CS 598MEB Computational Cancer Genomics Lecture 3

Mohammed El-Kebir

January 28, 2020



Course Project

- 1-2 students per project
- First write a proposal, which will receive feedback from instructor and fellow students
- Then, conduct research and write a paper
- Pick venue (conference/journal) and use LaTeX style for your paper

Lecture Outline

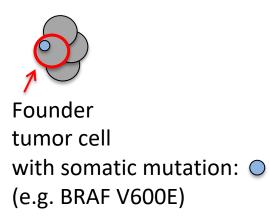
- Recap
- Two-state Perfect Phylogeny Mixtures

Reading

- M. El-Kebir, L. Oesper, H. Acheson-Field and B. J. Raphael. Reconstruction of clonal trees and tumor composition from multi-sample sequencing data. <u>Bioinformatics (Special Issue: Proceedings of ISMB), 31(12):i62-i70,</u> 2015
- Y. Qi, D. Pradhan and M. El-Kebir. Implications of non-uniqueness in phylogenetic deconvolution of bulk DNA samples of tumors. <u>Algorithms for</u> <u>Molecular Biology</u>, 14:19, 2019.

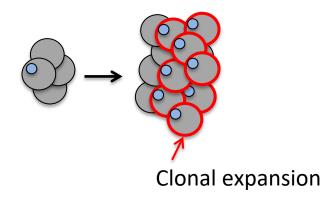
Tumorigenesis: Cell Mutation

Clonal Evolution Theory of Cancer [Nowell, 1976]



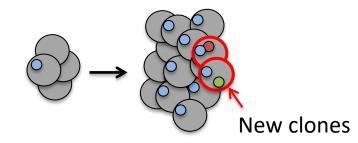
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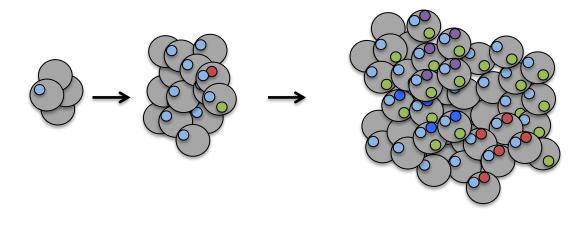
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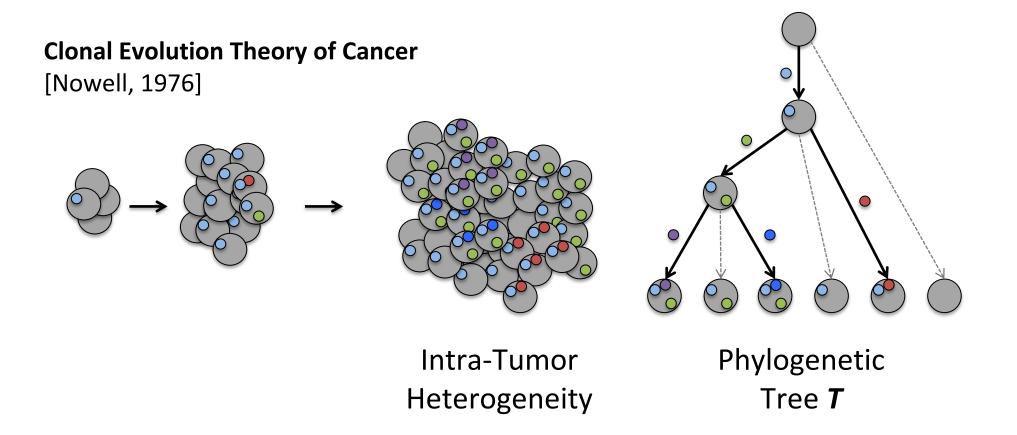
Tumorigenesis: Cell Mutation & Division

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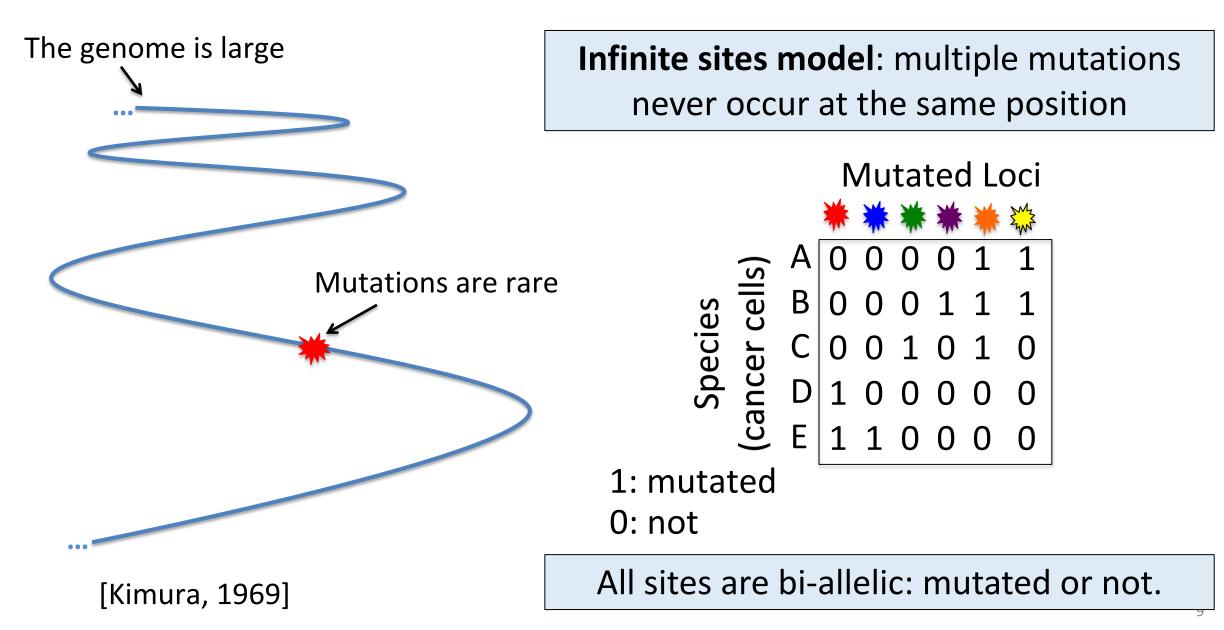


Intra-Tumor Heterogeneity

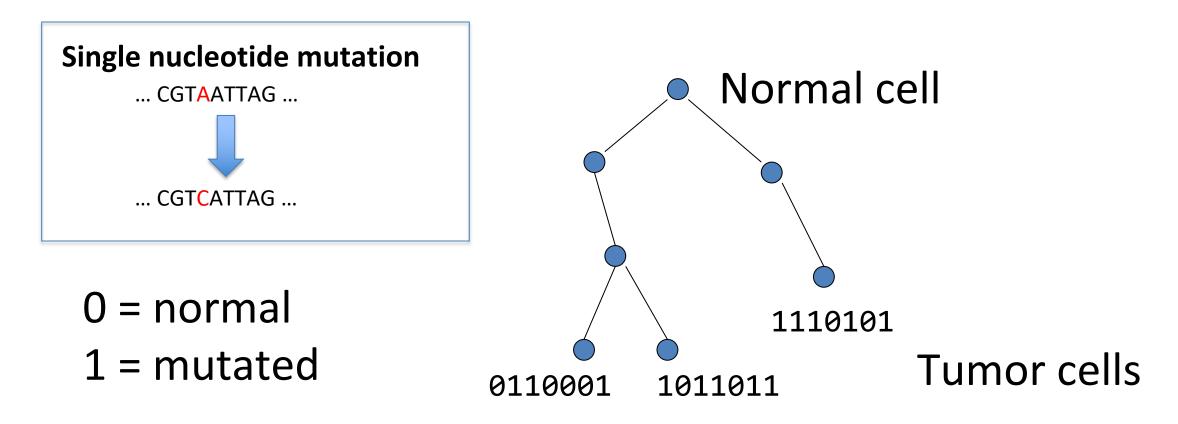
Tumorigenesis: Cell Mutation & Division



Infinite Sites Model



Progression of Somatic Mutations



Root is the normal, founder cell and leaves are cells in tumor.

Infinite sites assumption: each locus mutates only once.

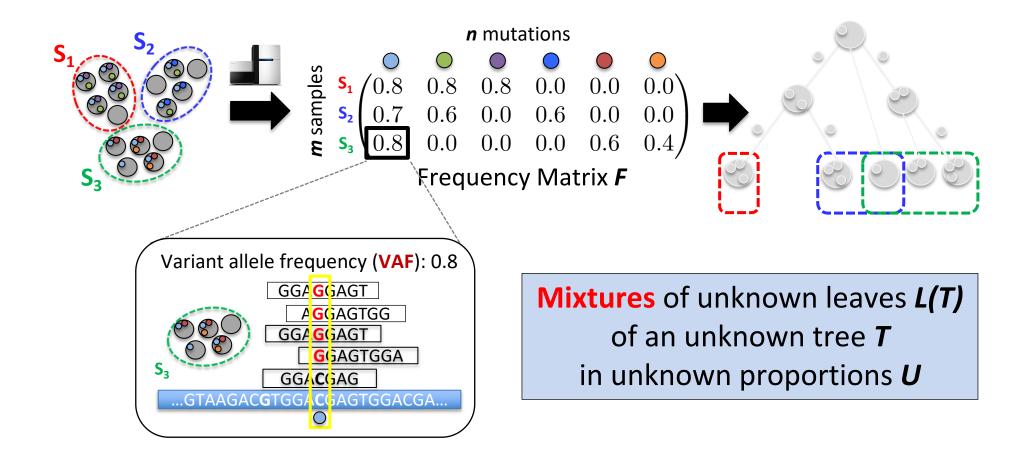
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- Recap
- Two-state Perfect Phylogeny Mixtures

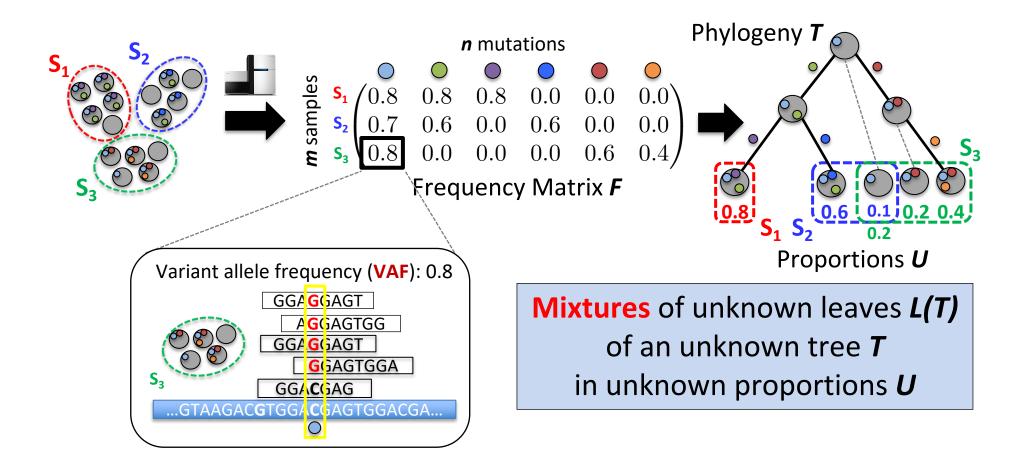
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Sequencing and Tumor Phylogeny Inference

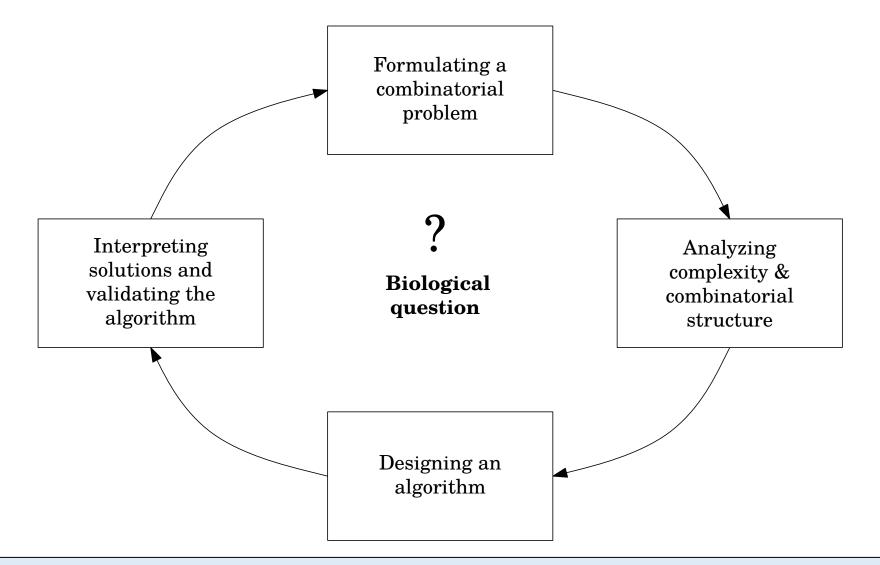


Sequencing and Tumor Phylogeny Inference

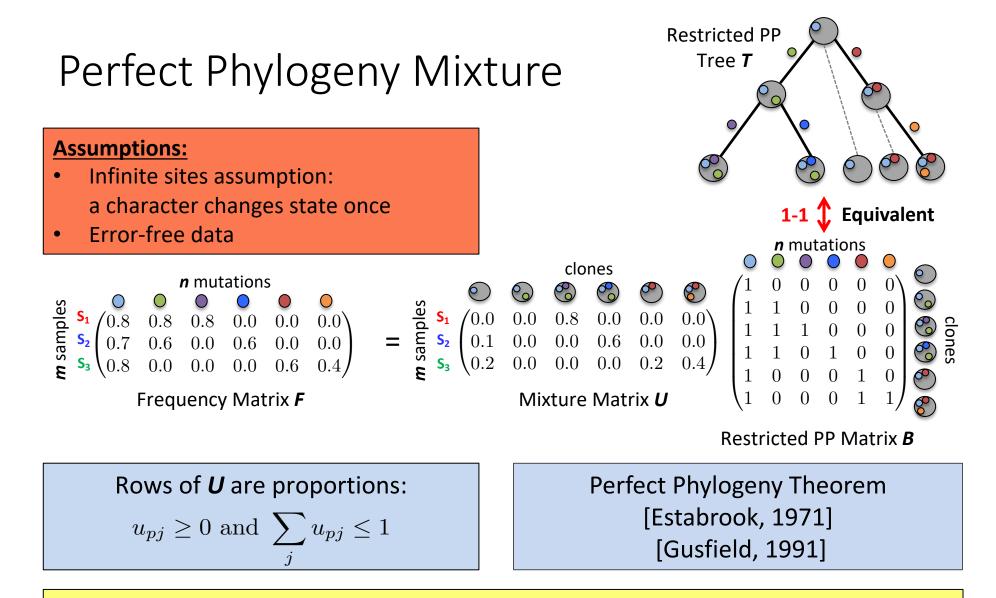


Tumor Phylogeny Inference: Given frequencies *F*, find phylogeny *T* and proportions *U*

Key Challenge in Computational Biology



Translating a biological problem into a computational biology



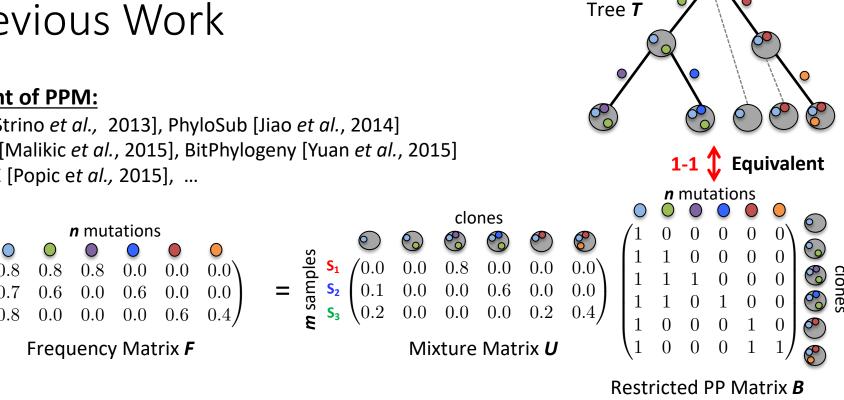
Perfect Phylogeny Mixture: [El-Kebir*, Oesper* et al., 2015] Given *F*, find *U* and *B* such that *F* = *U B*

Previous Work

Variant of PPM:

n samples

TrAp [Strino et al., 2013], PhyloSub [Jiao et al., 2014] CITUP [Malikic et al., 2015], BitPhylogeny [Yuan et al., 2015] LICHeE [Popic et al., 2015], ...



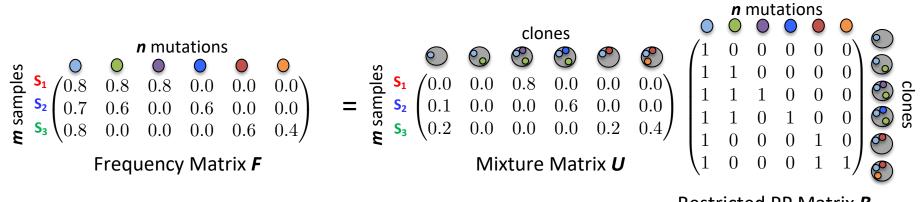
Restricted PP

Rows of **U** are proportions: Perfect Phylogeny Theorem [Estabrook, 1971] $u_{pj} \ge 0$ and $\sum u_{pj} \le 1$ [Gusfield, 1991]

Perfect Phylogeny Mixture: [El-Kebir*, Oesper* et al., 2015] Given **F**, find **U** and **B** such that **F** = **U B**

Combinatorial Characterization

Perfect Phylogeny Mixture: [El-Kebir*, Oesper* et al., 2015] Given *F*, find *U* and *B* such that *F* = *U B*

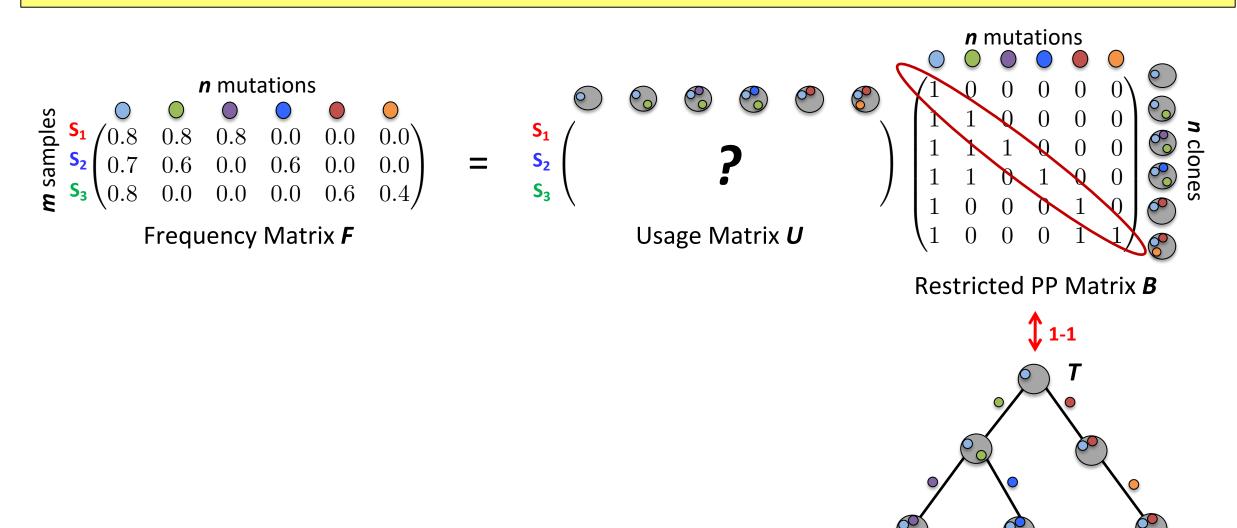


Restricted PP Matrix **B**

- Combinatorial characterization involves investigating what (optimal) solutions look like
- This starts by asking questions!

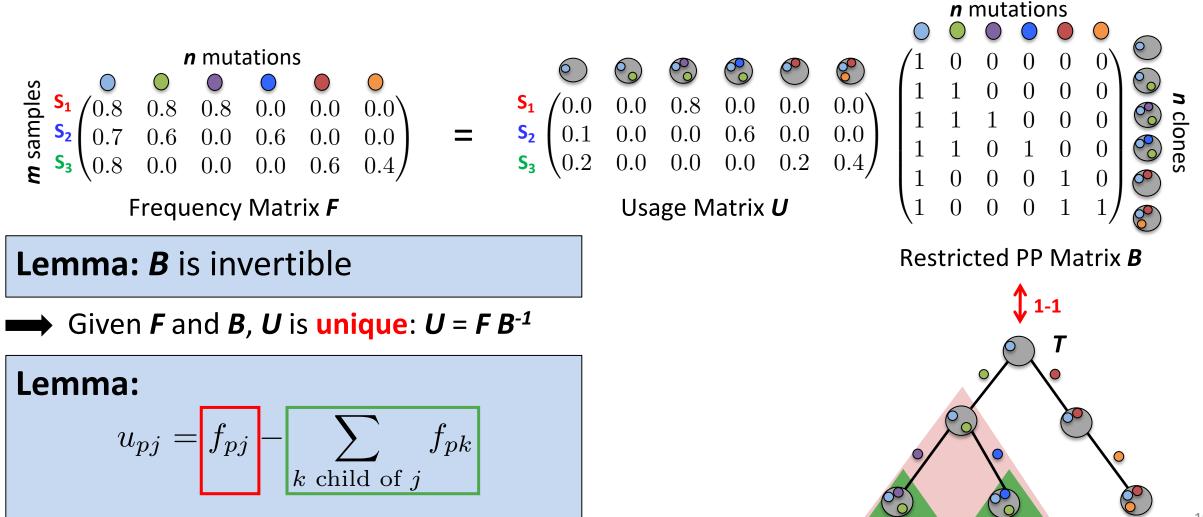
Given F and T (or B), is there a usage matrix U?

PPM: Given *F*, find *U* and *B* such that *F* = *U B*



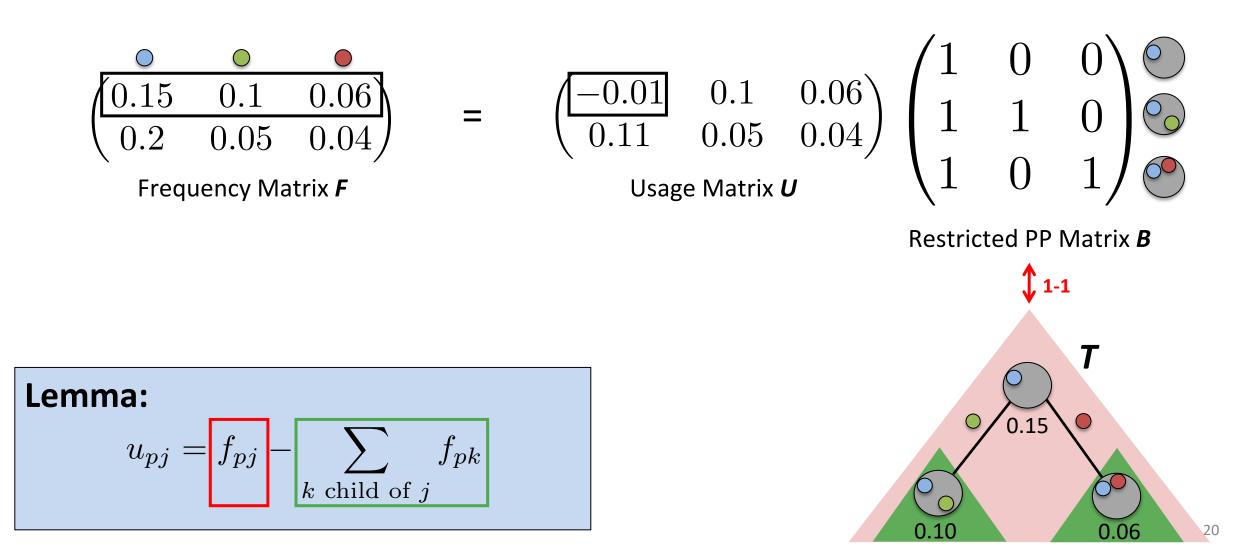
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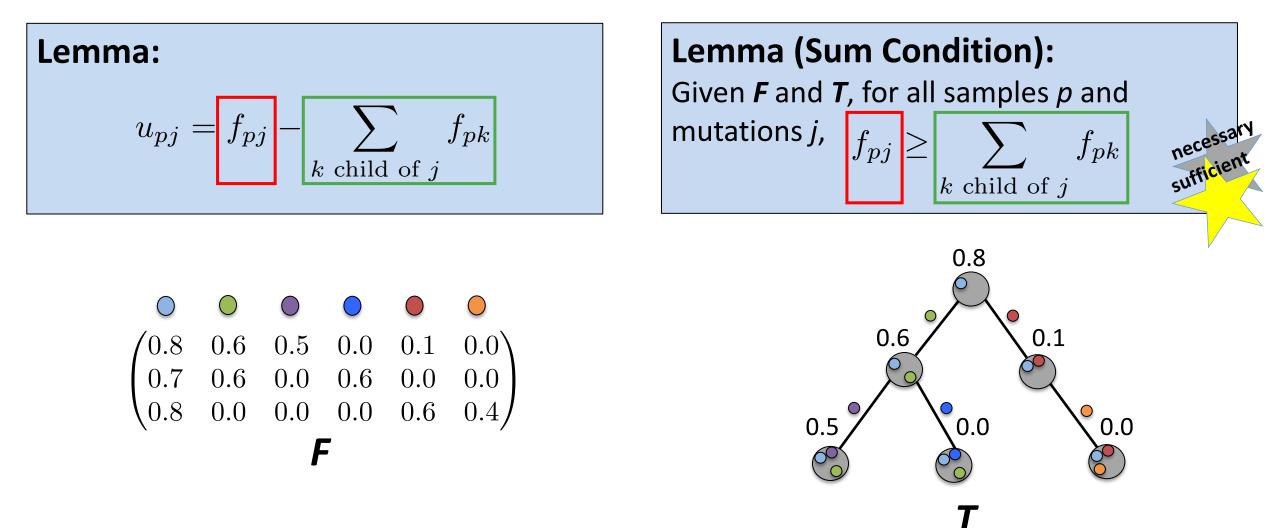
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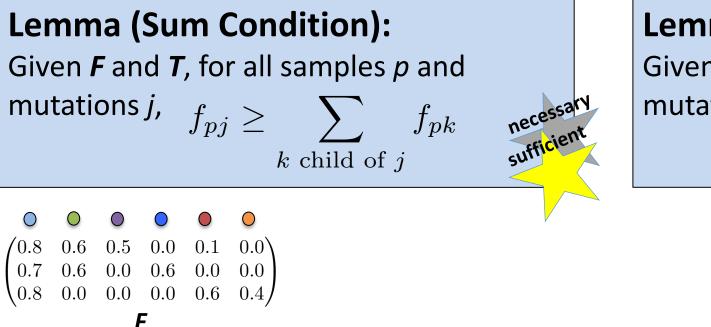


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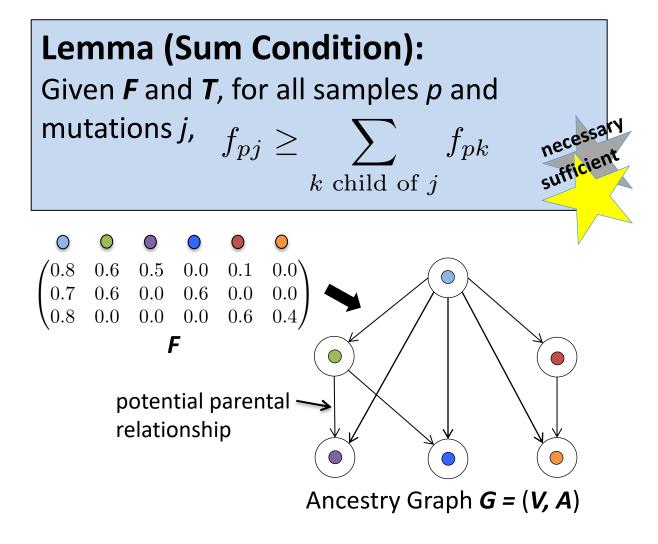






Lemma (Ancestry Condition): Given *F* and *T*, for all samples *p* and mutations *k* child of *j*, $f_{pj} \ge f_{pk}$

necessal



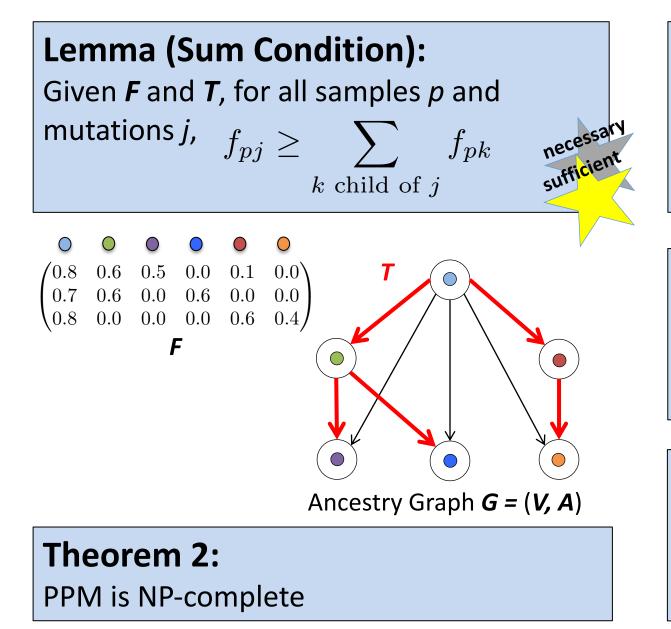
Lemma (Ancestry Condition): Given **F** and **T**, for all samples *p* and mutations *k* child of *j*, $f_{pj} \ge f_{pk}$

Ancestry graph G = (V, A); given F

- Vertex for every mutation
- Edge $(j,k) \in A$ iff $f_{pj} \geq f_{pk}$

for all samples p

necessal



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Theorem 1:

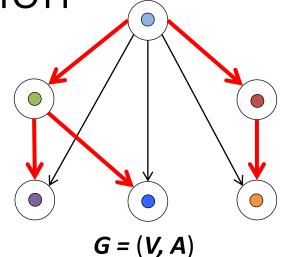
T is a solution to the PPM if and only if*T* is a spanning tree of *G* satisfying the SumCondition

Solving the PPM problem: ILP formulation

max $\sum x_{jk}$ $(v_j, v_k) \in A'$ s.t. $\sum x_{rj} = 1$ $v_i \in \delta^+(v_r)$ $x_{kl} \leq \sum x_{jk}$ $v_j \in \delta^-(v_k)$ $\sum \quad x_{jk} \le 1$ $v_j \in \delta^-(v_k)$

Find the largest set of edges in *G*

Exactly one root node



 $\forall (v_k, v_l) \in A$ Connectivity

 $\forall v_k \in V$ Tree

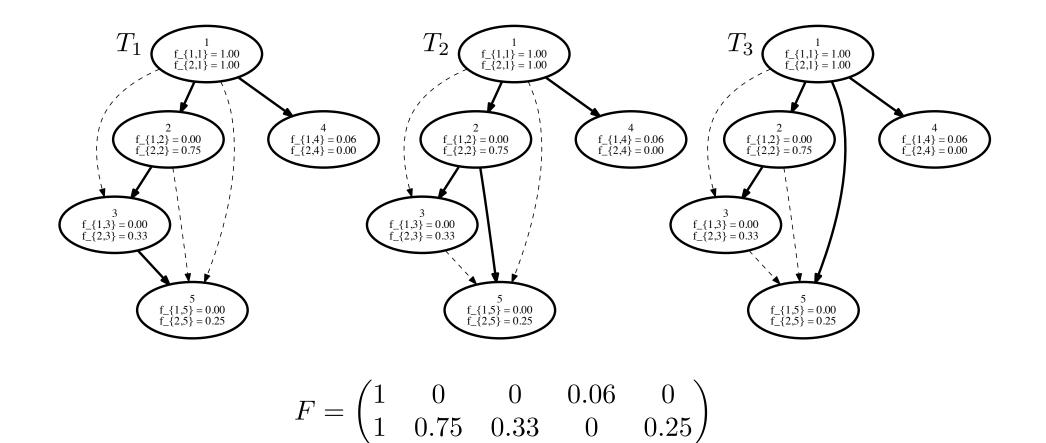
 $\sum_{v_j \in \delta^-(v_k)} f_{pk} x_{jk} \ge \sum_{v_l \in \delta^+(v_k)} f_{pl} x_{kl} \quad \forall p \in [m], v_k \in V \quad \text{Sum condition}$ $x_{jk} \in \{0, 1\} \qquad \qquad \forall (v_j, v_k) \in A'$

Non-uniqueness of Solutions to PPM

$$F = \begin{pmatrix} 1 & 0 & 0 & 0.06 & 0 \\ 1 & 0.75 & 0.33 & 0 & 0.25 \end{pmatrix}$$

Question 0: Reconstruct all solutions?

Non-uniqueness of Solutions to PPM



Question 1: Can we determine the number of solutions?

Question 2: Can sample solutions uniformly at random?

Question 3: How to enumerate solutions?

Recall: Different Types of Problems!

Problem Π with instance *X* and solution set $\Pi(X)$:

- Decision problem:
 - Is $\Pi(X) = \emptyset$?
- Optimization problem:
 - Find $y^* \in \Pi(X)$ s.t. $f(y^*)$ is optimum.
- Counting problem:
 - Compute $|\Pi(X)|$.
- Sampling problem:
 - Sample uniformly from $\Pi(X)$.
- Enumeration problem:
 - Enumerate all solutions in $\Pi(X)$

Algorithms:

Set of instructions for solving problem.

- Exact
- Heuristic

On the Complexity of #PPM (new results)

Question 1: Can we determine the number of solutions?

Question 2: Can sample solutions uniformly at random?

#PPM: Given *F*, count the number of pairs *(U, B)* composed of mixture matrix *U* and perfect phylogeny matrix *B* such that *F* = *U B*

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#P is the complexity class of counting problems whose decision problems are in NP

Every problem in #P can be reduced in polynomial time to any problem in #P-complete, preserving cardinalities

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Theorem: #PPM is #P-complete

Theorem: There is no FPRAS for #PPM

Theorem: There is no FPAUS for PPM



Yuanyuan Qi

Outline

Background and theory:

- Perfect Phylogeny Mixture (PPM) problem
- Combinatorial characterization of solutions
- #PPM: exact counting and uniform sampling

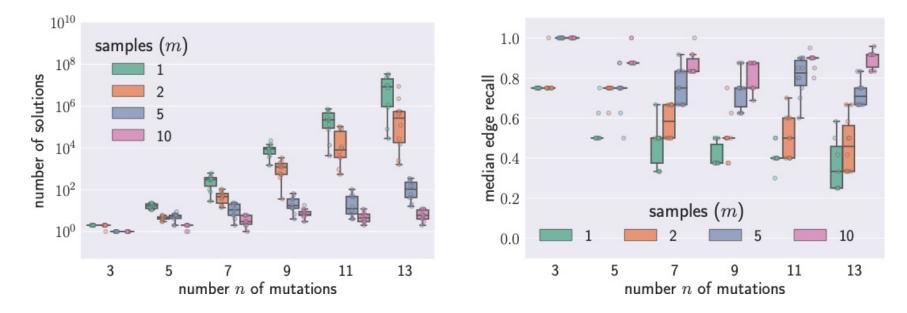
Simulation results:

- What contributes to non-uniqueness?
- How to reduce non-uniqueness?
- How does non-uniqueness affect current methods?

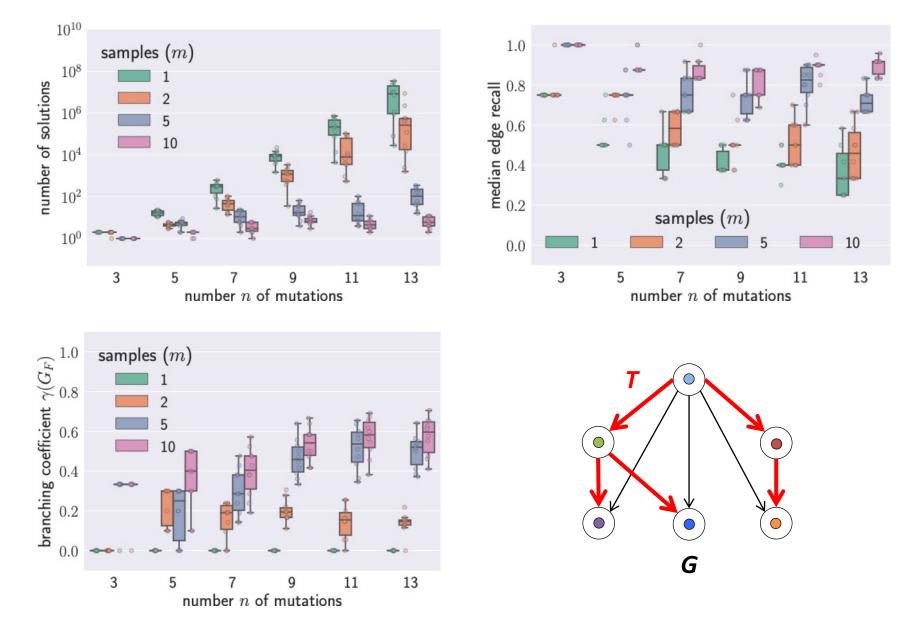


Dikshant Pradhan

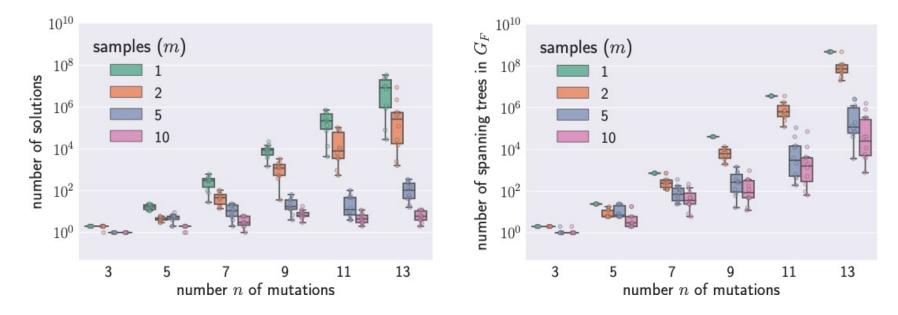
What Contributes to Non-uniqueness?

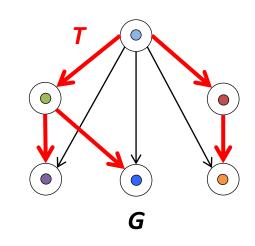


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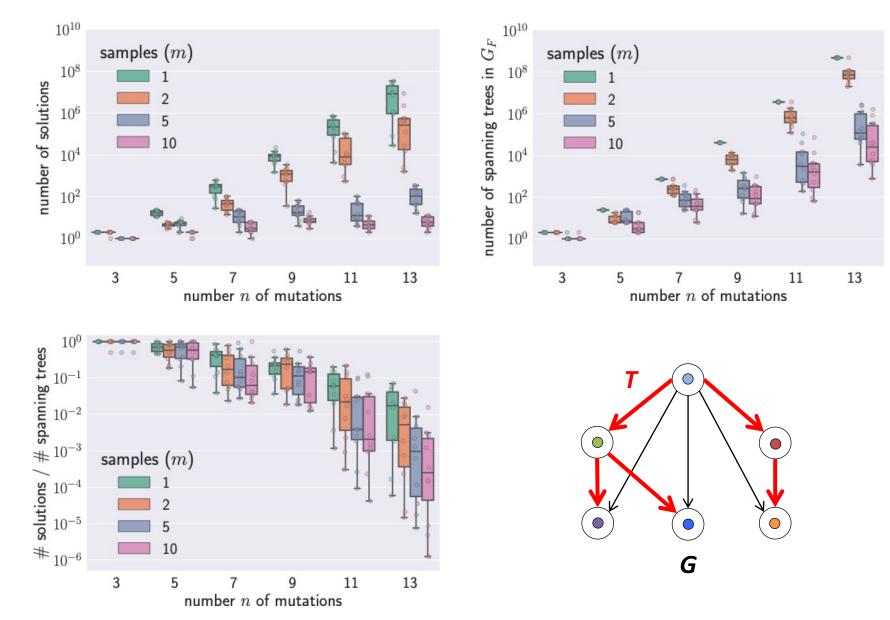


An Upper Bound for Number of Solutions

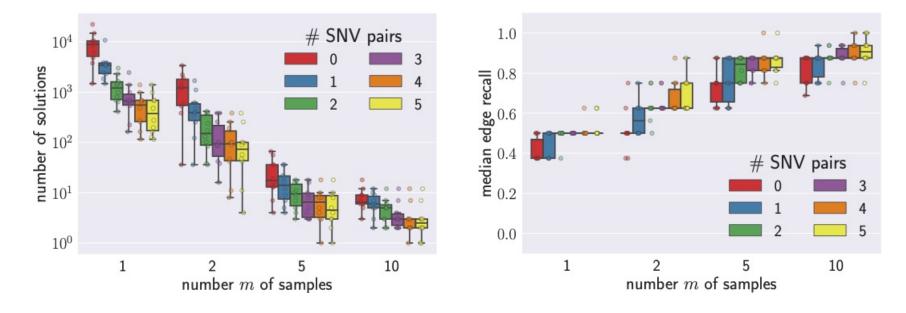


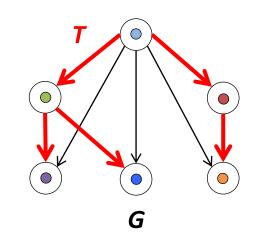


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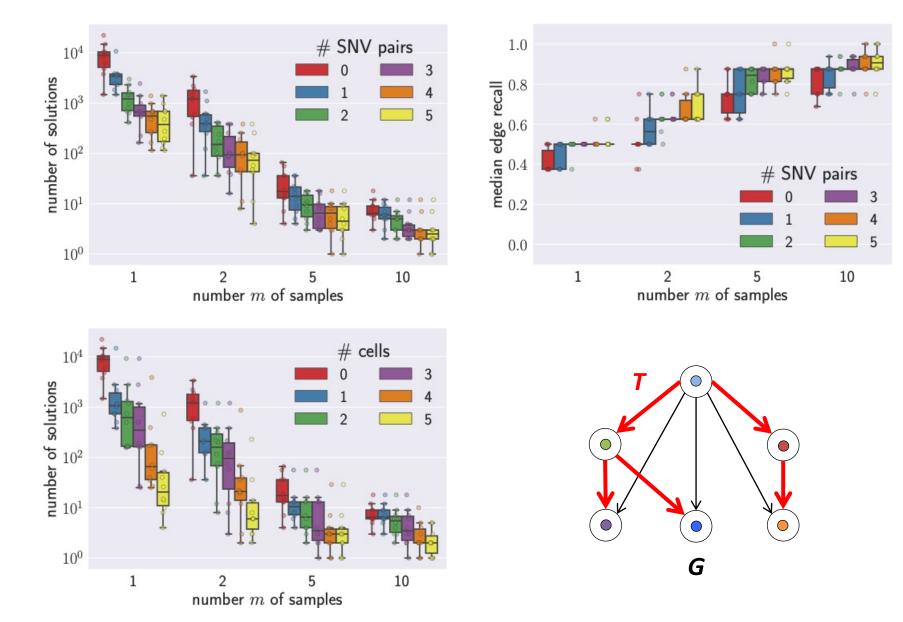


How to Reduce Non-Uniqueness?

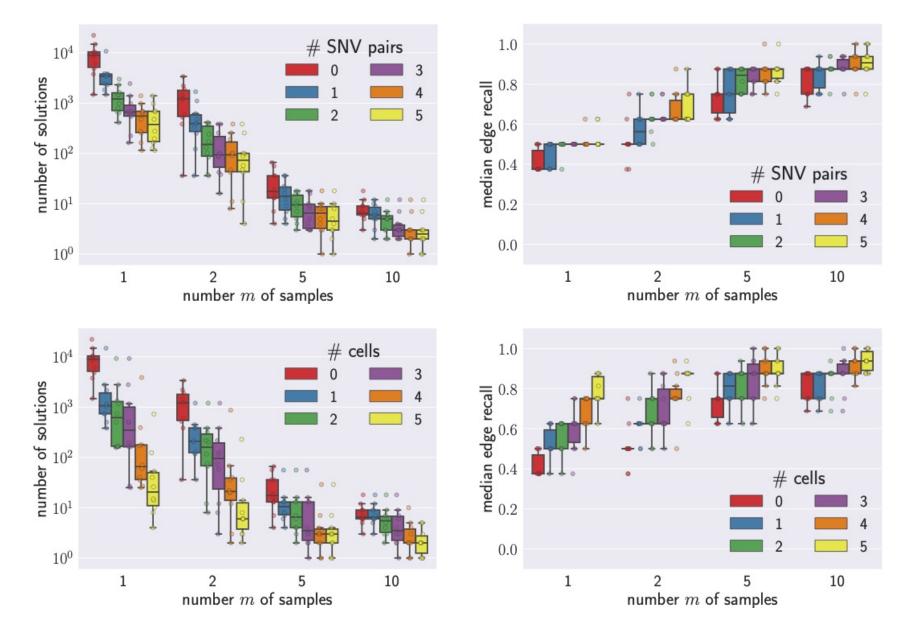




How to Reduce Non-Uniqueness?



How to Reduce Non-Uniqueness?

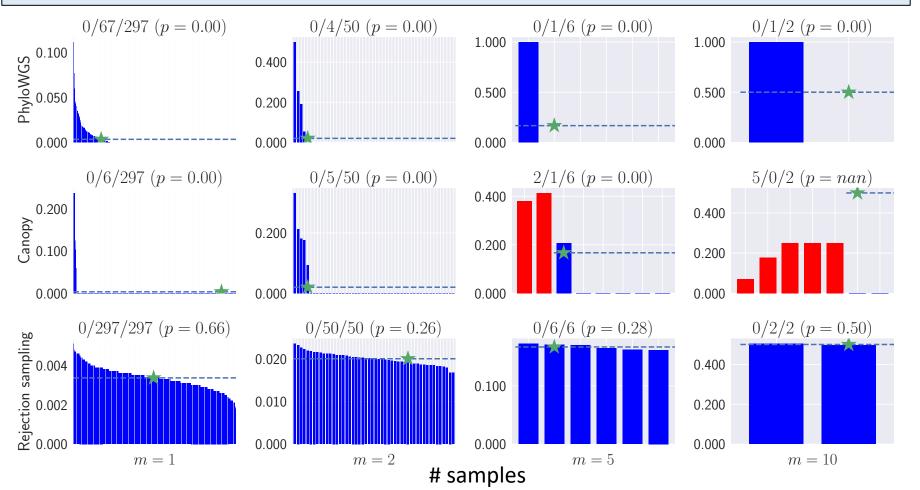


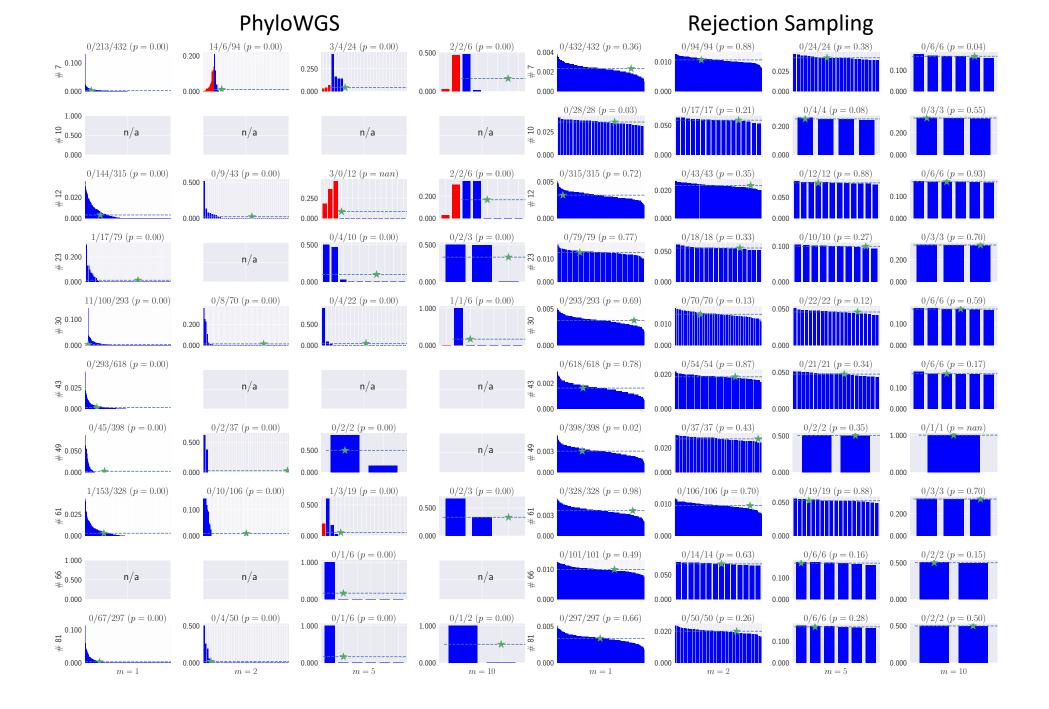
How Does Non-uniqueness affect Methods?

Two current MCMC methods using default parameters:

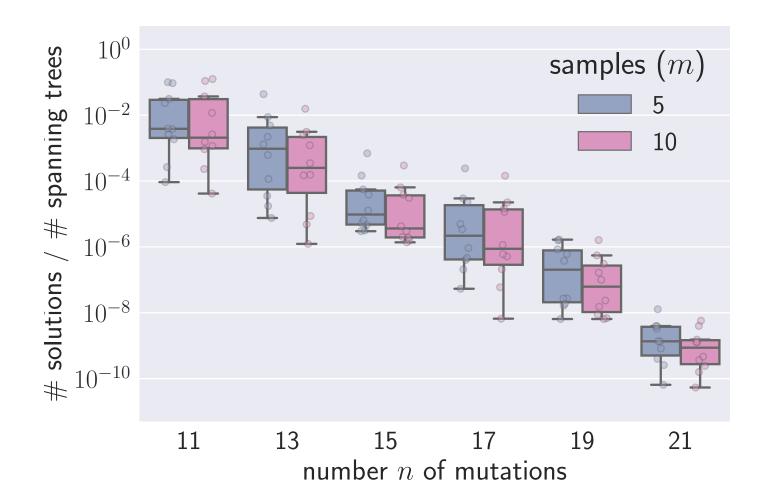
- PhyloWGS, Deshwar et al., Genom. Biol., 2015 [10,000 samples]
- Canopy, Jiang et al., PNAS, 2016

[~300 samples]

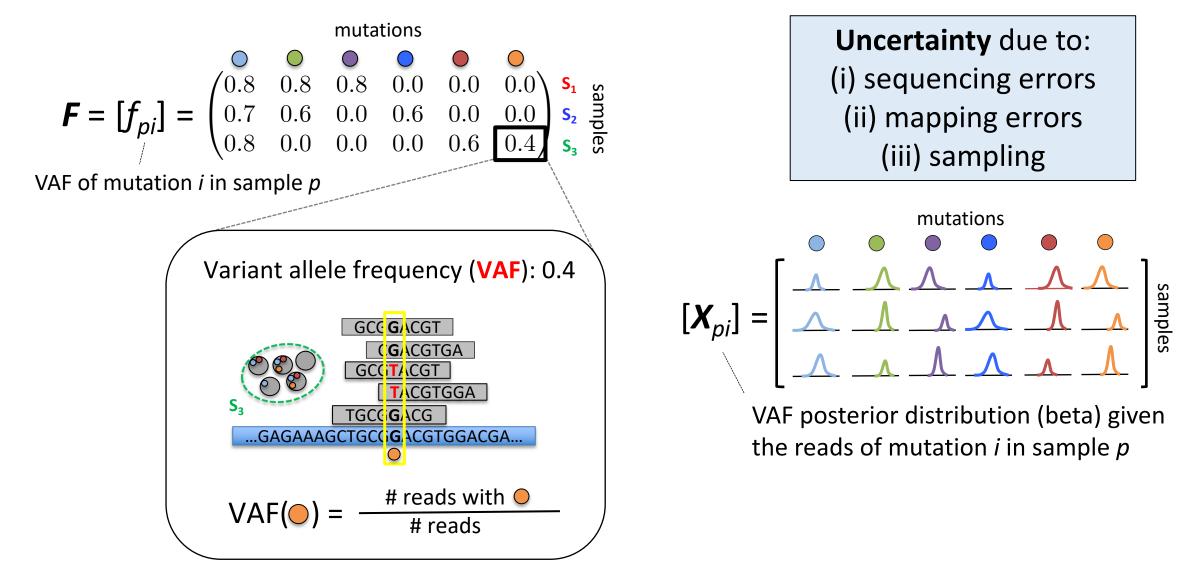




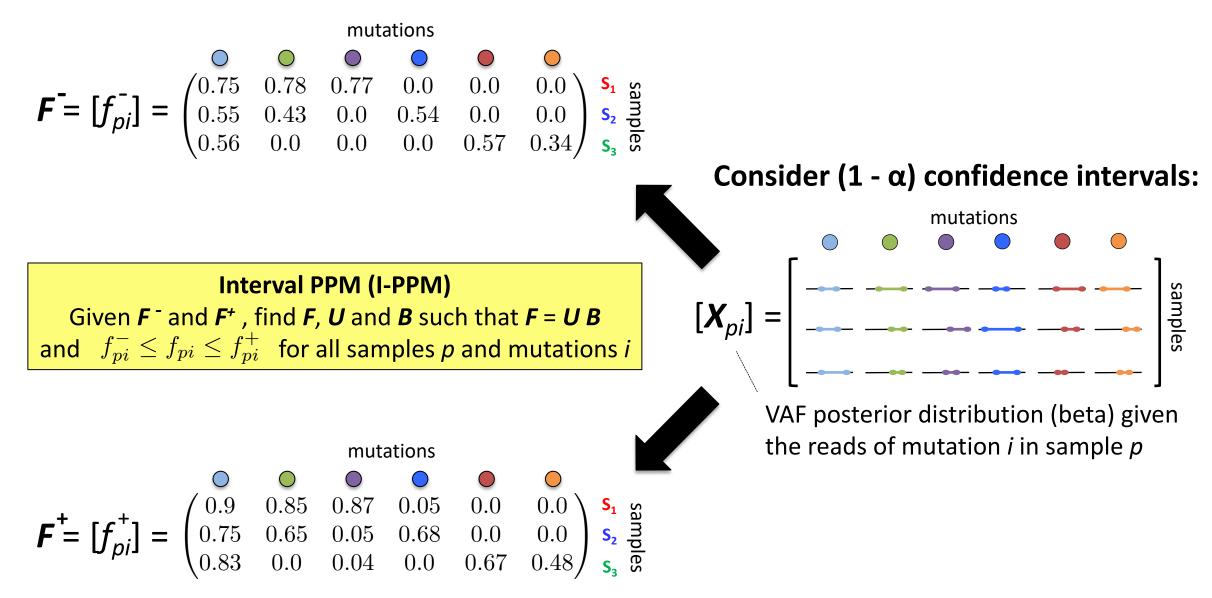
Rejection Sampling Does Not Scale



Probabilistic Model for Noisy Measurements

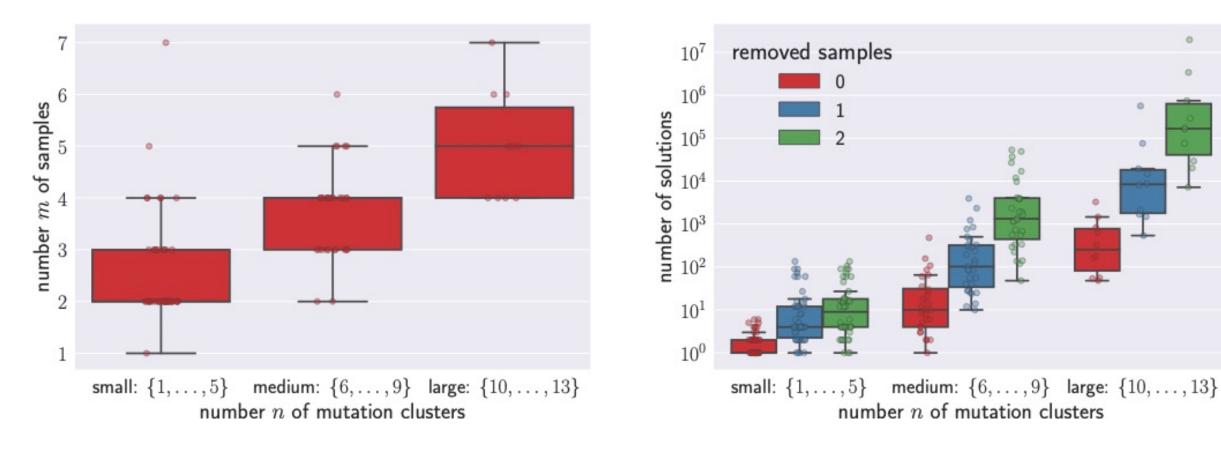


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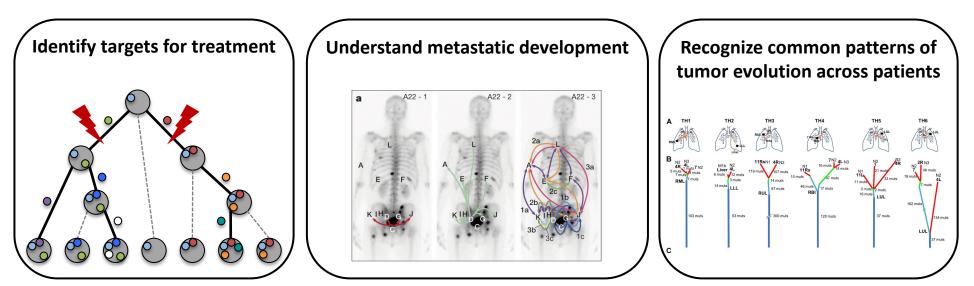


Real Data

- Cohort of 100 lung cancers [Jamal-Hanjani, NEJM 2017]
- 90% confidence intervals



Challenges



Downstream analyses in cancer genomics **critically rely** on accurate tumor phylogeny inference

Challenge I

Novel algorithms that sample uniformly at random from the space of PPM solutions

Challenge II

Algorithms to accurately summarizing solution space (consensus trees)

Conclusion

Background and theory:

- Perfect Phylogeny Mixture (PPM) problem
- Combinatorial characterization of solutions
- #PPM: exact counting and uniform sampling

Simulation results:

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Summary of Lectures 1, 2 and 3

- DNA, RNA and proteins are sequences
 - Central dogma of molecular biology: DNA -> RNA -> protein
- Problem != algorithm
- Key challenge in computational biology is translating a biological problem into a computational problem
- Cancer is a genetic disease caused by somatic mutations
- Inter-tumor heterogeneity and intra-tumor heterogeneity:
 - Not only is every tumor different, but so is every tumor cell...
- Non-uniqueness of solutions in phylogeny reconstruction from bulk DNA samples