrix} 1 & 0 & 0 & 0.06 & 0 \\ 1 & 0.75 & 0.33 & 0 & 0.25 \end{pmatrix}$$

**Question 1:** Can we determine the number of solutions?

**Question 2:** Can sample solutions uniformly at random?

# Summary of Lectures 1 & 2

- DNA, RNA and proteins are sequences
  - Central dogma of molecular biology: DNA -> RNA -> protein
- Problem != algorithm
- Key challenge in computational biology is translating a biological problem into a computational problem
- Cancer is a genetic disease caused by somatic mutations
- Inter-tumor heterogeneity and intra-tumor heterogeneity:
  - *Not only is every tumor different, but so is every tumor cell...*
- Reading:
  - “Biology for Computer Scientists” by Lawrence Hunter  
([http://www.el-kebir.net/teaching/CS466/Hunter\\_BIO\\_CS.pdf](http://www.el-kebir.net/teaching/CS466/Hunter_BIO_CS.pdf))